

# Discrete Vector Fields and Effective Homology

```
;; Clock  
Computing  
<TnPr <TnPr  
End of computing.  
  
;; Clock -> 2002-01-17, 19h 25m 36s.  
Computing the boundary of the generator 19 (dimension 7) :  
<TnPr <TnPr <TnPr S3 <<Abar[2 S1][2 S1]>>> <<Abar>>> <<Abar>>>  
End of computing.
```

Homology in dimension 6 :

Component Z/12Z

---done---

```
;; Clock -> 2002-01-17, 19h 27m 15s
```

*Francis Sergeraert, Institut Fourier, Grenoble  
Joint work with Ana Romero, Universidad de La Rioja  
Jagiellonian University, Krakow, 12 May 2016*

## Semantics of colours:

**Blue** = “Standard” Mathematics

**Red** = Constructive, effective,  
algorithm, machine object, ...

**Violet** = Problem, difficulty,  
obstacle, disadvantage, ...

**Green** = Solution, essential point,  
mathematicians, ...

## Plan.

- ● Introduction.
- Discrete vector fields.
- Homological Reductions.
- Product problem in Combinatorial Topology.
- Discrete Vector Field for Products.
- Free generalization to twisted products.
- Effective Eilenberg-Moore spectral sequences.

## Introduction.

Algebraic Topology is a translator:



# 1. Introduction.

Algebraic Topology is a translator:



Serre (1950): Up to homotopy

any **map** can be transformed into a **fibration**.

**Fibration** = **Twisted Product**

<b>Topology</b>	$\mapsto$	<b>Algebra</b>
<b>Product</b>	$\mapsto$	<b>Eilenberg-Zilber Theorem</b>
<b>Twisted product</b>	$\mapsto$	<b>Serre Spectral Sequence</b>

**Discrete vector fields**

$\Rightarrow$  **New understanding** of the **Eilenberg-Zilber Theorem**

$\Rightarrow$  An **effective** version of the **Serre Spectral Sequence**

as a **direct** consequence of this version of **Eilenberg-Zilber**.

Example: **Rubio-Morace** homotopy for **Eilenberg-Zilber**:

$$RM : C_*(X \times Y) \rightarrow C_{*+1}(X \times Y)$$

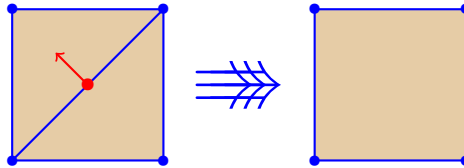
$$RM(x_p \times y_p) = \sum_{\substack{0 \leq r \leq p-1 \\ 0 \leq s \leq p-r-1 \\ (\eta, \eta') \in \text{Sh}(s+1, r)}} (-1)^{p-r-s} \varepsilon(\eta, \eta') \dots$$

$$\dots (\uparrow^{p-r-s}(\eta') \eta_{p-r-s-1} \partial_{p-r+1} \dots \partial_p x_p \times \uparrow^{p-r-s}(\eta) \partial_{p-r-s} \dots \partial_{p-r-1} y_p)$$

with  $\text{Sh}(p, q) = \{(p, q)\text{-shuffles}\} = \{(\eta_{i_{p-1}} \dots \eta_{i_0}, \eta_{j_{q-1}} \dots \eta_{j_0})\}$   
 for  $0 \leq i_0 < \dots < i_{p-1} \leq p + q - 1$   
 and  $0 \leq j_0 < \dots < j_{q-1} \leq p + q - 1$   
 and  $\{i_0, \dots, i_{p-1}\} \cap \{j_0, \dots, j_{q-1}\} = \emptyset$ .

and  $\uparrow^k (\eta_\alpha \eta_\beta \dots) = \eta_{\alpha+k} \eta_{\beta+k} \dots$  ( $\uparrow^k = k\text{-shift operator.}$ )

Simpler:



once the notion of **discrete vector field** is understood.



