

From a **constructiveness problem**

to

Discrete Vector Fields

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;; Clock  
Computing  
<TnPr <TnPr  
End of computing.
```

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;; Clock -> 2002-01-17, 19h 25m 36s.  
Computing the boundary of the generator 19 (dimension 7) :  
<TnPr <TnPr <TnPr S3 <<Abar[2 S1][2 S1]>>> <<Abar>>> <<Abar>>>  
End of computing.
```

Homology in dimension 6 :

Component Z/12Z

---done---

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;; Clock -> 2002-01-17, 19h 27m 15s
```

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Mulhouse, May 24, 2012*

Semantics of colours:

Blue = “Standard” Mathematics

Red = Constructive, effective,

algorithm, machine object, ...

Violet = Problem, difficulty,

obstacle, disadvantage, ...

Green = Solution, essential point,

mathematicians, ...

Dark Orange = Fuzzy objects.

Pale grey = Hyper-Fuzzy objects.

Plan.

1. Constructiveness.
2. Algebraic Topology and Constructiveness.
3. Homological Reduction.
4. Homotopy and Vector Fields.
5. Discrete Vector Fields.
6. Eilenberg-Zilber and Vector Fields.
7. Twisted Eilenberg-Zilber.
8. Results.

1/8. **Constructiveness** ?

Standard example:

Problem: Determine real numbers α and β satisfying:

1. $\alpha, \beta \in \mathbb{R} - \mathbb{Q}$,
2. $\alpha^\beta \in \mathbb{Q}$.

Solution: Consider $\gamma = \sqrt{2}^{\sqrt{2}}$.

Then: If $\gamma \in \mathbb{Q}$, $\alpha = \beta = \sqrt{2} = \text{solution}$.

Else $\alpha = \gamma$ and $\beta = \sqrt{2} = \text{solution}$. $[(\sqrt{2}^{\sqrt{2}})^{\sqrt{2}} = 2]$

But which one of $(\sqrt{2}, \sqrt{2})$ or $(\sqrt{2}^{\sqrt{2}}, \sqrt{2})$

finally is a solution ???

2/8. Standard Algebraic Topology is not constructive.

Standard Algebraic Topology:

Topological object \mapsto Algebraic object

Topological Object \mapsto Chain Complex \mapsto Homology group

Chain complex:

$$\begin{array}{ccccccc}
 \cdots & \xleftarrow{d_{p-2}} & C_{p-2} & \xleftarrow{d_{p-1}} & C_{p-1} & \xleftarrow{d_p} & C_p & \xleftarrow{d_{p+1}} & C_{p+1} & \xleftarrow{d_{p+2}} & \cdots \\
 & & & & & & & & & & \\
 & & & \underbrace{\hspace{1.5cm}} & & \underbrace{\hspace{1.5cm}} & & \underbrace{\hspace{1.5cm}} & & \underbrace{\hspace{1.5cm}} & \\
 & & & d_{p-2}d_{p-1} = 0 & & d_{p-1}d_p = 0 & & d_p d_{p+1} = 0 & & d_{p+1}d_{p+2} = 0 &
 \end{array}$$

$$dd = 0 \quad \Leftrightarrow \quad \ker d_p \supset \operatorname{im} d_{p+1} \quad \Rightarrow \quad H_p := \frac{\ker d_p}{\operatorname{im} d_{p+1}}$$

Standard work in Algebraic Topology:

1. Some topological space X .
2. \Rightarrow Chain complex $(C_*(X), d_*)$.
3. \Rightarrow Homology groups $\{H_p(X)\}_{p \in \mathbb{N}}$.

But most often $C_*(X)$ not of finite type.

\Rightarrow The groups $H_p(X)$ not directly calculable.

\Rightarrow How to determine these $H_p(X)$???

Standard “solution”:

Use the numerous exact sequences and spectral sequences
of the relevant textbooks.

Not enough !!

