

# Homological Perturbation Theorem and Eilenberg-Zilber vector field

```
;; Clock  
Computing  
<TnPr <Tn  
End of computing.  
  
;; Clock -> 2002-01-17, 19h 25m 36s.  
Computing the boundary of the generator 19 (dimension 7) :  
<TnPr <TnPr <TnPr S3 <<Abar[2 $1][2 $1]>>> <<Abar>>>  
End of computing.  
  
Homology in dimension 6 :  
  
Component 2/122  
  
---done---  
;; Clock -> 2002-01-17, 19h 27m 15s
```

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ETH Zurich, June 2012

$R =$  Unitary ring

$\varepsilon, \varphi, \psi, \beta \in R$  with  $\varepsilon$  invertible.

Gauss discussion of (1) + (2):

$$(1) \quad \varepsilon x + \varphi y = a$$

$$(2) \quad \psi x + \beta y = b$$

$$(2) - \psi \varepsilon^{-1} (1) \Rightarrow$$

$$(2') \quad (\beta - \psi \varepsilon^{-1} \varphi) y = (b - \psi \varepsilon^{-1} a)$$

$\Rightarrow$  (1) + (2) has a solution  $\Leftrightarrow$

$$(\beta - \psi \varepsilon^{-1} \varphi) \mid (b - \psi \varepsilon^{-1} a) \Rightarrow y = \dots$$

$$\Rightarrow x = \varepsilon^{-1} a - \varepsilon^{-1} \varphi y$$

Matrix translation:

$$\begin{pmatrix} \varepsilon & \varphi \\ \psi & \beta \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} a \\ b \end{pmatrix}$$

$\Leftrightarrow$

$$\underbrace{\begin{pmatrix} 1 & 0 \\ -\psi\varepsilon^{-1} & 1 \end{pmatrix} \begin{pmatrix} \varepsilon & \varphi \\ \psi & \beta \end{pmatrix} \begin{pmatrix} 1 & -\varepsilon^{-1}\varphi \\ 0 & 1 \end{pmatrix}}_{\begin{pmatrix} \varepsilon & 0 \\ 0 & \beta - \psi\varepsilon^{-1}\varphi \end{pmatrix}} \underbrace{\begin{pmatrix} 1 & \varepsilon^{-1}\varphi \\ 0 & 1 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix}}_{\begin{pmatrix} x + \varepsilon^{-1}\varphi y \\ y \end{pmatrix}} = \underbrace{\begin{pmatrix} 1 & 0 \\ -\psi\varepsilon^{-1} & 1 \end{pmatrix} \begin{pmatrix} a \\ b \end{pmatrix}}_{\begin{pmatrix} a \\ -\psi\varepsilon^{-1}a + b \end{pmatrix}}$$

$\Leftrightarrow$

$$\varepsilon(x + \varepsilon^{-1}\varphi y) = a$$

$$(\beta - \psi\varepsilon^{-1}\varphi) y = (b - \psi\varepsilon^{-1}a)$$

$\Leftrightarrow$

$$(\beta - \psi\varepsilon^{-1}\varphi) \mid (b - \psi\varepsilon^{-1}a) \Rightarrow \dots$$

Diagram translation:

$$\begin{array}{ccc}
 & \left( \begin{matrix} 1 & -\varepsilon^{-1}\varphi \\ 0 & 1 \end{matrix} \right) & \\
 R^2 & \xleftarrow{\hspace{1cm}} & R^2 \\
 \left( \begin{matrix} \varepsilon & \varphi \\ \psi & \beta \end{matrix} \right) \downarrow & \left( \begin{matrix} 1 & \varepsilon^{-1}\varphi \\ 0 & 1 \end{matrix} \right) & \downarrow \left( \begin{matrix} \varepsilon & 0 \\ 0 & \beta - \psi\varepsilon^{-1}\varphi \end{matrix} \right) \\
 & \left( \begin{matrix} 1 & -0 \\ \psi\varepsilon^{-1} & 1 \end{matrix} \right) & \\
 R^2 & \xleftarrow{\hspace{1cm}} & R^2 \\
 \left( \begin{matrix} 1 & -0 \\ -\psi\varepsilon^{-1} & 1 \end{matrix} \right) & &
 \end{array}$$

Combined with an obvious reduction:

$$\begin{array}{ccc}
 & \left( \begin{array}{cc} 1 & -\varepsilon^{-1}\varphi \\ 0 & 1 \end{array} \right) & \left( \begin{array}{c} 0 \\ 1 \end{array} \right) \\
 R^2 & \xleftarrow{\quad} & R^2 \xleftarrow{\quad} R \\
 & \left( \begin{array}{cc} 1 & \varepsilon^{-1}\varphi \\ 0 & 1 \end{array} \right) & \left( \begin{array}{cc} 0 & 1 \end{array} \right) \\
 \left( \begin{array}{cc} \varepsilon & \varphi \\ \psi & \beta \end{array} \right) & \downarrow & \downarrow \left( \begin{array}{cc} \varepsilon^{-1} & 0 \\ 0 & 0 \end{array} \right) \left( \begin{array}{cc} \varepsilon & 0 \\ 0 & \beta - \psi\varepsilon^{-1}\varphi \end{array} \right) \left( \begin{array}{c} 0 \\ 1 \end{array} \right) (\beta - \psi\varepsilon^{-1}\varphi) \\
 & \left( \begin{array}{cc} 1 & 0 \\ \psi\varepsilon^{-1} & 1 \end{array} \right) & \left( \begin{array}{c} 0 \\ 1 \end{array} \right) \\
 R^2 & \xleftarrow{\quad} & R^2 \xleftarrow{\quad} R \\
 & \left( \begin{array}{cc} 1 & 0 \\ -\psi\varepsilon^{-1} & 1 \end{array} \right) & \left( \begin{array}{cc} 0 & 1 \end{array} \right)
 \end{array}$$

$\Rightarrow$

$\Rightarrow$  Canonical reduction induced by  $\varepsilon$  invertible

$$\begin{array}{ccc}
 & g = \begin{pmatrix} -\varepsilon^{-1}\varphi \\ 1 \end{pmatrix} & \\
 R^2 & \xrightleftharpoons[f = \begin{pmatrix} 0 & 1 \end{pmatrix}]{} & R \\
 \uparrow & \left( \begin{array}{cc} \varepsilon^{-1} & 0 \\ 0 & 0 \end{array} \right) & \downarrow (\beta - \psi\varepsilon^{-1}\varphi) \\
 & \left( \begin{array}{cc} \varepsilon & \varphi \\ \psi & \beta \end{array} \right) & \\
 & g = \begin{pmatrix} 0 \\ 1 \end{pmatrix} & \\
 R^2 & \xrightleftharpoons[f = \begin{pmatrix} -\psi\varepsilon^{-1} & 1 \end{pmatrix}]{} & R
 \end{array}$$

The same is valid with

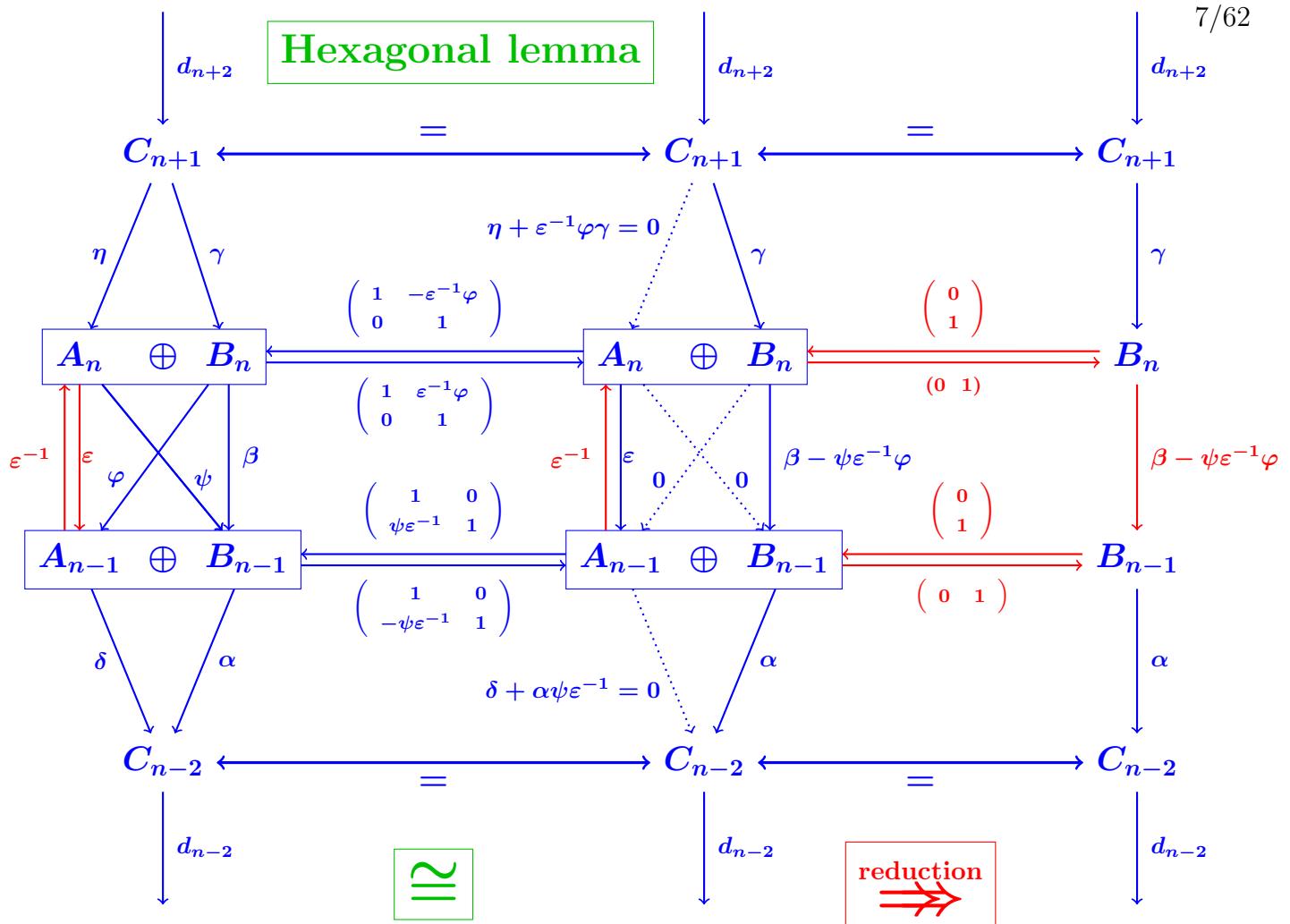
$$\begin{aligned} \mathbf{R}^2 = \mathbf{R} \oplus \mathbf{R} &\text{ replaced by } A_n \oplus B_n = C_n \\ &\text{or by } A_{n-1} \oplus B_{n-1} = C_{n-1} \end{aligned}$$