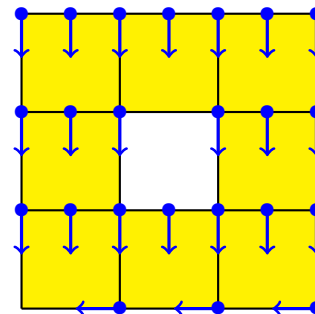
## 2. Discrete Vector Fields.

Definition: A **Discrete Vector Field** is a pairing:

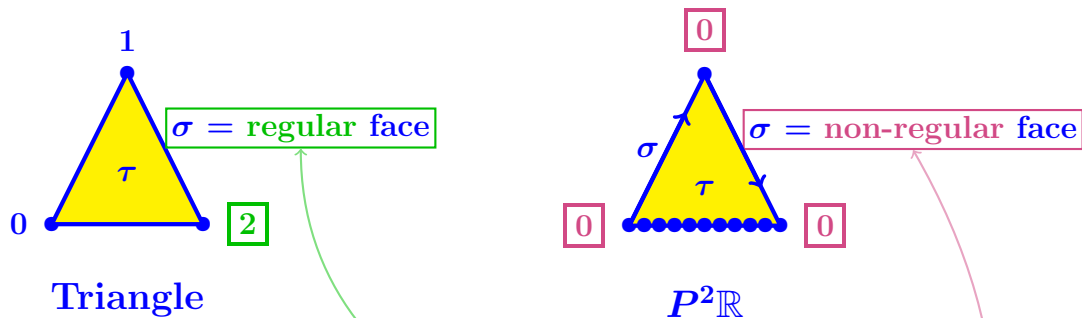
$$V = \{(\sigma_i, \tau_i)\}_{i \in v}$$

satisfying:

- $\forall i \in v$ ,  $\tau_i$  = some  $k_i$ -cell and  $\sigma_i$  = some  $(k_i - 1)$ -cell.
- $\forall i \in v$ ,  $\sigma_i$  is a **regular** face of  $\tau_i$ .
- $\forall i \neq j \in v$ ,  $\{\sigma_i, \tau_i\} \cap \{\sigma_j, \tau_j\} = \emptyset$ .
- The **vector field**  $V$  is **admissible**.



# Geometrical example of non-regular face:



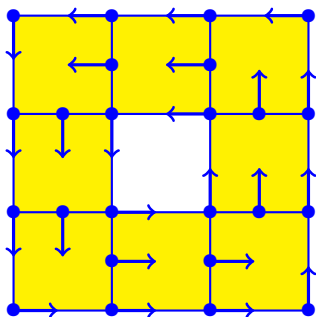
$$C_*(\text{Triangle}) = \{0 \longleftarrow \mathbb{Z}^3 \longleftarrow \mathbb{Z}^3 \longleftarrow \mathbb{Z} \longleftarrow 0\}$$

$$\begin{bmatrix} -1 & -1 & 0 \\ 1 & 0 & -1 \\ 0 & 1 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ -1 \\ 1 \end{bmatrix}$$

$$C_*(P^2\mathbb{R}) = \{0 \longleftarrow \mathbb{Z} \longleftarrow \mathbb{Z} \longleftarrow \mathbb{Z} \longleftarrow 0\}$$

$$\begin{bmatrix} 0 \\ 2 \end{bmatrix}$$

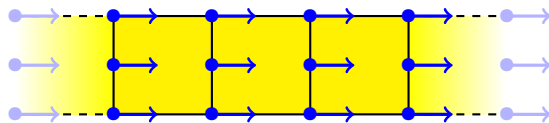
Typical examples of **non-admissible** vector fields.



???!!!

~

$\emptyset$



???!!!