Typical example: Serre computing $\pi_6(S^3) = ???$ in 1950:

Serre constructs a complicated space Y

and his famous spectral sequence

produces the exact sequence:

$$0 \longleftarrow \mathbb{Z}/6 \longleftarrow H_6(Y) \longleftarrow \mathbb{Z}/2 \longleftarrow 0$$

$$\Rightarrow H_6(Y) = \mathbb{Z}/12 \text{ or } \mathbb{Z}/2 \oplus \mathbb{Z}/6 \text{ ???}$$

Analysis of the problem:

Standard Algebraic Topology is non-constructive.

Typical example:

Exact translation of $H_6(Z) \text{ “=” } \mathbb{Z}/6$:

There exists an isomorphism $H_6(C_*(Z)) \cong \mathbb{Z}/6$,

but producing such an isomorphism

remains unsolved

in standard Algebraic Topology.

Elementary \mathbb{Z} -linear algebra \Rightarrow

such an isomorphism $\varphi : H_p(C_*) \cong A$

for A some \mathbb{Z} -module of finite type

is **elementary** when C_* is of finite type.

What about the usual case with $C_*(X)$ not of finite type ???

3/8. Homological Reduction.

Notion of Reduction

between **big** chain complexes and **small** chain complexes.

In particular between **chain complexes not of finite type**

and **chain complexes of finite type**.

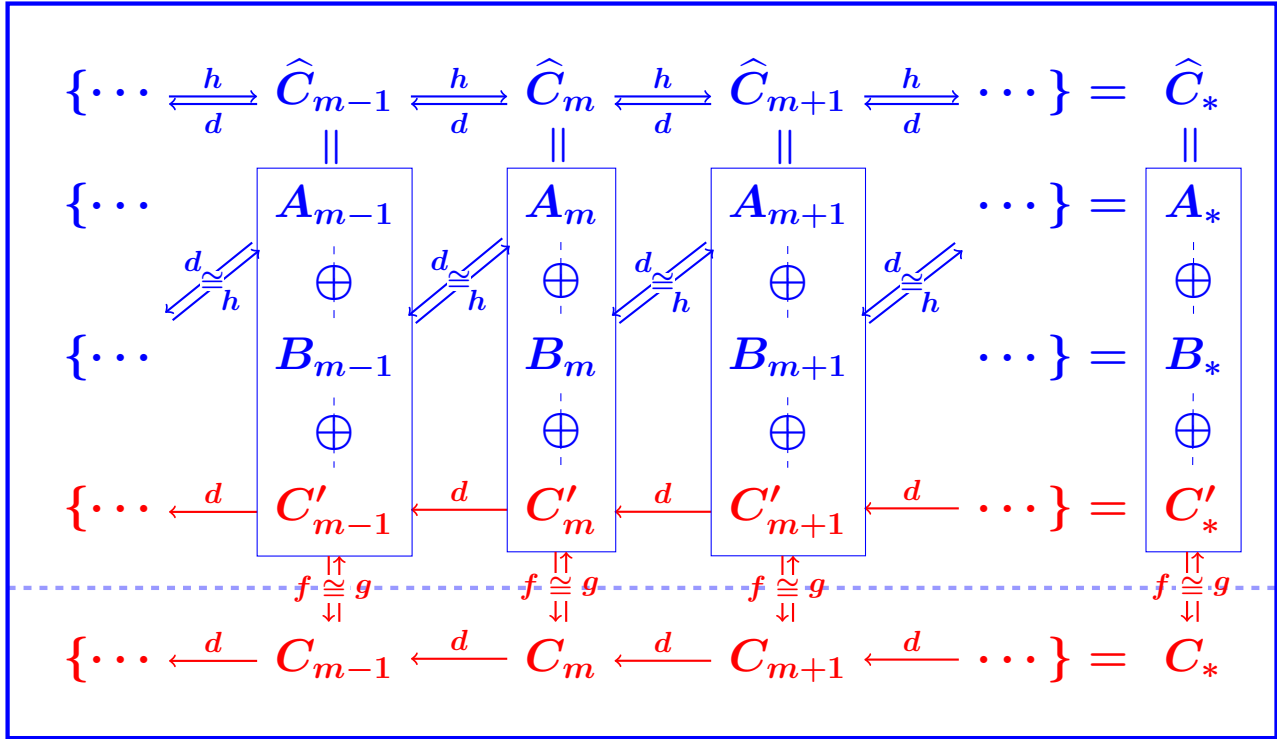
Definition: A (homological) reduction is a diagram:

$$\rho = (f, g, h) = \boxed{h \circlearrowleft (\widehat{C}_*, \widehat{d}_*) \begin{array}{c} \xleftarrow{g} \\ \xrightarrow{f} \end{array} (C_*, d_*)}$$

with:

1. \widehat{C}_* and C_* = chain complexes.