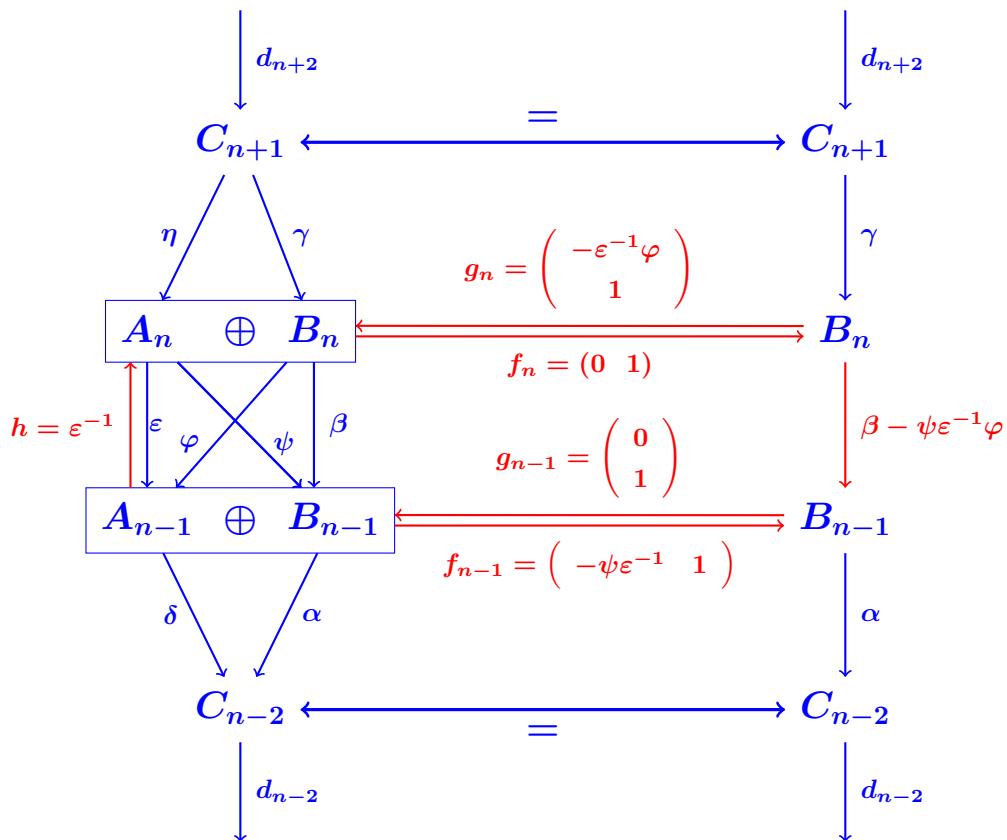
and:

$$\begin{pmatrix} \varepsilon & \varphi \\ \psi & \beta \end{pmatrix} : A_n \oplus B_n \rightarrow A_{n-1} \oplus B_{n-1}$$

with  $\varepsilon : A_n \rightarrow A_{n-1}$  isomorphism.

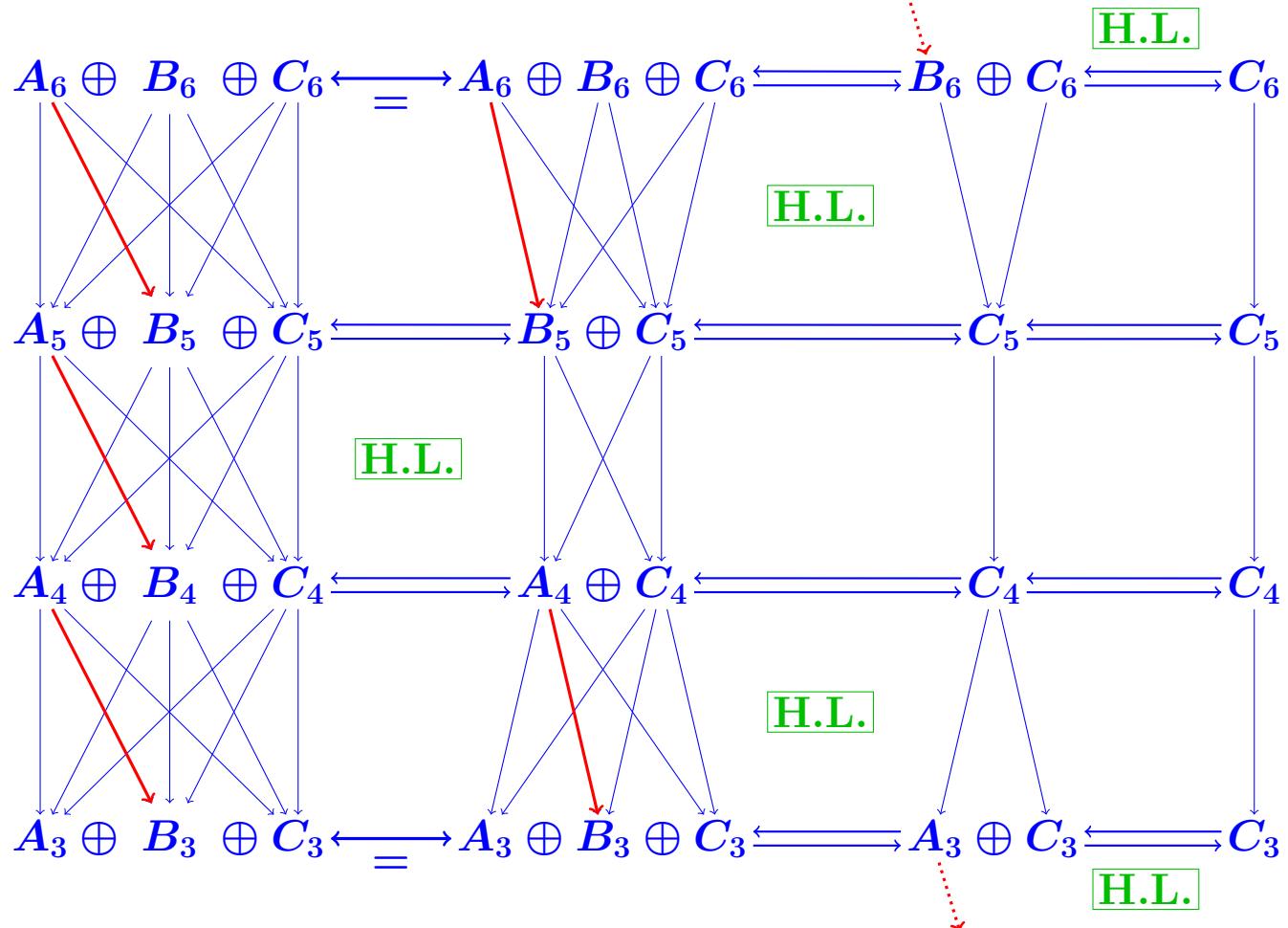
$\Rightarrow$  Hexagonal lemma.





Hexagonal lemma

## Iterating the Hexagonal Lemma:



Final result:

$$\begin{array}{c}
 g_5 = \begin{pmatrix} -d_{21}^{5^{-1}} d_{23}^5 \\ 0 \\ 1 \end{pmatrix} \\
 f_5 = (0 \quad -d_{31}^6 d_{21}^{6^{-1}} \quad 1) \\
 d^5 = d_{33}^5 - d_{31}^5 d_{21}^{5^{-1}} d_{23}^5 \\
 g_4 = \begin{pmatrix} -d_{21}^{4^{-1}} d_{23}^4 \\ 0 \\ 1 \end{pmatrix} \\
 f_4 = (0 \quad -d_{31}^5 d_{21}^{5^{-1}} \quad 1) \\
 d^5 = \begin{pmatrix} d_{11}^5 & d_{12}^5 & d_{13}^5 \\ d_{21}^5 & d_{22}^5 & d_{23}^5 \\ d_{31}^5 & d_{32}^5 & d_{33}^5 \end{pmatrix}
 \end{array}$$

Final result:

$$A_5 \oplus B_5 \oplus C_5 \xrightleftharpoons{f_5 = (0 \quad -d_{31}^6 d_{21}^{6-1} \quad 1)} C_5$$

$$d^5 = d_{33}^5 - d_{31}^5 d_{21}^{5-1} d_{23}^5$$

$$g_5 = \begin{pmatrix} -d_{21}^5 & -1 \\ 0 & 1 \end{pmatrix}$$

$$A_4 \oplus B_4 \oplus C_4 \xrightleftharpoons{f_4 = (0 \quad -d_{31}^5 d_{21}^{5-1} \quad 1)} C_4$$

$$d^5 = \begin{pmatrix} d_{11}^5 & d_{12}^5 & d_{13}^5 \\ d_{21}^5 & d_{22}^5 & d_{23}^5 \\ d_{31}^5 & d_{32}^5 & d_{33}^5 \end{pmatrix}$$

$$g_4 = \begin{pmatrix} -d_{21}^4 & -1 \\ 0 & 1 \end{pmatrix}$$

Final result:

$$\begin{array}{c}
 g_5 = \begin{pmatrix} -d_{21}^5{}^{-1}d_{23}^5 \\ 0 \\ 1 \end{pmatrix} \\
 A_5 \oplus B_5 \oplus C_5 \xleftrightarrow{f_5 = (0 \quad -d_{31}^6 d_{21}^6{}^{-1} \quad 1)} C_5 \\
 \downarrow \\
 d^5 \\
 \text{---} \\
 d_{21}^5 \quad d_{21}^{5-1} \quad d_{23}^5 \\
 A_4 \oplus B_4 \oplus C_4 \xleftrightarrow{f_4 = (0 \quad -d_{31}^5 d_{21}^5{}^{-1} \quad 1)} C_4 \\
 \downarrow \\
 d^5 = d_{33}^5 - d_{31}^5 d_{21}^5{}^{-1} d_{23}^5
 \end{array}$$

$$d^5 = \begin{pmatrix} d_{11}^5 & d_{12}^5 & d_{13}^5 \\ d_{21}^5 & d_{22}^5 & d_{23}^5 \\ d_{31}^5 & d_{32}^5 & d_{33}^5 \end{pmatrix}$$

