

Effectiveness vs Gödel

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;; Clock  
Computing  
<TnPr <TnPr  
End of computing.  
  
;; Clock -> 2002-01-17, 19h 25m 36s.  
Computing the boundary of the generator 19 (dimension 7) :  
<TnPr <TnPr <TnPr S3 <<Abar[2 S1][2 S1]>>> <<Abar>>> <<Abar>>>  
End of computing.  
  
Homology in dimension 6 :  
  
Component Z/12Z  
  
---done---  
  
;; Clock -> 2002-01-17, 19h 27m 15s
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SupMéca, Saint-Ouen, May-2013*

Semantics of colours:

Blue = “Standard” Mathematics

Red = Constructive, effective,

algorithm, machine object, ...

Violet = Problem, difficulty,

obstacle, disadvantage, ...

Green = Solution, essential point,

mathematicians, ...

First significant computability result in Algebraic Topology:

ANNALS OF MATHEMATICS
 Vol. 65, No. 1, January, 1957
 Printed in U.S.A.

FINITE COMPUTABILITY OF POSTNIKOV COMPLEXES¹

BY EDGAR H. BROWN, JR.

(Received March 3, 1956)

In [4] Postnikov associates with each arcwise connected space X a sequence of algebraically defined complexes $P_1(X)$, $P_2(X)$, \dots , and in [5] he asserts

 have a finite number of non-degenerate simplexes.)

It must be emphasized that although the procedures developed for solving these problems are finite, they are much too complicated to be considered practical.

In the first section of this paper we give some preliminary definitions, con-

Problem: Tractable algorithms ???

Main obstacle: Infinite objects !!!

Didactic example: $\pi_3(S^2) := \pi_0(\mathcal{C}(S^3, S^2)) = ???$

Hurewicz Theorem $\Rightarrow \pi_2(S^2) = H_2(S^2) = \mathbb{Z}$. (easy)

$X_3 = K(\mathbb{Z}, 1) \times_{\tau} S^2$ with appropriate $\tau \Rightarrow$

$$\pi_3(S^2) \stackrel{\tau}{=} \pi_3(X_3) \stackrel{\text{Hurewicz}}{=} H_3(X_3)$$

But $K(\mathbb{Z}, 1) =$ infinite simplicial set

\Rightarrow Implementation of $K(\mathbb{Z}, 1)$ on a machine ???

Functional trick:

The **functional trick** often allows one
to **implement** infinite objects.

Example: $K(\mathbb{Z}, 1)$:

$\{n\text{-simplices}\} := \{\text{lists of } n \text{ integers}\}$

$\sigma = (12 \ 23 \ 34 \ 45 \ 56) = 5\text{-simplex of } K(\mathbb{Z}, 1).$

Example of geometric structure: 3-face of $\sigma = ???$

$\partial_3 \sigma := (12 \ 23 \ 79 \ 56) \quad (34+45=79)$

$$H_*(K(\mathbb{Z}, 1)) = ???$$

Theorem (Gödel, Church, Turing, Post):

∄ general algorithm:

$$[\text{Simplicial set functionally coded } X] \longmapsto H_*(X)$$

E. Brown solution: Approximation of infinite objects
by sufficiently large finite approximations.

Effective Homology solution: Defining pairs:

Functional Infinite objects $\xleftrightarrow{\varepsilon}$ Finite objects

connected by functional objects.

Definition: An object with effective homology X is a 4-tuple:

$$X = \boxed{X, C_*(X), EC_*, \varepsilon}$$

with:

1. X = an arbitrary object (simplicial set, simplicial group, differential graded algebra, ...)
2. $C_*(X)$ = “the” chain complex “traditionally” associated with X to define the homology groups $H_*(X)$.
3. EC_* = some effective chain complex.
4. ε = some equivalence $C_*(X) \overset{\varepsilon}{\rightleftarrows} EC_*$.

Main result of **effective homology**:

Meta-theorem: Let X_1, \dots, X_n be a collection of **objects** with **effective homology** and ϕ be a **reasonable construction process**:

$$\phi : (X_1, \dots, X_n) \mapsto X.$$

Then **there exists a version with effective homology** ϕ_{EH} :

$$\phi_{EH}: \left(\boxed{X_1, C_*(X_1), EC_{1*}, \varepsilon_1}, \dots, \boxed{X_n, C_*(X_n), EC_{n*}, \varepsilon_n} \right) \mapsto \boxed{X, C_*(X), EC_*, \varepsilon}$$

The process is **perfectly stable**

and can be **again used** with X for **further calculations**.

Example:

Julio Rubio's solution of Adams' problem:

$\Omega X := \mathcal{C}(S^1, X)$. Adams: Computation of $H_*(\Omega^n X)$???

$$X = (X, C_*(X), EC_*^X, \varepsilon^X)$$

$$\Downarrow\Downarrow\Downarrow\Downarrow \quad \Omega_{EH}$$

$$\Omega X = (\Omega X, C_*(\Omega X), EC_*^{\Omega X}, \varepsilon^{\Omega X})$$

\implies Trivial iteration now available.