~

$\emptyset$

$\mathbb{R} \times I$

### 3. Homological Reductions.

Definition: A (homological) reduction is a diagram:

$$\rho: \boxed{h \circlearrowleft \hat{C}_* \begin{matrix} \xleftarrow{g} \\ \xrightarrow{f} \end{matrix} C_*}$$

with:

1.  $\hat{C}_*$  and  $C_* =$  chain complexes.
2.  $f$  and  $g =$  chain complex morphisms.
3.  $h =$  homotopy operator (degree +1).
4.  $fg = \text{id}_{C_*}$  and  $d_{\hat{C}}h + hd_{\hat{C}} + gf = \text{id}_{\hat{C}_*}$ .
5.  $fh = 0$ ,  $hg = 0$  and  $hh = 0$ .

$$\begin{array}{c}
 \{ \cdots \xleftarrow[h]{d} \widehat{C}_{m-1} \xleftarrow[h]{d} \widehat{C}_m \xleftarrow[h]{d} \widehat{C}_{m+1} \xleftarrow[h]{d} \cdots \} = \widehat{C}_* \\
 \parallel \\
 \{ \cdots \xleftarrow[h]{d} A_{m-1} \xleftarrow[h]{d} A_m \xleftarrow[h]{d} A_{m+1} \xleftarrow[h]{d} \cdots \} = A_* \\
 \parallel \\
 \{ \cdots \xleftarrow[h]{d} B_{m-1} \xleftarrow[h]{d} B_m \xleftarrow[h]{d} B_{m+1} \xleftarrow[h]{d} \cdots \} = B_* \\
 \parallel \\
 \{ \cdots \xleftarrow[d]{d} C'_{m-1} \xleftarrow[d]{d} C'_m \xleftarrow[d]{d} C'_{m+1} \xleftarrow[d]{d} \cdots \} = C'_* \\
 \parallel \\
 \{ \cdots \xleftarrow[d]{d} C_{m-1} \xleftarrow[d]{d} C_m \xleftarrow[d]{d} C_{m+1} \xleftarrow[d]{d} \cdots \} = C_*
 \end{array}$$

$$A_* = \ker f \cap \ker h$$

$$B_* = \ker f \cap \ker d$$

$$C'_* = \text{im}(g)$$

$$\widehat{C}_* = A_* \oplus B_{*\text{exact}} \oplus C'_* \cong C_*$$

## Fundamental Theorem:

Given:  $C_* = (C_p, \beta_p, d_p)_{p \in \mathbb{Z}} =$  Cellular chain complex.

$V = (\sigma_i, \tau_i)_{i \in v} =$  Admissible Discrete Vector Field.

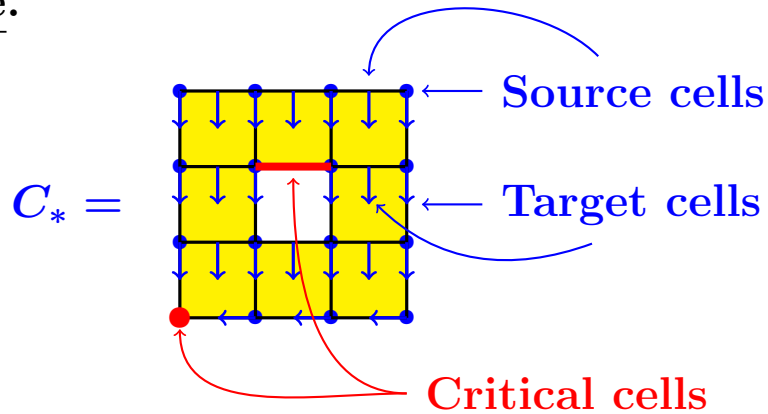
$\Rightarrow$  Canonical Reduction:

$$\rho_V = \boxed{h \circlearrowleft (C_p, \beta_p, d_p)_{p \in \mathbb{Z}} \begin{matrix} \xrightarrow{g} \\ \xleftarrow{f} \end{matrix} (C_p^c, \beta_p^c, d_p^c)_{p \in \mathbb{Z}}}$$

$$\boxed{\text{Initial Complex}} \xRightarrow{\rho_V} \boxed{\text{Critical complex}}$$

Proof: Homological Perturbation Theorem.