2. f and g = chain complex morphisms.
3. h = homotopy operator (degree +1).
4. $fg = \text{id}_{C_*}$ and $d_{\widehat{C}}h + hd_{\widehat{C}} + gf = \text{id}_{\widehat{C}_*}$.
5. $fh = 0$, $hg = 0$ and $hh = 0$.

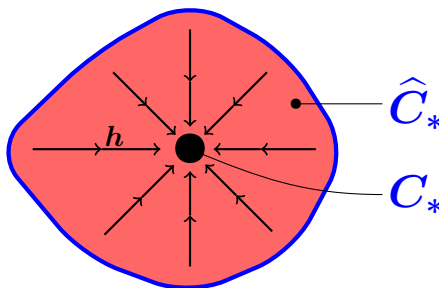
Meaning = Reduction Diagram:



$$\text{In } \rho = (f, g, h) = \boxed{h \circlearrowleft (\hat{C}_*, \hat{d}_*) \xrightleftharpoons[f]{g} (C_*, d_*)},$$

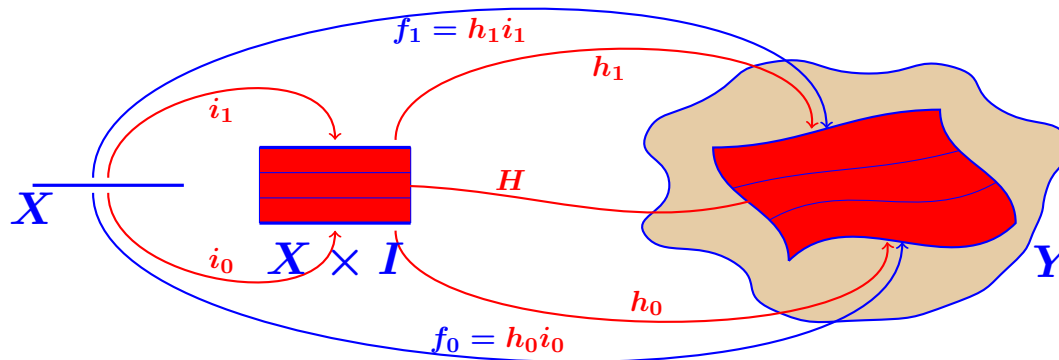
the **most important** component is h .

To be understood as a **contraction** $\hat{C}_* \Rightarrow C_*$.



The **homotopy** h **contracts** the “cherry” \hat{C}_*

onto the “stone” C_* .



General notion of **homotopy**.

Given: $f_0, f_1 : X \rightarrow Y$.