Note :  $\text{im}(g) \subset A_* \oplus C_*$   
 $= \text{graph } d_{21}^{5-1} d_{23}^5 : C_* \rightarrow A_*$

Final result:

$$\begin{array}{c}
 g_5 = \begin{pmatrix} -d_{21}^{5^{-1}} d_{23}^5 \\ 0 \\ 1 \end{pmatrix} \\
 f_5 = (0 \quad -d_{31}^6 d_{21}^{6^{-1}} \quad 1) \\
 d^5 = d_{33}^5 - d_{31}^5 d_{21}^{5^{-1}} d_{23}^5 \\
 g_4 = \begin{pmatrix} -d_{21}^{4^{-1}} d_{23}^4 \\ 0 \\ 1 \end{pmatrix} \\
 f_4 = (0 \quad -d_{31}^5 d_{21}^{5^{-1}} \quad 1) \\
 d^5 = \begin{pmatrix} d_{11}^5 & d_{12}^5 & d_{13}^5 \\ d_{21}^5 & d_{22}^5 & d_{23}^5 \\ d_{31}^5 & d_{32}^5 & d_{33}^5 \end{pmatrix}
 \end{array}$$

## Global Hexagonal Theorem:

Input: A **chain complex**  $(C_*, d_*)$

with for every  $n \in \mathbb{Z}$  a decomposition:

$$C_n = C_n^1 \oplus C_n^2 \oplus C_n^3 \quad d_n = \begin{pmatrix} d_{n,11} & d_{n,12} & d_{n,13} \\ d_{n,21} & d_{n,22} & d_{n,23} \\ d_{n,31} & d_{n,32} & d_{n,33} \end{pmatrix}$$

with  $d_{n,21} : C_n^1 \rightarrow C_{n-1}^2$  isomorphism  $\forall n$ .

Output: A canonical **reduction**:

$$(C_*, d_*) = (C_*^1 \oplus C_*^2 \oplus C_*^3, d_*) \Rightarrow (C_*^3, d'_*)$$

Application: **Basic Perturbation Lemma (BPL)**

Definition:  $(C_*, d) =$  given chain complex.

A perturbation  $\delta : C_* \rightarrow C_{*-1}$  is an operator of degree -1

satisfying  $(d + \delta)^2 = 0$  ( $\Leftrightarrow (d\delta + \delta d + \delta^2) = 0$ ):  
 $(C_*, d) + (\delta) \mapsto (C_*, d + \delta)$ .

Let  $\rho : h \curvearrowright (\widehat{C}_*, \widehat{d}_*) \xleftarrow[f]{g} (C_*, d_*)$  be a given reduction

and  $\widehat{\delta}$  a perturbation of  $\widehat{d}$

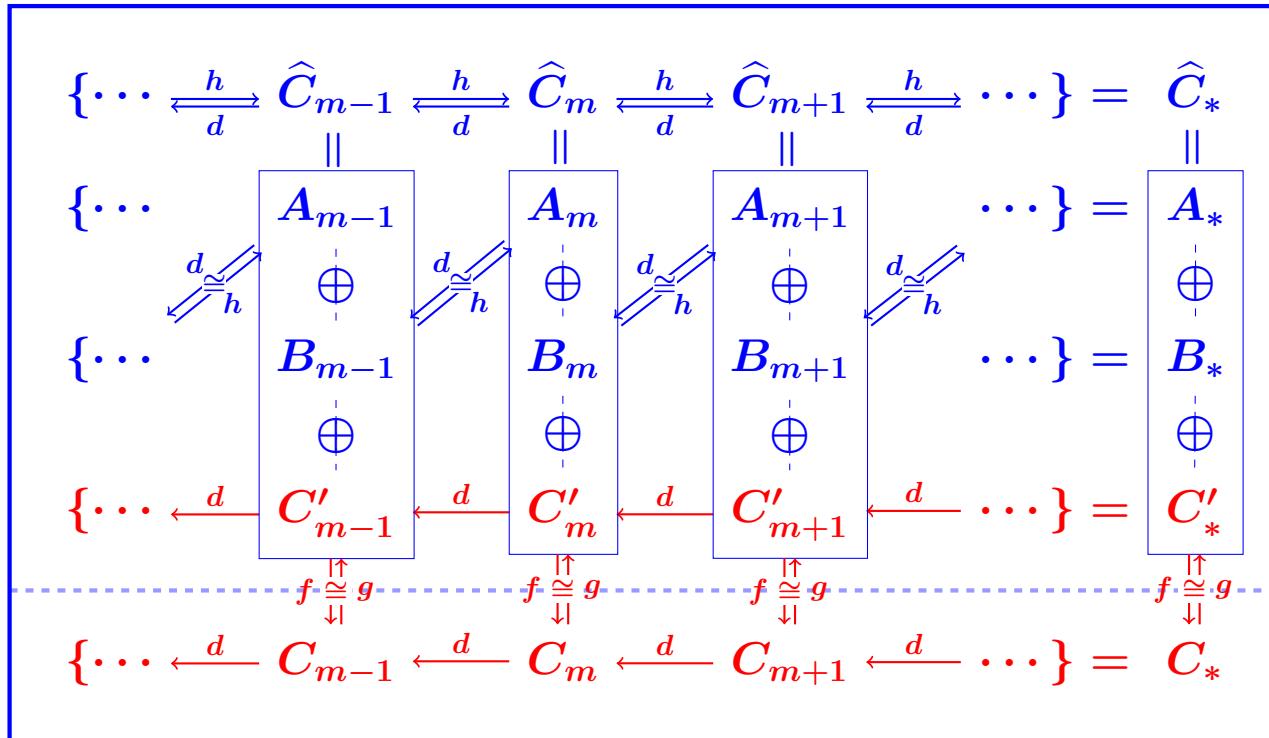
satisfying  $h\widehat{\delta}$  pointwise nilpotent.

Theorem: The BPL determines a new reduction:

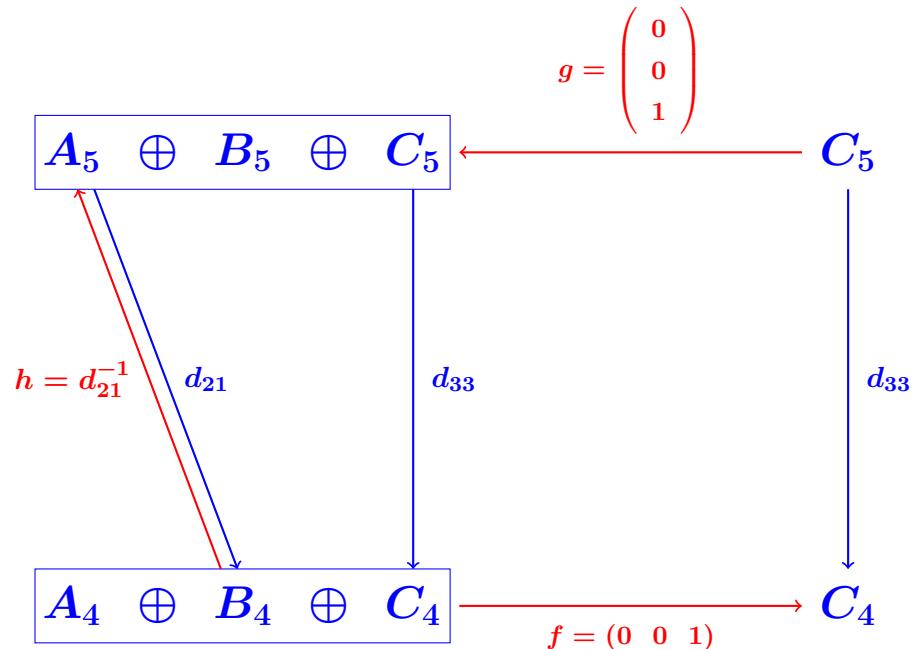
$\rho' : h + \delta_h \curvearrowright (\widehat{C}_*, \widehat{d}_* + \widehat{\delta}_*) \xleftarrow[f + \delta_f]{g + \delta_g} (C_*, d_* + \delta_{d*})$

Proof:

Reduction Diagram:



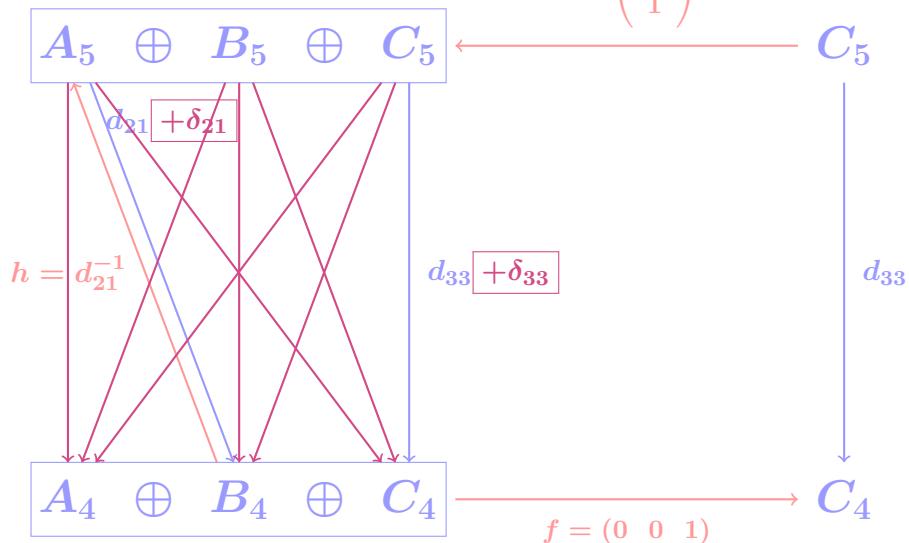
Main part:



with  $d_{21} = \text{isomorphism.}$

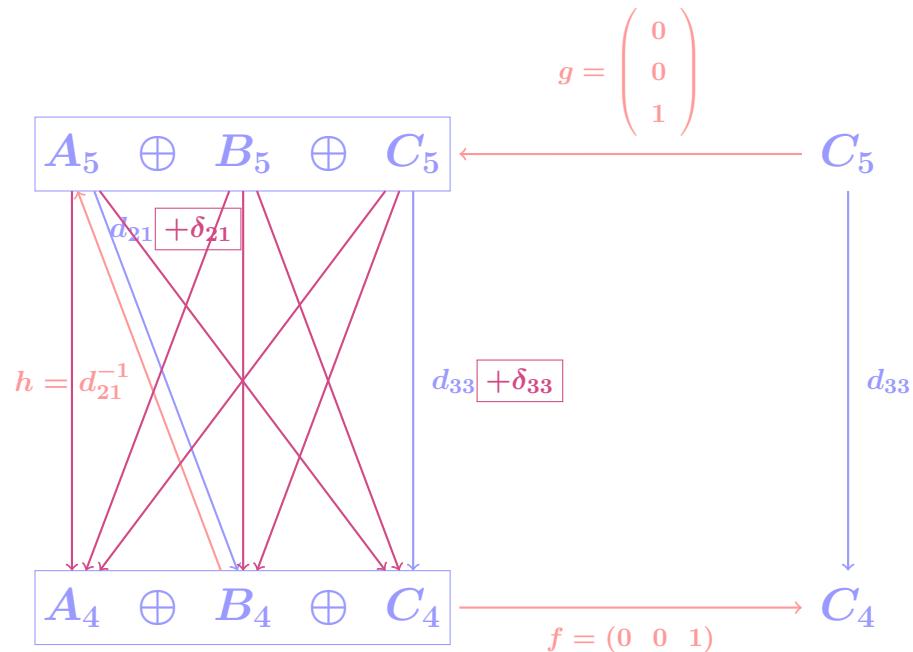
$$\text{Perturbation} = \begin{pmatrix} \delta_{11} & \delta_{12} & \delta_{13} \\ \delta_{21} & \delta_{22} & \delta_{23} \\ \delta_{31} & \delta_{32} & \delta_{33} \end{pmatrix} :$$

$$g = \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix}$$



Question:  $(d_{21} + \delta_{21})$  again isomorphism?

(applying the **Global Hexagonal Theorem** possible?)



But  $d_{21}$  invertible with  $d_{21}h = 1 \Rightarrow$

$$d_{21} + \delta_{21} = d_{21} + d_{21}h\delta_{21} = d_{21}(1 + h\delta_{21})$$

$\Rightarrow d_{21} + \delta_{21}$  invertible  $\Leftrightarrow (1 + h\delta_{21})$  invertible.

A sufficient condition is  $h\delta_{21}$  nilpotent, in which case:

$$(1 + h\delta_{21})^{-1} = \sum_{i=0}^{\infty} (-1)^i (h\delta_{21})^i$$

Then:

$$(d_{21} + \delta_{21})^{-1} =: h' := \left( \sum_{i=0}^{\infty} (-1)^i (h\delta_{21})^i \right) h$$

Remark:

$$\left( \sum_{i=0}^{\infty} (-1)^i (h\delta_{21})^i \right) h = \left( \sum_{i=0}^{\infty} (-1)^i (h\delta)^i \right) h$$

## Global Hexagonal Theorem:

$A_5 \oplus B_5 \oplus C_5 \xleftrightarrow{f_5 = (0 \ -d_{31}^6 d_{21}^{6-1} \ 1)} C_5$

$A_4 \oplus B_4 \oplus C_4 \xleftrightarrow{f_4 = (0 \ -d_{31}^5 d_{21}^{5-1} \ 1)} C_4$

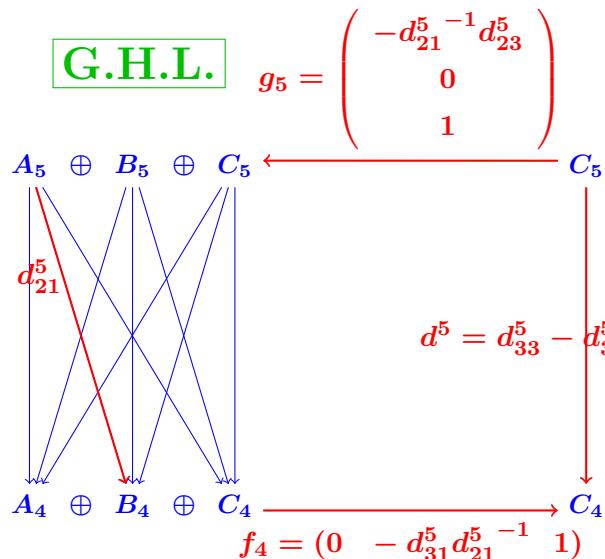
$d^5 = \begin{pmatrix} d_{11}^5 & d_{12}^5 & d_{13}^5 \\ d_{21}^5 & d_{22}^5 & d_{23}^5 \\ d_{31}^5 & d_{32}^5 & d_{33}^5 \end{pmatrix}$

$g_5 = \begin{pmatrix} -d_{21}^{5-1} d_{23}^5 \\ 0 \\ 1 \end{pmatrix}$

$g_4 = \begin{pmatrix} -d_{21}^{4-1} d_{23}^4 \\ 0 \\ 1 \end{pmatrix}$

$d^5 = d_{33}^5 - d_{31}^5 d_{21}^{5-1} d_{23}^5$

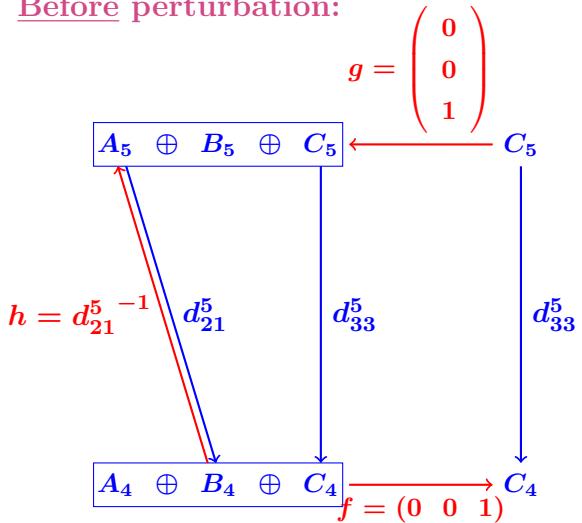
Applying to our situation:



$$d_{21}^5 \mapsto d_{21}^5 + \delta_{21}^5$$

$$d_{21}^{5\ -1} \mapsto h' = \left( \sum_{i=0}^{\infty} (-1)^i (h\delta)^i \right) h$$

Before perturbation:



$$g_5 \mapsto (1 - h'\delta)g$$

$$f_4 \mapsto f(1 - \delta h')$$

$$\begin{aligned} d^5 &\mapsto (d_{33}^5 + \delta_{33}^5) - f\delta h'\delta g \\ &= d_{33}^5 + f\delta g - f\delta h'\delta g \end{aligned}$$

= Basic Perturbation Lemma

QED

## Application to Vector Fields.

$(C_*, d_*, \beta_*)$  = Cellular chain complex.

$V = \{(\sigma_i, \tau_i)\}_{i \in I}$  = Admissible discrete vector field.

⇒ Canonical reduction:

$$(C_*, d_*, \beta_*) \Rightarrow (C^c_*, d^c_*, \beta^c_*)$$

on the critical chain complex  $C^c_*$ :

$$\rho : h \hookrightarrow (C_*, d_*) \xleftarrow[f]{g} (C^c_*, d^c_*)$$

Forman's Morse subcomplex =  $g(C^c_*) \subset C_*$

Definition: Incidence number  $\varepsilon(\sigma, \tau) :=$  coefficient of  $\sigma$  in  $d\tau$ .

Remark:  $V$  vector field  $\Rightarrow \varepsilon(\sigma, \tau)$  invertible if  $(\sigma, \tau) \in V$ .

Definition: The *gradient*  $v$

associated to the vector field  $V = \{(\sigma_i, \tau_i)\}_{i \in I}$

is the *codifferential*  $v_* : C_* \rightarrow C_{*+1}$  defined by:

$$v(\sigma) = \begin{cases} \varepsilon(\sigma, \tau)^{-1}\tau & \text{if } (\sigma, \tau) \in V, \\ 0 & \text{otherwise.} \end{cases}$$

Definition: The *cogradient*  $v^{-1}$

associated to the vector field  $V = \{(\sigma_i, \tau_i)\}_{i \in I}$

is the *differential*  $v_*^{-1} : C_* \rightarrow C_{*-1}$  defined by:

$$v^{-1}(\tau) = \begin{cases} \varepsilon(\sigma, \tau)\sigma & \text{if } (\sigma, \tau) \in V, \\ 0 & \text{otherwise.} \end{cases}$$

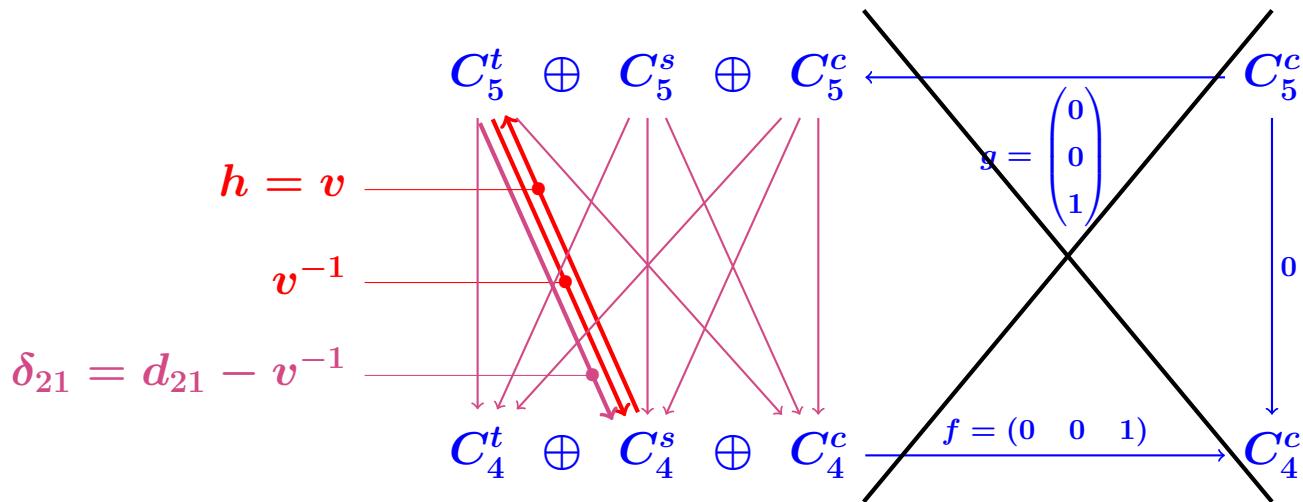
## Applying the BPL.