## Toy Example.



## Fundamental Reduction Theorem $\Rightarrow$

$$\rho : C_* \rightsquigarrow C_*^c = \begin{array}{c} \boxed{\begin{array}{c} \text{---} \\ \curvearrowright \\ \bullet \end{array}} \end{array} = \mathbb{Z} \xleftarrow{d_1^c=0} \mathbb{Z} = \text{Circle}$$

$\text{Rank}(C_*) = (16, 24, 8)$	vs	$\text{Rank}(C_*^c) = (1, 1, 0)$
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#### 4. Product problem in Combinatorial Topology.

1. **Simplicial** organisation **necessary**

for example for **Eilenberg-MacLane spaces**.

2.  $\Rightarrow$  **Elementary models** =  $\Delta^n$  for  $n \in \mathbb{N}$ .

3. **Fact:**

**No direct simplicial structure** for a product  $\Delta^p \times \Delta^q$ .

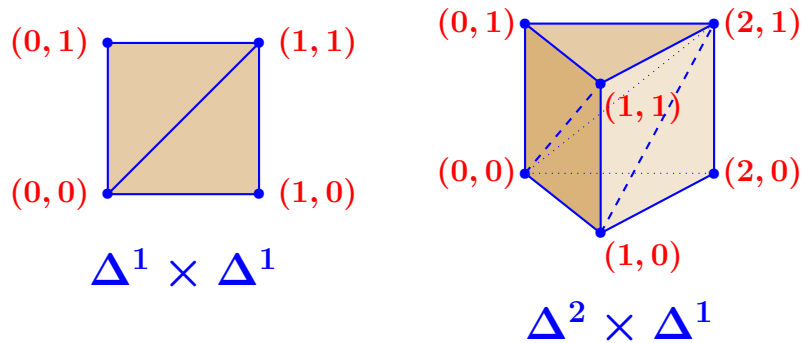
4. What about **twisted products = Fibrations** ??

5. Classical solution = **Eilenberg-Zilber + Kan + RM**

+ **Serre and Eilenberg-Moore Spectral sequences**.

6. Other **solution** = **Discrete Vector Fields**.

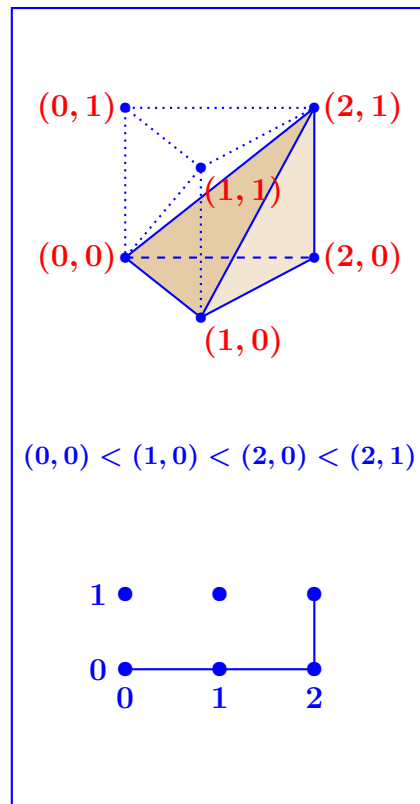
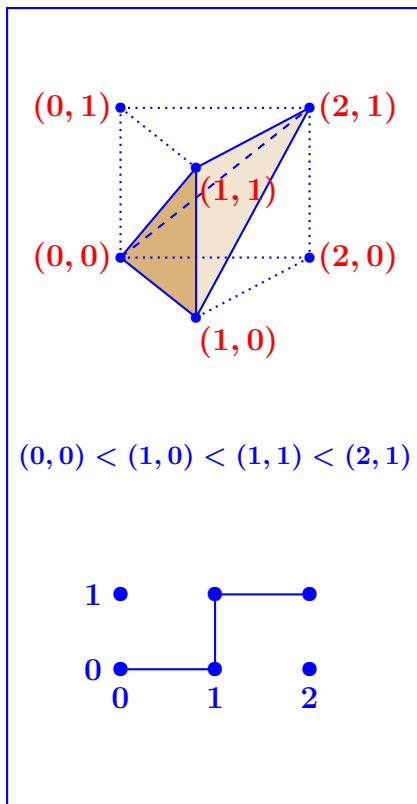
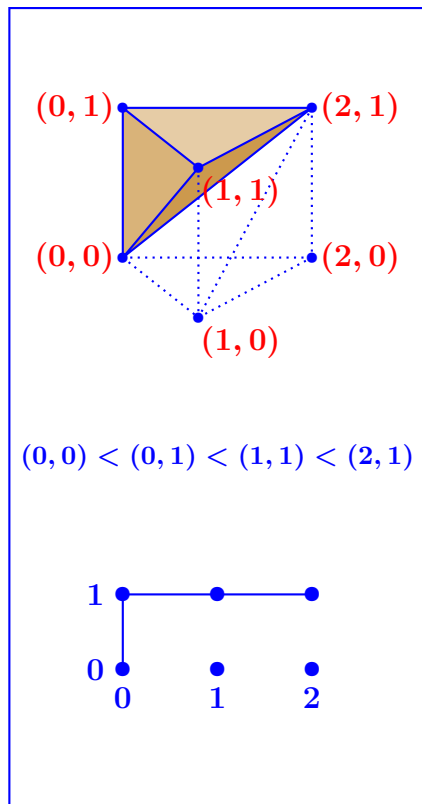