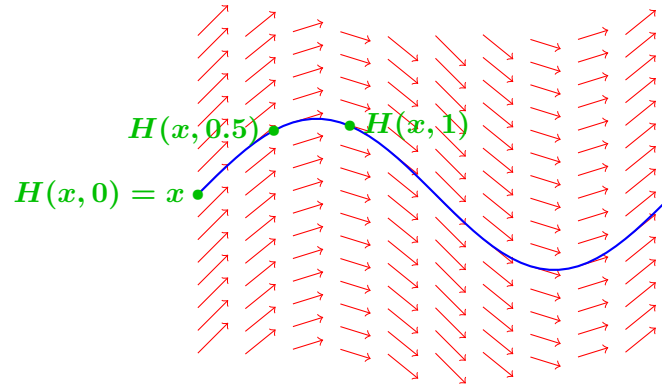
Definition: f_0 **homotope** to $f_1 \iff$

$\exists H : X \times I \rightarrow Y$ satisfying:

$$f_0(x) = H(x, 0)$$

$$f_1(x) = H(x, 1)$$

4/8. Vector field \Rightarrow Flow \Rightarrow Homotopy.



Vector field = $V : x \mapsto V(x) \in T_x X$.

\Rightarrow Corresponding flow $\Phi(x, t)$:

$$\text{solution of } \left\{ \begin{array}{l} \frac{\partial \Phi(x, t)}{\partial t} = V(\Phi(x, t)) \\ \Phi(x, 0) = x \end{array} \right\}.$$

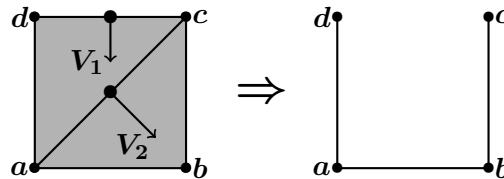
\Rightarrow Auto-Homotopy $H(x, t) := \Phi(x, t)$ for $t \in [0, 1]$.

5/8. Corresponding **combinatorial** notion ??

Example of **contraction**: 

Appropriate **discrete** vector field

on a **triangulation** of the square:



V_1 **contracts** $acd \Rightarrow da || ac =$ cancels the pair (cd, dac) .

V_2 **contracts** $abc \Rightarrow ab || bc =$ cancels the pair (ac, abc) .

\Rightarrow There remains only $da || ab || bc$.

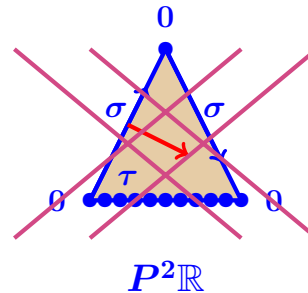
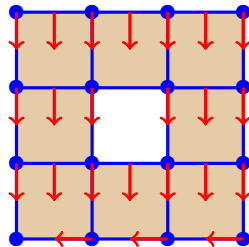
C = cellular complex = collection of cells satisfying...

Definition: A **vector field**

is a set of cells $V = \{(\sigma_i, \tau_i)_{i \in v}\}$ satisfying:

- $(\forall i)$ σ_i is **regular** face of τ_i .
- $(\forall i \neq j)$ $\sigma_i \neq \sigma_j \neq \tau_i \neq \tau_j$.

Examples:



σ_i **regular** face of $\tau_i \Leftrightarrow$ incidence number $\varepsilon(\sigma_i, \tau_i) = \pm 1$.

C = Cellular chain complex.

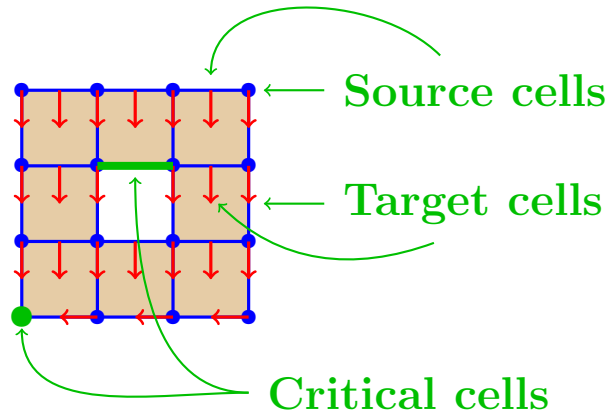
$V = \{(\sigma_i, \tau_i)_{i \in v}\} =$ Vector field.