Initial situation:

$$\begin{array}{c}
 C_5^t \oplus C_5^s \oplus C_5^c \xleftarrow{\quad} C_5^c \\
 \downarrow v^{-1} \quad \downarrow \approx \quad \downarrow v = h \\
 C_4^t \oplus C_4^s \oplus C_4^c \xrightarrow{f = (0 \quad 0 \quad 1)} C_4^c
 \end{array}$$

$g = \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix}$

Restoring the differential  $d$  of  $C_*$ :



$$d_{21} = v^{-1} + (\delta_{21} - v^{-1}) = v^{-1} + \delta_{21}$$

But  $h \ \delta_{21} = v(d_{21} - v^{-1})$  nilpotent ???

Examining:  $\xrightarrow{v} \xrightarrow{d_{21}-v^{-1}} \xrightarrow{v} \xrightarrow{d_{21}-v^{-1}} \xrightarrow{v} \xrightarrow{d_{21}-v^{-1}} \xrightarrow{v} \dots$

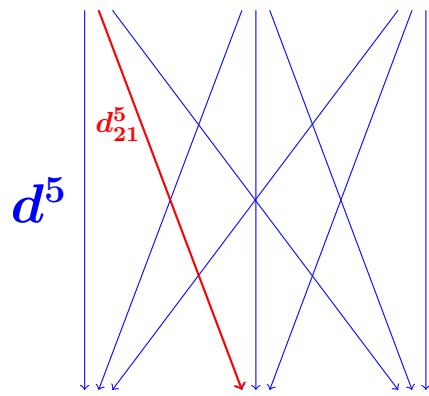
1. Starting **cell** =  $\sigma \in C_4^s$ .
2.  $v(\sigma) = \varepsilon(\sigma, \tau)^{-1}\tau \sim$  target cell associated by  $v$ .
3.  $v^{-1}(\varepsilon(\sigma, \tau)^{-1}\tau) = \sigma$   
 $\Rightarrow (d_{21} - v^{-1})(\tau) =$  Comb. of {sources faces of  $\tau \neq \sigma$ }
- $\Rightarrow$  beginning all the **V-paths** starting from  $\sigma$ .
4.  $v \Rightarrow$  trying to extend **V-paths**.
5.  $(d_{21} - v^{-1})(\tau) \Rightarrow$  trying to extend **V-paths**.
6. . . . . . . . .

$\Rightarrow \delta_{21} h$  nilpotent  $\Leftrightarrow V$  admissible.

Remember :

$$g_5 = \begin{pmatrix} -d_{21}^{5^{-1}} d_{23}^5 \\ 0 \\ 1 \end{pmatrix}$$

$$A_5 \oplus B_5 \oplus C_5 \xrightleftharpoons[f_5 = (0 \quad -d_{31}^6 d_{21}^{6^{-1}} \quad 1)]{} C_5$$



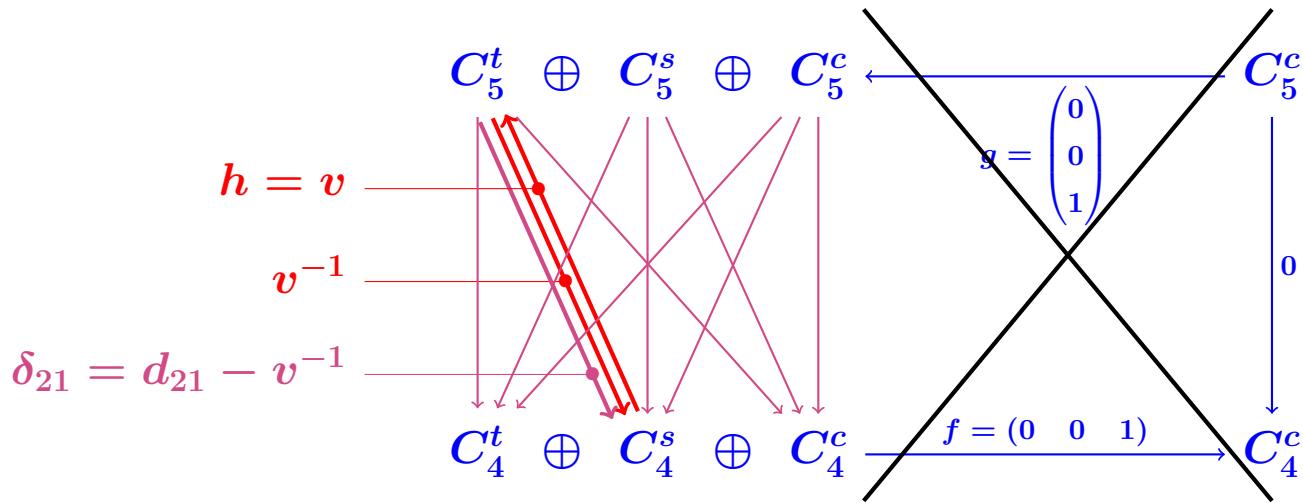
$$d^5 = d_{33}^5 - d_{31}^5 d_{21}^{5^{-1}} d_{23}^5$$

$$g_4 = \begin{pmatrix} -d_{21}^{4^{-1}} d_{23}^4 \\ 0 \\ 1 \end{pmatrix}$$

$$A_4 \oplus B_4 \oplus C_4 \xrightleftharpoons[f_4 = (0 \quad -d_{31}^5 d_{21}^{5^{-1}} \quad 1)]{} C_4$$

$$d^5 = \begin{pmatrix} d_{11}^5 & d_{12}^5 & d_{13}^5 \\ d_{21}^5 & d_{22}^5 & d_{23}^5 \\ d_{31}^5 & d_{32}^5 & d_{33}^5 \end{pmatrix}$$

Applying in our situation  $\Rightarrow$



$$h \mapsto h' = \left( \sum_{i=0}^{\infty} (-1)^i (v \delta_{21})^i \right) v \quad \text{and}$$

Initial situation:

$$\begin{array}{c}
 C_5^t \oplus C_5^s \oplus C_5^c \xleftarrow{\quad} C_5^c \\
 \downarrow v^{-1} \quad \downarrow \cong \quad \downarrow v = h \\
 C_4^t \oplus C_4^s \oplus C_4^c \xrightarrow{f = (0 \ 0 \ 1)} C_4^c
 \end{array}$$

$g = \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix}$

Final situation:

$$\begin{array}{c}
 C_5^t \oplus C_5^s \oplus C_5^c \xleftarrow{\quad} C_5^c \\
 \downarrow \quad \quad \quad \quad \quad \quad \downarrow \\
 C_4^t \oplus C_4^s \oplus C_4^c \xrightarrow{f = (0 \quad -d_{31}h' \quad 1)} C_4^c
 \end{array}$$

$g = \begin{pmatrix} -h'd_{23} \\ 0 \\ 1 \end{pmatrix}$

## Forman's Morse Subcomplex.

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Theorem:  $(C_*, d, \beta_*) =$  cellular complex.

$$V = \{(\sigma_i, \tau_i)\}_{i \in I} =$$

admissible discrete vector field on  $C_*$ .

$\rho = (f, g, h) : C_* \Rightarrow C_*^c$  the induced reduction.

Then  $\ker(dh + hd) = \ker(dv + vd) = \text{im}(g)$

is Forman's Morse subcomplex  $C_*^{\text{inv}}$ .

Forman's definition:  $C_*^{\text{inv}} := \{\sigma \text{ st } (1 - dv - vd)(\sigma) = \sigma\}$ .

Proposition:

$$\begin{aligned}
 h(\sigma) &= v(\sigma) - h\delta_{21}v(\sigma) \\
 &= v(\sigma) - h(d_{21} - v^{-1})v(\sigma) \\
 &= v(\sigma) - h(d - v^{-1})v(\sigma) \\
 hdv(\sigma) &= v(\sigma)
 \end{aligned}$$

Proof: For a source cell  $\sigma$ :

$$\begin{aligned}
 h(\sigma) &= v(\sigma) - v\delta_{21}v(\sigma) + v\delta_{21}v\delta_{21}v(\sigma) - \dots \\
 &= v(\sigma) - (v - v\delta_{21}v + \dots)\delta_{21}v(\sigma) \\
 &= v(\sigma) - h\delta_{21}v(\sigma)
 \end{aligned}$$

+ Everything null for  $\sigma$  target or critical cell  $\Rightarrow$  QED

Proof of Theorem:

$$\begin{aligned}
 & \sigma \in \ker(dv + vd) \\
 \Rightarrow & dv\sigma + vd\sigma = 0 \\
 \Rightarrow & hdv\sigma + hvd\sigma = 0 \\
 hv = 0 \Rightarrow & hdv\sigma = 0 \\
 hdv = v \Rightarrow & v\sigma = 0 \\
 \Rightarrow & vd\sigma = 0 \\
 \Rightarrow & dh\sigma + hd\sigma = 0 \\
 \Rightarrow & \sigma \in \ker(dh + hd) \\
 & = \text{im}(g)
 \end{aligned}$$

Conversely:

$$\begin{aligned}
 & \sigma \in \ker(dh + hd) \\
 \Rightarrow & \quad \sigma = g(\chi) \\
 \Rightarrow & \quad \sigma \text{ made of target and critical cells} \\
 dg = gd^c \Rightarrow & \quad d\sigma = g(\chi) \\
 \Rightarrow & \quad d\sigma \text{ made of target and critical cells} \\
 \Rightarrow & \quad dv\sigma + vd\sigma = 0 \\
 \Rightarrow & \quad \sigma \in \ker(dv + vd)
 \end{aligned}$$

Finally:  $\ker(dv + vd) = \ker(dh + hd)$  QED

## Vector Fields and Morphisms

Problem: Let  $C_*$  and  $C'_*$  be two cellular chain complexes respectively provided with vector fields  $V$  and  $V'$ .