Two  $\Delta^2$  in  $\Delta^1 \times \Delta^1$ :  $(0,0) < (0,1) < (1,1)$   
 $(0,0) < (1,0) < (1,1)$

Three  $\Delta^3$  in  $\Delta^2 \times \Delta^1$ :  $(0,0) < (0,1) < (1,1) < (2,1)$   
 $(0,0) < (1,0) < (1,1) < (2,1)$   
 $(0,0) < (1,0) < (2,0) < (2,1)$

# Rewriting the triangulation of $\Delta^2 \times \Delta^1$ .

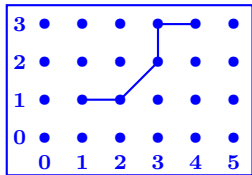


Increasing chain in the lattice  $\longleftrightarrow$  Simplex in the Product

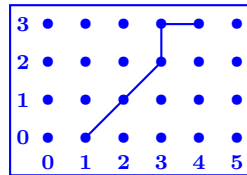
“Seeing” the **triangulation** of  $\Delta^5 \times \Delta^3$ .

Example of 5-simplex : =  $\sigma \in (\Delta^5 \times \Delta^3)_5$

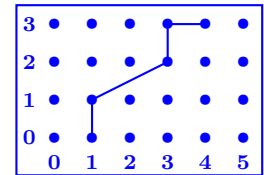
$\Rightarrow$  6 faces:



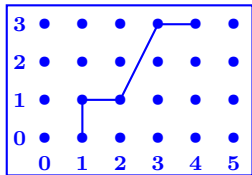
$\partial_0 \sigma$



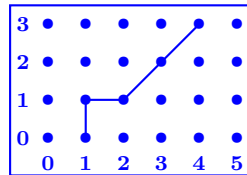
$\partial_1 \sigma$



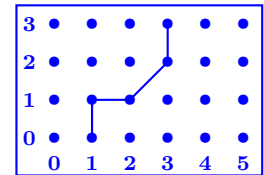
$\partial_2 \sigma$



$\partial_3 \sigma$

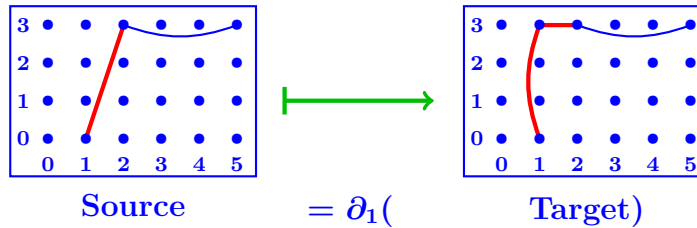
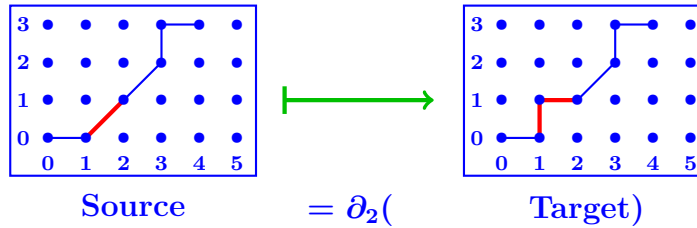




$\partial_4 \sigma$



$\partial_5 \sigma$

⇒ **Canonical discrete vector field** for  $\Delta^5 \times \Delta^3$ .