Definition: A **critical p -cell** is a p -cell

which **does not** occur in V .

Other **cells** divided in **source cells** and **target cells**.

Example:



C = Cellular complex.

$V = \{(\sigma_i, \tau_i)_{i \in v}\}$ = vector field.

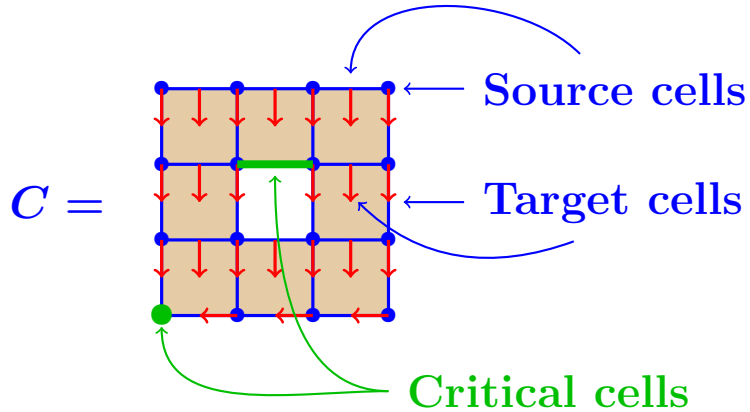
Theorem: A “critical” chain complex $C_*^c = ((\beta_p^c)_{p \in \mathbb{Z}}, d^c)$
can be constructed:

- β_p^c = the set of critical p -cells of V .
- $d_p^c : \mathbb{Z}[\beta_p^c] \rightarrow \mathbb{Z}[\beta_{p-1}^c]$
an appropriate “critical” differential
deduced from the initial differential d
and the vector field V .

Also a canonical reduction $\rho : C_* \Rightarrow C_*^c$ is provided.

\Rightarrow Any homological problem in C_* can be solved in C_*^c .

Simple example.



$$C_* = \{0 \leftarrow \mathbb{Z}^{16} \leftarrow \mathbb{Z}^{24} \leftarrow \mathbb{Z}^8 \leftarrow 0\}$$

Theorem \Rightarrow

$$\rho : C_* \twoheadrightarrow C_*^c = \begin{array}{|c|} \hline \begin{array}{c} \text{---} \\ \curvearrowright \\ \bullet \end{array} \\ \hline \end{array} \begin{array}{l} d_1^c \\ d_1^c \end{array} = \mathbb{Z} \xleftarrow{d_1^c=0} \mathbb{Z} = \text{Circle}$$

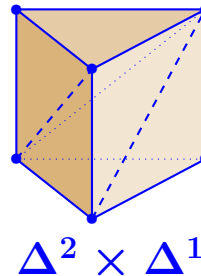
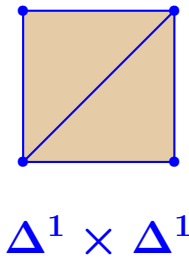
6/8. Eilenberg-Zilber Vector Field.

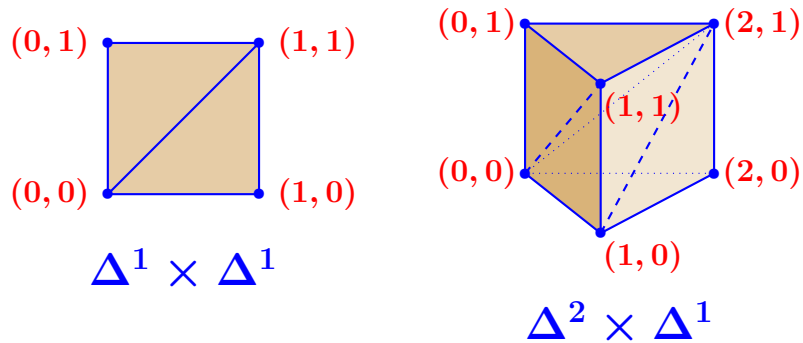
Vector Fields \Rightarrow Eilenberg-Zilber

\Rightarrow Twisted Eilenberg-Zilber \Rightarrow Serre spectral sequence

\Rightarrow Eilenberg-Moore spectral sequence.