Question: Right notion

of morphism  $\varphi : (C_*, V) \rightarrow (C'_*, V')$  ???

1. Not trivial.
2. Essential to master the Eilenberg-Zilber vector fields.
3. Quite amazing !!

Definition: A **cellular** morphism:

$$\varphi : (C_*, d_*, \beta_*) \rightarrow (C'_*, d'_*, \beta'_*)$$

is a chain complex morphism  $\varphi : (C_*, d_*) \rightarrow (C'_*, d'_*)$   
satisfying the extra condition:

For every  **$p$ -cell**  $\sigma \in \beta_p$ ,

$$\varphi(\sigma) \text{ is null or } \in \beta'_p.$$

$(C_*, d, \beta, V)$  and  $(C'_*, d', \beta', V')$

= cellular chain complexes

with respective admissible discrete vector fields  $V$  and  $V'$ .

Definition: A **vectorious morphism**:

$$\varphi : (C_*, d, \beta, V) \rightarrow (C'_*, d', \beta', V')$$

is a **cellular morphism**  $\varphi := (C_*, d, \beta) \rightarrow (C'_*, d', \beta')$

satisfying the **extra conditions**:

1. For every **critical cell**  $\chi \in \beta_p^c$ ,  $\varphi(\chi)$  is **null** or  $\in \beta_p'^c$ .
2. For every **target cell**  $\tau \in \beta_p^t$ ,  $\varphi(\tau)$  is **null** or  $\in \beta_p'^t$ .
3. **No condition at all for the source cells !!**

**Theorem:**  $\varphi : (C_*, d, \beta, V) \rightarrow (C'_*, d', \beta', V')$   
 $=$  vectorious morphism.

Then  $\varphi$  defines a morphism  $(\varphi, \varphi^c)$   
 between the corresponding reductions:

$$\begin{array}{ccc}
 C_* & \xrightarrow{\varphi} & C'_* \\
 h \text{ (loop)} & & h' \text{ (loop)} \\
 \uparrow f \quad g & & \uparrow f' \quad g' \\
 C_*^c & \xrightarrow{\varphi^c} & C'^c_*
 \end{array}
 \quad \text{with:}$$

$$\begin{aligned}
 d'^c \varphi^c &= \varphi^c d^c \\
 f' \varphi &= \varphi^c f \\
 g' \varphi^c &= \varphi g \\
 h' \varphi &= \varphi h
 \end{aligned}$$

Proof:Definition:

$$\lambda_\sigma = \begin{cases} 0 & \text{for target and critical cells,} \\ & \text{maximal length of a } V\text{-path} \\ & \text{starting from the source cell } \sigma. \end{cases}$$

Remember: Recursive formula:

$$\begin{aligned}
 h(\sigma) &= v(\sigma) - h(d - v^{-1})v(\sigma) \\
 &\Rightarrow hdv(\sigma) = v(\sigma) \\
 &\Rightarrow hd\tau = \tau \text{ for every target cell } \tau
 \end{aligned}$$

$$1. h'\varphi\sigma = \varphi h\sigma ??$$

Obvious for  $\sigma$  target or critical cell.

Assumed known for  $\lambda_\sigma < k$ .

Let  $\sigma$  be a source cell with  $\lambda_\sigma = k$ .

$$\varphi h\sigma = \varphi v\sigma - \varphi h(d - v^{-1})v\sigma$$

$$\text{OK for } (d - v^{-1})v\sigma \Rightarrow \quad = \varphi v\sigma - h'\varphi(d - v^{-1})v\sigma$$

$$\varphi d = d'\varphi \Rightarrow \quad = \varphi v\sigma - h'd'\varphi v\sigma + h'\varphi\sigma$$

$$\varphi v\sigma = \text{target cell} \Rightarrow \quad = h'\varphi\sigma$$

QED

$$2. \ g' \varphi^c = \varphi g \ ??$$

Remember:

$$\begin{array}{c}
 C_5^t \oplus C_5^s \oplus C_5^c \xleftarrow{\quad} C_5^c \\
 \downarrow \begin{matrix} \cong \\ d_{21} \\ h \end{matrix} \quad \downarrow \quad \downarrow \\
 C_4^t \oplus C_4^s \oplus C_4^c \xrightarrow{f = (0 \quad -d_{31}h \quad 1)} C_4^c
 \end{array}$$

$$g = \begin{pmatrix} -hd_{23} \\ 0 \\ 1 \end{pmatrix}$$

$$d_{33} - d_{31}hd_{23}$$

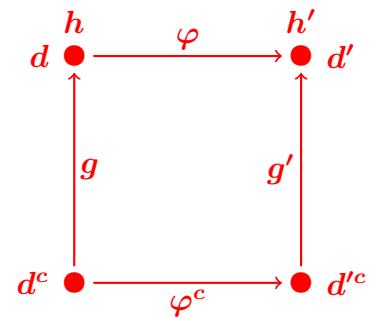
$$\text{For a critical cell } \chi: \quad g\chi = \chi - hd\chi = (1 - hd)\chi \quad \Rightarrow$$

$$\begin{aligned}
 &\Rightarrow \varphi g \chi = \varphi(1 - hd)\chi \\
 \varphi hd = h'd'\varphi &\Rightarrow = (1 - h'd')\varphi \chi \\
 \varphi \chi = \varphi^c \chi &\Rightarrow = g' \varphi^c \chi
 \end{aligned}$$

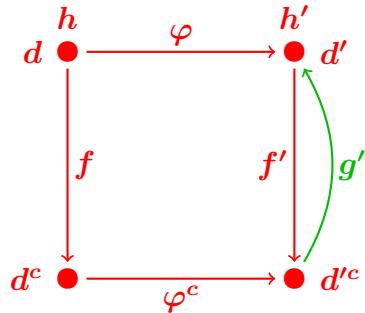
QED

$\varphi^c d^c = d'^c \varphi^c ??$

$$\begin{aligned}
 g' \varphi^c = \varphi g &\Rightarrow g' \varphi^c d^c = \varphi g d^c \\
 gd^c = dg &\Rightarrow = \varphi dg \\
 \varphi d = d' \varphi &\Rightarrow = d' \varphi g \\
 \varphi g = g' \varphi^c &\Rightarrow = d' g' \varphi^c \\
 d' g' = g' d'^c &\Rightarrow = g' d'^c \varphi^c \\
 g' \text{ injective} &\Rightarrow \varphi^c d^c = d'^c \varphi^c
 \end{aligned}$$



QED



$$f'\varphi = \varphi^c f \quad ??$$

$$g' \text{ injective} \Rightarrow [(f'\varphi = \varphi^c f) \Leftrightarrow (g'f'\varphi = g'\varphi^c f)]$$

$$\begin{aligned}
 g'f'\varphi &= (1 - d'h' - h'd')\varphi \\
 (d'\varphi = \varphi d) + (h'\varphi = \varphi h) \Rightarrow &= \varphi(1 - dh - hd) \\
 &= \varphi gf \\
 &= g'\varphi^c f
 \end{aligned}$$

QED

## Application:

Theorem: Given:  $\begin{array}{|c|} \varphi : X \rightarrow X' \\ \varphi' : Y \rightarrow Y' \end{array}$  = simplicial morphisms.

Then:  $\varphi$  and  $\varphi'$  induce a **morphism** between the **reductions**  
defined by the Eilenberg-Zilber vector fields:

$$\begin{array}{ccccc}
 h & \textcircled{\text{S}} & C_*(X \times Y) & \xrightarrow{\varphi \times \varphi'} & C_*(X' \times Y') & \textcircled{\text{S}} & h' \\
 & \uparrow & & & \uparrow & & \text{with:} \\
 & f & & & f' & & \\
 & \downarrow & & & \downarrow & & \\
 C_*(X \times Y)^c & \xrightarrow{(\varphi \times \varphi')^c} & C_*(X' \times Y')^c & & & &
 \end{array}$$

## Important Note:

$$\begin{array}{ccccc}
 h & \textcircled{\text{C}} & C_*(X \times Y) & \xrightarrow{\varphi \times \varphi'} & C_*(X' \times Y') & \textcircled{\text{C}} & h' \\
 & \uparrow & f & \downarrow g & & f' & \downarrow g' & \\
 & & & & & & & \\
 & & C_*(X \times Y)^c & \xrightarrow{(\varphi \times \varphi')^c} & C_*(X' \times Y')^c & & 
 \end{array}$$

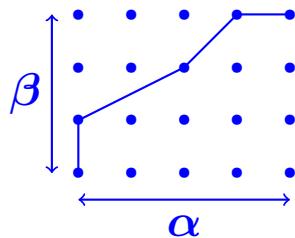
= same diagram as:

$$\begin{array}{ccccc}
 RM & \textcircled{\text{C}} & C_*(X \times Y) & \xrightarrow{\varphi \times \varphi'} & C_*(X' \times Y') & \textcircled{\text{C}} & RM' \\
 & \uparrow & AW & \downarrow EML & & AW' & \downarrow EML' & \\
 & & & & & & & \\
 & & C_*(X) \otimes C_*(Y) & \xrightarrow{\varphi \otimes \varphi'} & C_*(X') \otimes C_*(Y') & & 
 \end{array}$$

But not yet known!