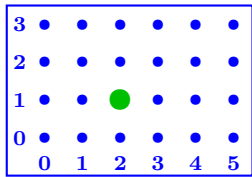
Recipe: First “event” = **Diagonal step** =  ⇒ **Source cell**.  
 = **(-90°)-bend** =  ⇒ **Target cell**.

# Critical cells ??

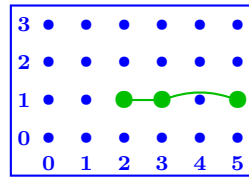
Critical cell = cell without any “event”

= without any diagonal or  $-90^\circ$ -bend.

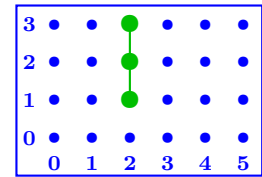
Examples.



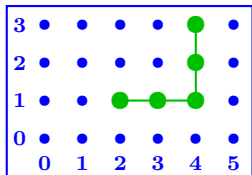
$$\Delta_2^0 \otimes \Delta_1^0$$



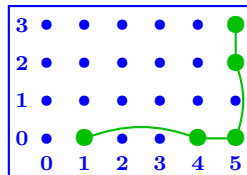
$$\Delta_{2,3,5}^2 \otimes \Delta_1^0$$



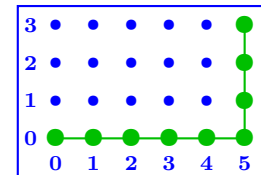
$$\Delta_2^0 \otimes \Delta_{1,2,3}^2$$



$$\Delta_{2,3,4}^2 \otimes \Delta_{1,2,3}^2$$



$$\Delta_{1,4,5}^2 \otimes \Delta_{0,2,3}^2$$



$$\Delta_{0,1,2,3,4,5}^5 \otimes \Delta_{0,1,2,3}^3$$

Conclusion:

$$C_*^c = C_*(\Delta^5) \otimes C_*(\Delta^3)$$

Fundamental theorem of vector fields  $\Rightarrow$

Canonical Homological Reductions:

$$\rho : C_*(\Delta^5 \times \Delta^3) \rightrightarrows C_*(\Delta^5) \otimes C_*(\Delta^3)$$

$$\rho : C_*(\Delta^p \times \Delta^q) \rightrightarrows C_*(\Delta^p) \otimes C_*(\Delta^q)$$

$$p = q = 10 \Rightarrow 16,583,583,743 \text{ vs } 4,190,209$$

More generally:  $X$  and  $Y =$  simplicial sets.

An admissible discrete vector field

is canonically defined on  $C_*(X \times Y)$ .

$\Rightarrow$  Critical chain complex  $C_*^c(X \times Y) = C_*(X) \otimes C_*(Y)$ .

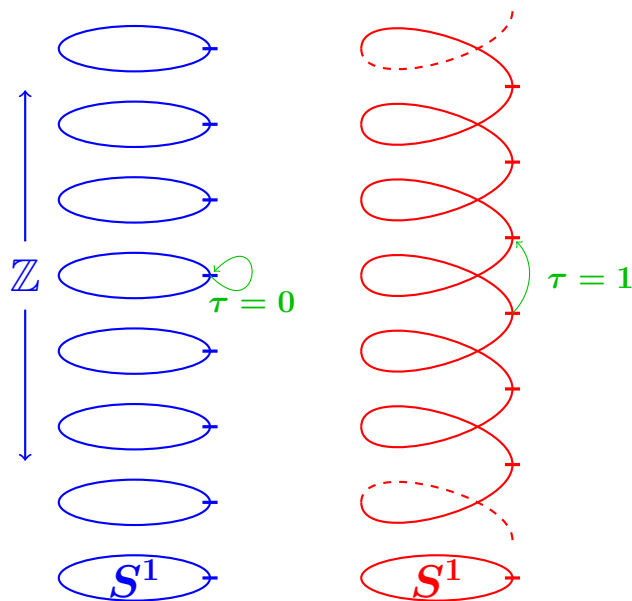
Eilenberg-Zilber Theorem: Canon. homological reduction:

$$\rho_{EZ} : C_*(X \times Y) \xrightarrow{\cong} C_*^c(X \times Y) = C_*(X) \otimes C_*(Y)$$

$\Rightarrow$  Künneth theorem to compute  $H_*(X \times Y)$ .

## 5. Extension to twisted product.

Simplest example:  $\mathbb{Z} \times S^1$  vs  $\mathbb{Z} \times_{\tau} S^1 = \mathbb{R}$ :





General notion of **twisted product**:  $B$  = base space.

$F$  = fibre space.

$G$  = structural group.

Action  $G \times F \rightarrow F$ .

$\tau : B \rightarrow G$  = Twisting function.

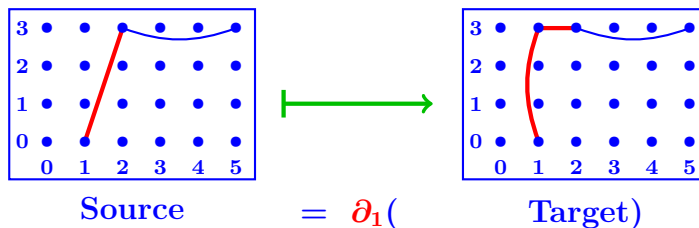
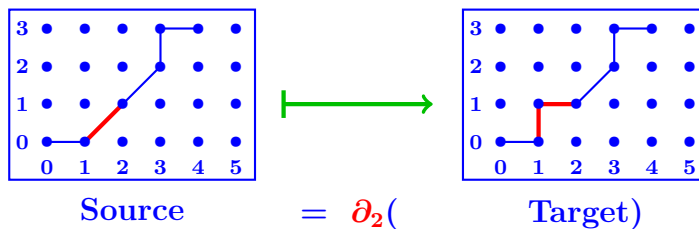
Structure of  $F \times_{\tau} B$ :

$$\partial_i(\sigma_f, \sigma_b) = (\partial_i \sigma_f, \partial_i \sigma_b) \quad \text{for } i > 0$$

$$\partial_0(\sigma_f, \sigma_b) = (\tau(\sigma_b) \cdot \partial_0 \sigma_f, \partial_0 \sigma_b)$$

$\Rightarrow$  Only the 0-face is modified in the **twisted product**.

Reminder about the **EZ-vector field** of  $\Delta^5 \times \Delta^3$ .



The **vector field** is concerned by faces  $\partial_i$  only if  $i > 0$ .

1. The **twisting function**  $\tau$  modifies only  $\boxed{0}$ -faces.
2. The **EZ-vector field**  $V_{EZ}$  of  $X \times Y$   
uses only  $\boxed{i}$ -faces with  $i \geq 1$ .

$\Rightarrow V_{EZ}$  is **defined** and **admissible** as well on  $X \times_{\boxed{\tau}} Y$ .

Fundamental theorem of admissible vector fields  $\Rightarrow$

$$\begin{array}{ccc}
 C_*(X \times Y) & & C_*(X \times_{\boxed{\tau}} Y) \\
 V_{EZ} \Rightarrow \Downarrow & & V_{EZ} \Rightarrow \Downarrow \\
 C_*(X) \otimes C_*(Y) & & C_*(X) \otimes_{\boxed{t}} C_*(Y)
 \end{array}$$

Known as the **twisted Eilenberg-Zilber Theorem**.