Main problem: Triangulation of $\Delta^p \times \Delta^q$???

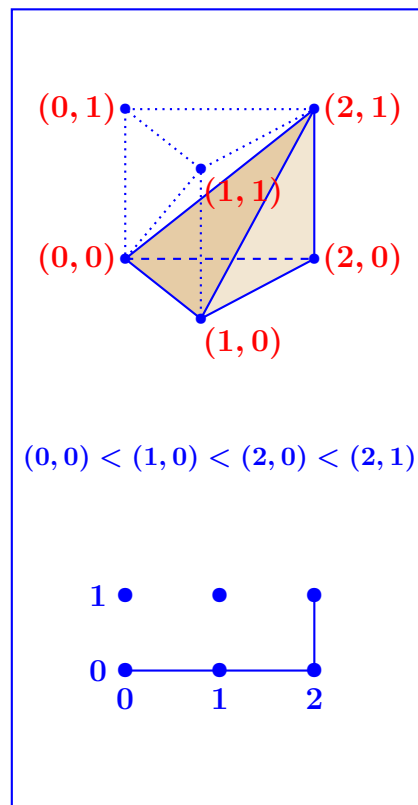
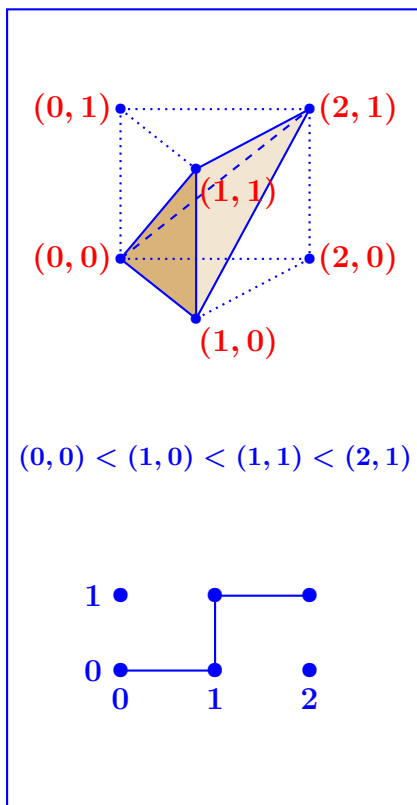
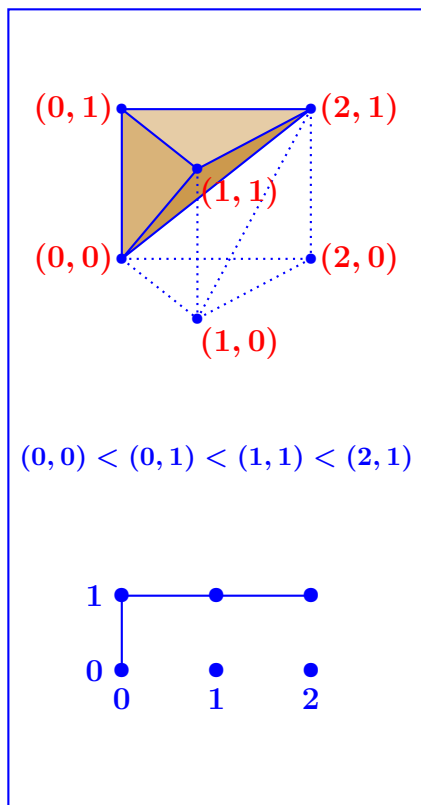




Two Δ^2 in $\Delta^1 \times \Delta^1$: $(0,0) < (0,1) < (1,1)$
 $(0,0) < (1,0) < (1,1)$

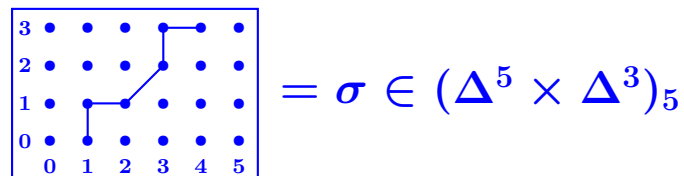
Three Δ^3 in $\Delta^2 \times \Delta^1$: $(0,0) < (0,1) < (1,1) < (2,1)$
 $(0,0) < (1,0) < (1,1) < (2,1)$
 $(0,0) < (1,0) < (2,0) < (2,1)$

Rewriting the triangulation of $\Delta^2 \times \Delta^1$.