## Proof:

Representation of a simplex of  $X \times Y$  via an s-path.



= subsimplex of  $\alpha \times \beta \subset (X \times Y)_7$

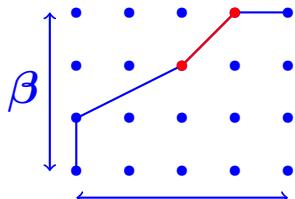
spanned by the vertices  $(0,0) - (0,1) - (2,2) - (3,3) - (4,3)$ .

The game first event “diagonal ↗”

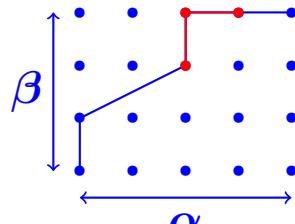
or “right-angle bend ↛”

determines the nature source, target or critical.

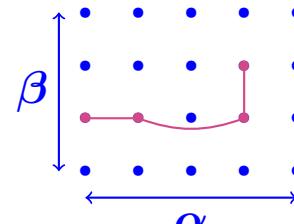
Examples:



source =  $\sigma$



target =  $\tau$



critical =  $\chi$

Here:

$$\partial_3(\tau) = \sigma$$

$$\Rightarrow v(\sigma) = -\tau \quad (\text{gradient})$$

$$v^{-1}(\tau) = -\sigma \quad (\text{cogradient})$$

Two maps  $\begin{vmatrix} \varphi : X \rightarrow X' \\ \varphi' : Y \rightarrow Y' \end{vmatrix}$  = simplicial morphisms.

Claim:

$\tau$  target cell in  $X \times Y \Rightarrow$

$(\varphi \times \varphi')(\tau)$  target or degenerate cell in  $X' \times Y'$

$\chi$  critical cell in  $X \times Y \Rightarrow$

$(\varphi \times \varphi')(\chi)$  critical or degenerate cell in  $X' \times Y'$

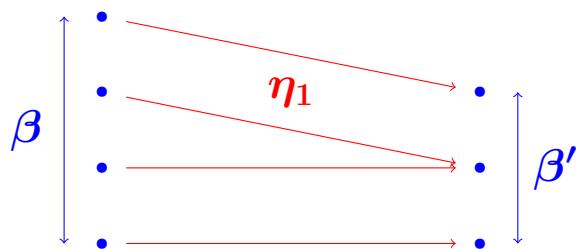
$$\alpha \in X^{ND} \Rightarrow \varphi(\alpha) = \eta\alpha'$$

for some multi-degeneracy  $\eta$  and some  $\alpha' \in X'^{ND}$ .

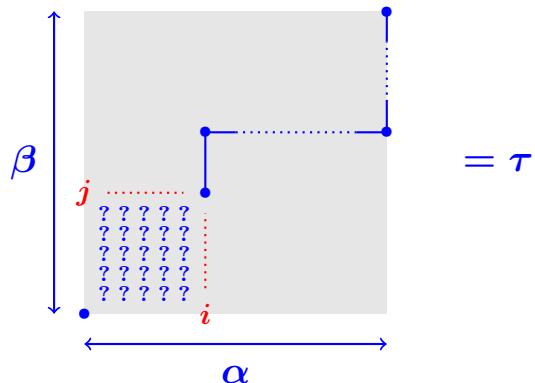
$$\beta \in Y^{ND} \Rightarrow \psi(\beta) = \theta\beta'$$

for some multi-degeneracy  $\theta$  and some  $\beta' \in Y'^{ND}$ .

Example:  $Y_3 \ni \beta \mapsto \theta\beta' = \eta_1\beta' \in Y'_3$ :



General **shape** of an Eilenberg-Zilber target cell:



$$(\varphi \times \psi)(\alpha \times \beta) = (\eta \alpha' \times \theta \beta')$$

If no index of  $\eta$  is  $\geq i$  and no index of  $\theta$  is  $\geq j$ ,

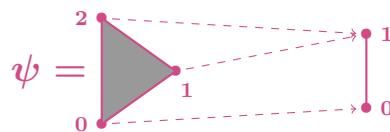
then  $(\varphi \times \psi)(\tau)$  has the same shape and therefore is a target cell.

Otherwise  $(\varphi \times \psi)(\tau)$  is degenerate.

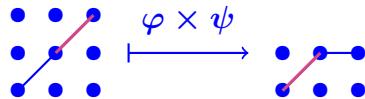
Same argument for critical simplices.  $\Rightarrow$  QED

Typical accidents with source cells.

$\alpha = \text{id} : \Delta^2 \rightarrow \Delta^2$  and  $\psi : \Delta^2 \rightarrow \Delta^1$  as below:



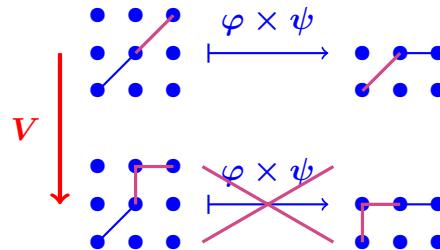
1)



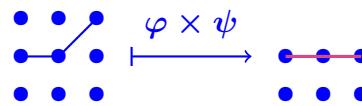
Then  $(\varphi \times \psi)(\text{source}) = \text{source}$

but for reasons which do not match!

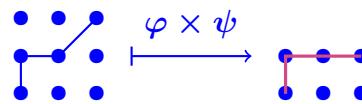
Compare corresponding target cells.



2) The image of a source cell can be a critical cell:



or target:



But we don't care about source cells!

## Eilenberg-Zilber formulas

Theorem: The Eilenberg-Zilber vector field  
 previously described  
 gives the standard Eilenberg-Zilber reduction.

Standard Eilenberg-Zilber reduction:

$$EZ : \textcolor{red}{RM} \hookrightarrow C_*(X \times Y) \xleftarrow[\textcolor{red}{AW}]{} C_*(X) \otimes C_*(Y)$$

$AW$  = Alexander-Whitney

$EML$  = Eilenberg-MacLane

$RM$  = Rubio-Morace

$$\textcolor{red}{EZ} = \textcolor{red}{AW} + \textcolor{red}{EML} + \textcolor{red}{RM}:$$

$$\textcolor{red}{AW}(x_p \times y_p) = \sum_{i=0}^p \partial_{p-i+1} \cdots \partial_p x_p \otimes \partial_0 \cdots \partial_{p-i-1} y_p$$

$$\textcolor{red}{EML}(x_p \otimes y_q) = \sum_{(\eta, \eta') \in \text{Sh}(p, q)} \varepsilon(\eta, \eta') (\eta' x_p \times \eta y_q)$$

$$\textcolor{red}{RM}(x_p \times y_p) = \sum_{\substack{0 \leq r \leq p-1 \\ 0 \leq s \leq p-r-1 \\ (\eta, \eta') \in \text{Sh}(s+1, r)}} (-1)^{p-r-s} \varepsilon(\eta, \eta') \dots$$

$$\dots (\uparrow^{p-r-s}(\eta') \eta_{p-r-s-1} \partial_{p-r+1} \cdots \partial_p x_p \times \dots$$

$$\dots \uparrow^{p-r-s}(\eta) \partial_{p-r-s} \cdots \partial_{p-r-1} y_p)$$

Plan of proof:

- 1) Prove the  $\textcolor{red}{RM}$ -formula for the homotopy induced by the Eilenberg-Zilber vector field.
- 2) Use the diagram ( $\textcolor{blue}{h}' = \textcolor{red}{RM}$ ,  $\textcolor{violet}{f} = ?AW$ ,  $\textcolor{blue}{g} = ?EML$ ):

$$\begin{array}{c}
 C_5^t \oplus C_5^s \oplus C_5^c \xleftarrow{\quad g = \begin{pmatrix} -\textcolor{red}{h}'d_{23} \\ 0 \\ 1 \end{pmatrix}} C_5^c \\
 \downarrow \textcolor{red}{d}_{21} \qquad \qquad \qquad \downarrow \textcolor{red}{d}_{33} - d_{31}\textcolor{red}{h}'d_{23} \\
 C_4^t \oplus C_4^s \oplus C_4^c \xrightarrow{\quad f = (0 \quad -d_{31}\textcolor{red}{h}' \quad 1) \quad} C_4^c
 \end{array}$$

The diagram illustrates a commutative square of chain complexes. The top row consists of three copies of the complex  $C_5$  (with components  $t$ ,  $s$ , and  $c$ ) joined at their middle term. The bottom row consists of three copies of the complex  $C_4$  (with components  $t$ ,  $s$ , and  $c$ ) joined at their middle term. Red arrows indicate differentials:  $d_{21}$  from  $C_5^t$  to  $C_4^s$ ,  $d_{33} - d_{31}\textcolor{red}{h}'d_{23}$  from  $C_5^c$  to  $C_4^c$ , and  $\textcolor{red}{h}'$  from  $C_5^t$  to  $C_4^s$ . A red double-headed arrow between  $C_5^t$  and  $C_4^s$  is labeled  $\cong$ .

which produces:  $\textcolor{red}{h}' \Rightarrow \textcolor{violet}{f}$  and  $\textcolor{blue}{g}$ .

## Plan for the $RM$ -formula:

1) True for target and critical simplices.

2) Use the recursive formula:

$$h'(\sigma) = v(\sigma) - h'\delta_{21}v(\sigma)$$

$$\begin{array}{c}
 C_5^t \oplus C_5^s \oplus C_5^c \xleftarrow{\quad} C_5^c \\
 \downarrow \begin{matrix} \cong \\ d_{21} \\ h' \end{matrix} \quad \downarrow \quad \downarrow \quad \downarrow \\
 C_4^t \oplus C_4^s \oplus C_4^c \xrightarrow{\quad f = (0 \quad -d_{31}h' \quad 1) \quad} C_4^c
 \end{array}$$

$$g = \begin{pmatrix} -h'd_{23} \\ 0 \\ 1 \end{pmatrix}$$

$$d_{33} - d_{31}h'd_{23}$$

with  $\delta_{21} = d_{21} - v^{-1}$ .