Corollary: Base  $B$  1-reduced  $\Rightarrow$  Algorithm:

$$[(F, C_*(F), EC_*^F, \varepsilon_F) + (B, C_*(B), EC_*^B, \varepsilon_B) + G + \tau] \\ \longmapsto (F \times_\tau B, C_*(F \times_\tau B), EC_*^{F \times_\tau B}, \varepsilon_{F \times_\tau B}).$$

Version of  $F$  with effective homology

+ Version of  $B$  with effective homology

+  $G + \tau$  describing the fibration  $F \hookrightarrow F \times_\tau B \rightarrow B$

$\Rightarrow$  Version with effective homology of the total space  $F \times_\tau B$ .

= Version with effective homology

of the Serre Spectral Sequence

## 6. The Eilenberg-Moore spectral sequence.

Key results:

$G = \text{Simplicial group} \Rightarrow BG = \text{classifying space.}$

$$BG = \dots ((SG \times_{\tau} SG) \times_{\tau} SG) \times_{\tau} SG \times_{\tau} \dots$$

$X = \text{Simplicial set} \Rightarrow KX = \text{Kan loop space.}$

$$KX = \dots ((S^{-1}X \times_{\tau} S^{-1}X) \times_{\tau} S^{-1}X) \times_{\tau} S^{-1}X \times_{\tau} \dots$$

Analogous process  $\Rightarrow$  Algorithms:

$$\begin{aligned} (G, C_*G, EC_*^G, \varepsilon_G) &\mapsto (BG, C_*BG, EC_*^{BG}, \varepsilon_{BG}) \\ (X, C_*X, EC_*^X, \varepsilon_X) &\mapsto (KX, C_*KX, EC_*^{KX}, \varepsilon_{KX}) \end{aligned}$$

More generally:

$$[\alpha : E \rightarrow B] + [\alpha' : E' \rightarrow B] + [\alpha \text{ fibration}]$$

$$\Rightarrow \text{algorithm: } (B_{EH}, E_{EH}, E'_{EH}, \alpha, \alpha') \mapsto (E \times_B E')_{EH}.$$

$$\begin{array}{ccc} E' \times_B E & \longrightarrow & E \\ \downarrow & & \downarrow \alpha \\ E' & \xrightarrow{\alpha'} & B \end{array}$$

= Version with effective homology

of Eilenberg-Moore spectral sequence I.

Also:

[ $G$  simplicial group] + [ $\alpha : G \times E \rightarrow E$ ] +  
 [ $\alpha' : E' \times G \rightarrow E'$ ] + [ $\alpha$  principal fibration]  
 $\Rightarrow$  **algorithm:**  $(G_{EH}, E_{EH}, E'_{EH}, \alpha, \alpha') \mapsto (E' \times_G E)_{EH}$ .

$$\begin{array}{ccc}
 G & \xrightarrow{\alpha} & E \\
 \alpha' \swarrow & & \searrow \\
 E' & \longrightarrow & E' \times_G E
 \end{array}
 \quad (\alpha'(x', g), x) \sim (x', \alpha(g, x))$$

= Version **with effective homology**

of **Eilenberg-Moore spectral sequence II.**

Integrating the **Vector Field technology**

in the **Kenzo program**

⇒ **Faster program!**

Example:  $\pi_5(\Omega(S^3) \cup_2 D^3) = ??$



Integrating the **Vector Field technology**

in the **Kenzo program**

⇒ **Faster program!**

Example:  $\pi_5(\Omega(S^3) \cup_2 D^3) = (\mathbb{Z}/2)^4$

On the same machine:

**Old version** ⇒ **1h32m**

**New version** ⇒ **0h05m**

with the **same result !**

**Computing time divided by 18.**

Most recent **calculation**:

$$\pi_6(\Omega(S^3) \cup_2 D_3) = ??$$

14 days long calculation on an idle machine

of our cryptography team (1To RAM):

$$\pi_6(\Omega(S^3) \cup_2 D_3) = (\mathbb{Z}/2)^5 + \mathbb{Z}$$

(14 days  $\times$  18 = 252 days)

The END

```
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---done---  
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Joint work with Ana Romero, Universidad de La Rioja  
Jagiellonian University, Krakow, 12 May 2016*