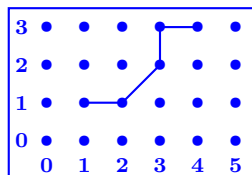


“Seeing” the **triangulation** of $\Delta^5 \times \Delta^3$.

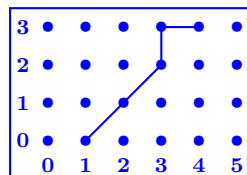
Example of 5-simplex :



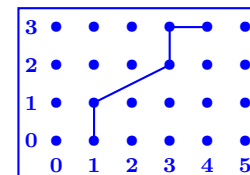
\Rightarrow 6 faces:



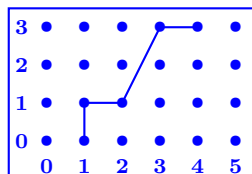
$\partial_0 \sigma$



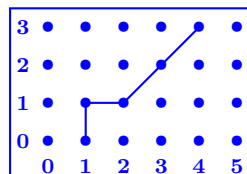
$\partial_1 \sigma$



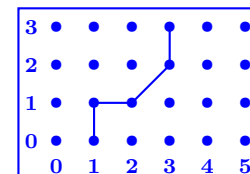
$\partial_2 \sigma$



$\partial_3 \sigma$

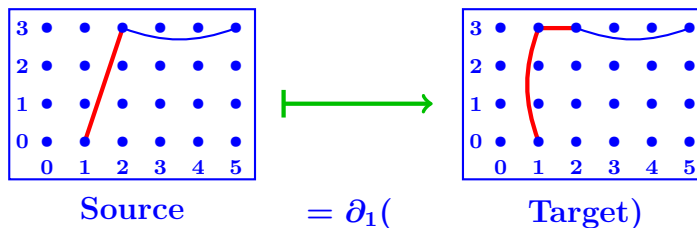
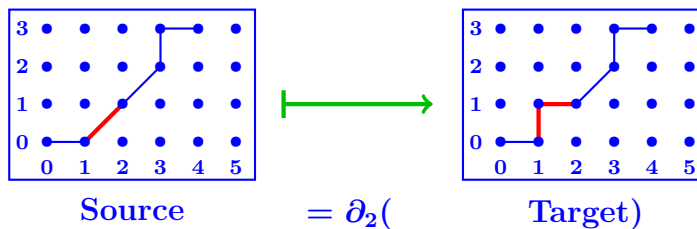




$\partial_4 \sigma$



$\partial_5 \sigma$

⇒ **Canonical discrete vector field** for $\Delta^5 \times \Delta^3$.



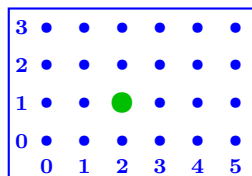
Recipe: First “event” = **Diagonal step** =  ⇒ **Source cell**.
 = **(-90°)-bend** =  ⇒ **Target cell**.

Critical cells ??

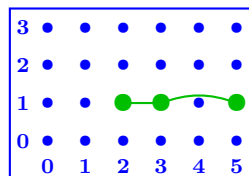
Critical cell = cell without any “event”

= without any diagonal or -90° -bend.

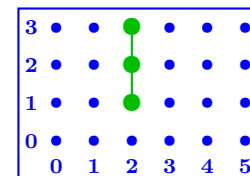
Examples.