Study of the generic term of the  $RM$ -formula:

$$(\uparrow^{p-r-s}(\eta')\eta_{p-r-s-1}\partial_{p-r+1} \cdots \partial_p x_p \times \dots$$

$$\dots \uparrow^{p-r-s}(\eta)\partial_{p-r-s} \cdots \partial_{p-r-1} y_p)$$

$\uparrow^{p-r-s}(\eta')\eta_{p-r-s-1}$  and  $\uparrow^{p-r-s}(\eta) =$  known degeneracy part.

$\partial_{p-r+1} \cdots \partial_p x_p$  and  $\partial_{p-r-s} \cdots \partial_{p-r-1} y_p =$  problematic part.

$$(\eta, \eta') \in \text{Sh}(s+1, r)$$

$$\Rightarrow \text{Ind}_{\eta} \cup \text{Ind}_{\eta'} = \{0, \dots, r+s\}$$

$$\Rightarrow \text{Ind}_{\uparrow^{p-r-s}(\eta')\eta_{p-r-s-1}} \cup \text{Ind}_{\uparrow^{p-r-s}(\eta)} = \{p-r-s-1, \dots, p\}$$

Collision lemma:  $(\eta x, \eta'y) \in (X \times Y)_p$

$$+ \text{ Ind}_{\eta} \cup \text{Ind}_{\eta'} = \{k, \dots, p-1\}$$

+  $\eta_i$  present in  $x$  or  $y$  with  $i \geq k$

$\Rightarrow (\eta x, \eta'y)$  degenerate.

Proof by examples:  $p = 10$ ,  $k = 6$ ,  $i = 6$ ,

$(\eta_9\eta_7x, \eta_8\eta_6y)$  degenerate?

$$\eta_6 \in x \Rightarrow (\eta_9\eta_7x, \eta_8\eta_6y) = (\eta_9\eta_7\eta_6x', \eta_8\eta_6y) = \eta_6(\eta_8\eta_6x', \eta_7y)$$

$$\eta_6 \in y \Rightarrow (\eta_9\eta_7x, \eta_8\eta_6y) = (\eta_9\eta_7x, \eta_8\eta_6\eta_6y') = \eta_7(\eta_8x, \eta_7\eta_6y)$$

Putting both components in **canonical form**

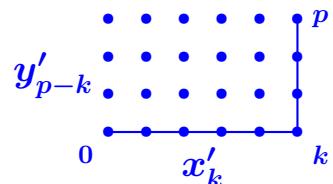
$\Rightarrow$  a **common factor**.

QED

Application:  $RM(x_p \times y_p) = 0$  if  $(x_p \times y_p)$  = critical cell.

Proof:  $(x_p, y_p) = (\eta_{p-1} \cdots \eta_k x'_k \times \eta_{k-1} \cdots \eta_0 y'_{p-k})$

1) Examine  $\underbrace{\partial_{p-r+1} \cdots \partial_p}_{r} \underbrace{\eta_{p-1} \cdots \eta_k}_{p-k} x'_k$ .



$r < p - k \Rightarrow$  there remains  $\eta_{p-r-1} \cdots x'_k$

but  $\{p - r - 1 \geq p - r - s - 1 \Rightarrow \text{collision}\} \Rightarrow r \geq p - k$ .

2) Examine  $\underbrace{\partial_{p-r-s} \cdots \partial_{p-r-1}}_s \underbrace{\eta_{k-1} \cdots \eta_0}_{k} y'_{p-k}$

$s < k \Rightarrow$  there remains  $\eta_{k-s-1} \cdots y'_{p-k}$

$\Rightarrow k - s - 1 < p - r - s - 1 \Rightarrow r < p - k$ .



QED

Same sort of argument  $\Rightarrow$

$$\textcolor{blue}{RM}(x_p \times y_p) = 0 \text{ for } (x_p \times y_p) = \text{target cell}.$$

$$\Rightarrow \textcolor{green}{h}(x_p \times y_p) = \textcolor{blue}{RM}(x_p \times y_p) = 0$$

for  $(x_p \times y_p) = \text{critical or target cells.}$

If  $(x_p \times y_p) = \text{source cell}$ ,  $\textcolor{red}{h}(x_p \times y_p) \stackrel{??}{=} \textcolor{blue}{RM}(x_p \times y_p)$ .

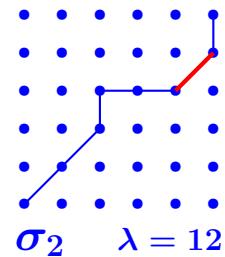
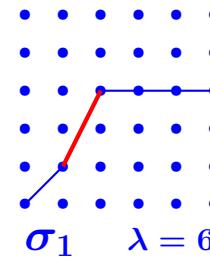
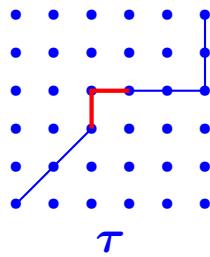
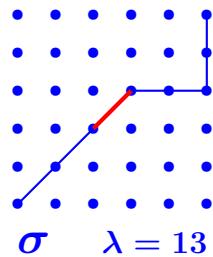
Remember:

$$\lambda_\sigma = \begin{cases} 0 & \text{for target and critical cells,} \\ & \text{maximal length of a } V\text{-path} \\ & \text{starting from the source cell } \sigma. \end{cases}$$

Remember: Recursive formula for source cells:

$$h(\sigma) = v(\sigma) - h(d_{21} - v^{-1})v(\sigma)$$

## Example:



$$v(\sigma) = -\tau$$

$$d_{21}(\tau) = -\sigma + \sigma_1 + \sigma_2$$

$$v^{-1}(\tau) = -\sigma$$

$$\Rightarrow h(\sigma) = -\tau + h(\sigma_1) + h(\sigma_2)$$

Then we prove the summands of  $RM(\sigma_1)$

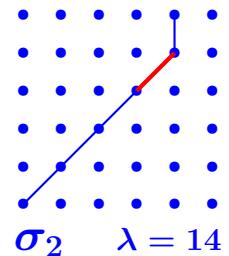
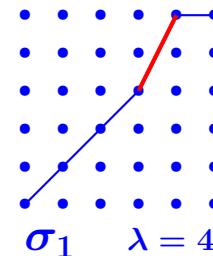
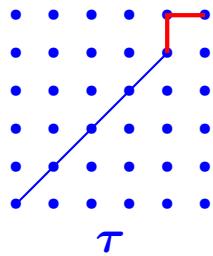
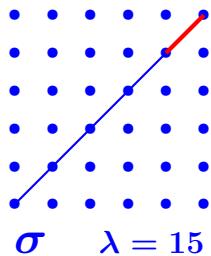
are exactly divided into

$-\tau$	{
the summands of $h(\sigma_1)$	

the summands of $h(\sigma_2)$	}
-------------------------------	---

$\Rightarrow$  Recursive proof

Easiest and last case:



$$x_p = y_p = \Delta^p$$

$$\sigma = (x_p \times y_p)$$

$$\tau = (\eta_{p-1} x_p \times \eta_p y_p)$$

$$\sigma_1 = (x_p \times \eta_{p-1} \partial_{p-1} y_p)$$

$$\sigma_2 = (\eta_{p-1} \partial_p x_p \times y_p)$$

We recursively assume

$$h(\sigma_1) = RM(\sigma_1) \text{ and } h(\sigma_2) = RM(\sigma_2)$$

Then we prove  $RM(\sigma) = \pm\tau \pm h(\sigma_1) \pm h(\sigma_2) =: h(\sigma)$  QED

***RM*-formula:**

Indices =

$$\{0 \leq r \leq p-1, 0 \leq s \leq p-r-1, (\eta|\eta') \in \text{Sh}(s+1, r)\}$$

Generic term of the ***RM*-formula**:

$$\pm(\uparrow^{p-r-s}(\eta')\eta_{p-r-s-1}\partial_{p-r+1}\cdots\partial_p x_p \times \dots)$$

$$\dots\uparrow^{p-r-s}(\eta)\partial_{p-r-s}\cdots\partial_{p-r-1}y_p)$$

We must prove  $\text{RM}(\sigma) = \pm\tau \pm \text{RM}(\sigma_1) \pm \text{RM}(\sigma_2)$

$$\sigma = (x_p \times y_p)$$

$$\tau = (\eta_{p-1}x_p \times \eta_p y_p)$$

$$\sigma_1 = (x_p \times \eta_{p-1}\partial_{p-1}y_p)$$

$$\sigma_2 = (\eta_{p-1}\partial_p x_p \times y_p)$$

Elementary applications of Collision Lemma  $\Rightarrow$

$$\tau = \mathbf{RM}_{r=s=0}(\sigma)$$

$$\mathbf{RM}(\sigma_1) = \mathbf{RM}_{r=0,s>0}(\sigma) \quad \mathbf{RM}(\sigma_2) = \mathbf{RM}_{r>0}(\sigma)$$

$$\mathbf{RM} = \mathbf{RM}_{r=s=0} + \mathbf{RM}_{r=0,s>0} + \mathbf{RM}_{r>0}(\sigma) \quad \Rightarrow \quad QED$$

Analogous proof for the general case  $\Rightarrow$

OK for  $\mathbf{RM} = h_{\text{EZ-vector-field}}$

$h_{\text{EZ-vector-field}} \Rightarrow f_{\text{EZ-vector-field}}$  and  $g_{\text{EZ-vector-field}}$

the EZ-vector-field defines the standard EZ-reduction.

QED

# The END

```
;; Clock  
Computing  
<TnPr <Tn  
End of computing.  
  
;; Clock -> 2002-01-17, 19h 25m 36s.  
Computing the boundary of the generator 19 (dimension 7) :  
<TnPr <TnPr <TnPr S3 <<Abar[2 $1][2 $1]>>> <<Abar>>>  
End of computing.  
  
Homology in dimension 6 :  
  
Component 2/122  
  
---done---  
;; Clock -> 2002-01-17, 19h 27m 15s
```

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ETH Zurich, June 2012