⇒ Very simple solution of Adam's problem :

Indefinite iteration of the Cobar construction ???

$$X = (X, C_*(X), EC_*^X, \varepsilon^X)$$

$$\Downarrow \Omega_{EH}$$

$$\Omega X = (\Omega X, C_*(\Omega X), EC_*^{\Omega X}, \varepsilon^{\Omega X})$$

$$\Downarrow \Omega_{EH}$$

$$\Omega^2 X = (\Omega^2 X, C_*(\Omega^2 X), EC_*^{\Omega^2 X}, \varepsilon^{\Omega^2 X})$$

$$\Downarrow \Omega_{EH}$$

$$\Omega^3 X = (\Omega^3 X, C_*(\Omega^3 X), EC_*^{\Omega^3 X}, \varepsilon^{\Omega^3 X})$$

$$\Downarrow \Omega_{EH}$$

$$\Omega^4 X = \dots$$



Simple example of **calculation** now **possible**.

$$X := P^\infty \mathbb{R} / P^3 \mathbb{R}$$

X is not a suspension \Rightarrow **computation** of $H_* \Omega^3 X = ???$

X has **finite type** \Rightarrow

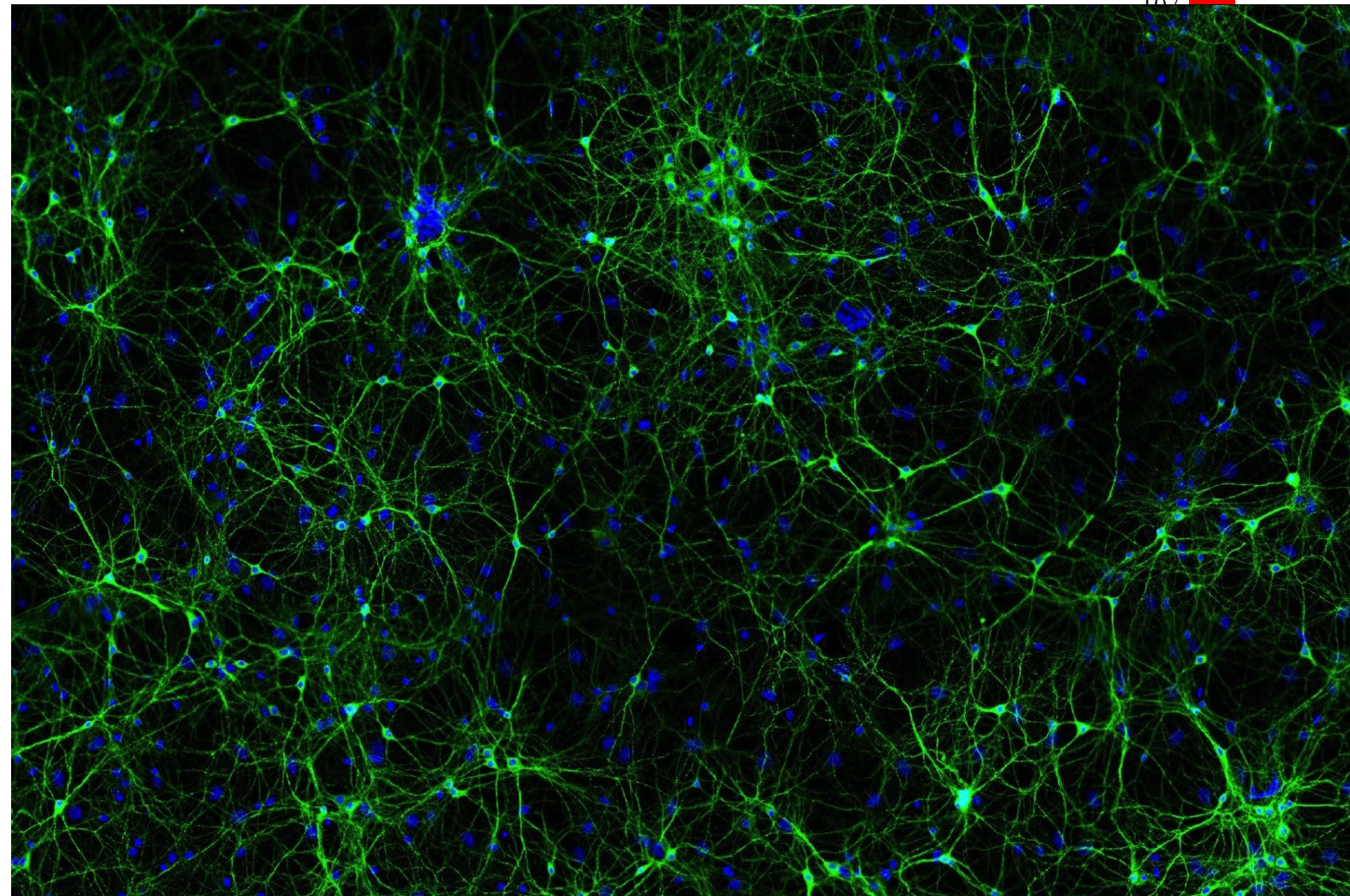
trivial object with eff. hom.: $(X, C_* X, C_* X, \varepsilon_0 = \text{id})$.

Julio Rubio $\Rightarrow (\Omega^3 X, C_* \Omega^3 X, EC_* \Omega^3 X, \varepsilon_3)$

$EC_* \Omega^3 X$ has **finite type** \Rightarrow

Computing $H_*(EC_* \Omega^3 X)$ **is elementary**.

Example: $H_5 \Omega^3 X = H_5(EC_* \Omega^3 X) = (\mathbb{Z}/3) \oplus (\mathbb{Z}/2)^5 \oplus \mathbb{Z}$



The END

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