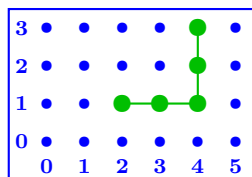
$$\Delta_2^0 \times \Delta_1^0$$



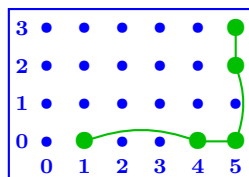
$$\Delta_{2,3,5}^2 \times \Delta_1^0$$



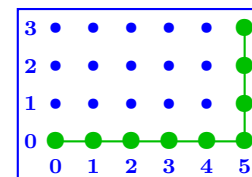
$$\Delta_2^0 \times \Delta_{1,2,3}^2$$



$$\Delta_{2,3,4}^2 \times \Delta_{1,2,3}^2$$



$$\Delta_{1,4,5}^2 \times \Delta_{0,2,3}^2$$



$$\Delta_{0,1,2,3,4,5}^5 \times \Delta_{0,1,2,3}^3$$

Conclusion:

$$C_*^c = C_*(\Delta^5) \otimes C_*(\Delta^3)$$

Fundamental theorem of vector fields \Rightarrow

Canonical Homological Reductions:

$$\rho : C_*(\Delta^5 \times \Delta^3) \Rightarrow C_*(\Delta^5) \otimes C_*(\Delta^3)$$

$$\rho : C_*(\Delta^p \times \Delta^q) \Rightarrow C_*(\Delta^p) \otimes C_*(\Delta^q)$$

$$p = q = 10 \Rightarrow 16,583,583,743 \text{ vs } 4,190,209$$

More generally: X and $Y =$ simplicial sets.

An admissible discrete vector field

is canonically defined on $C_*(X \times Y)$.

\Rightarrow Critical chain complex $C_*^c(X \times Y) = C_*(X) \otimes C_*(Y)$.

Eilenberg-Zilber Theorem: Canon. homological reduction:

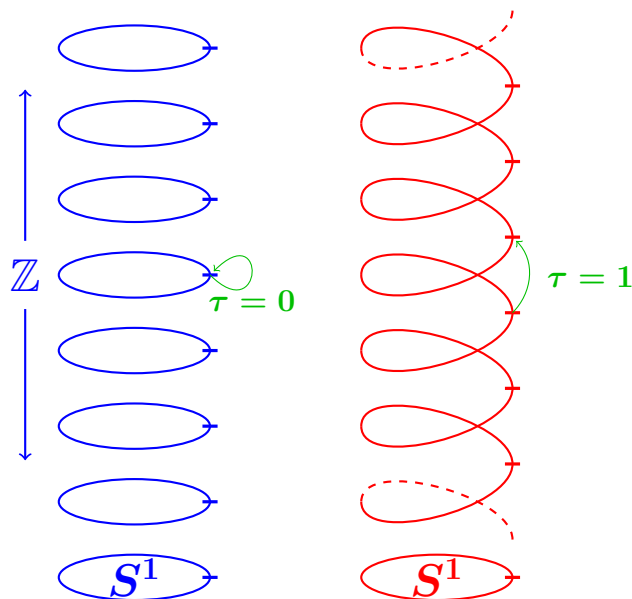
$$\rho_{EZ} : C_*(X \times Y) \twoheadrightarrow C_*^c(X \times Y) = C_*(X) \otimes C_*(Y)$$

\Rightarrow Künneth theorem to compute $H_*(X \times Y)$.

7/8. Twisted Eilenberg-Zilber.

Notion of **twisted product**.

Simplest example: $\mathbb{Z} \times S^1$ vs $\mathbb{Z} \times_{\tau} S^1 = \mathbb{R}$:



General notion of **twisted product**: B = base space.

F = fibre space.

G = structural group.

Action $G \times F \rightarrow F$.

$\tau : B \rightarrow G$ = Twisting function.

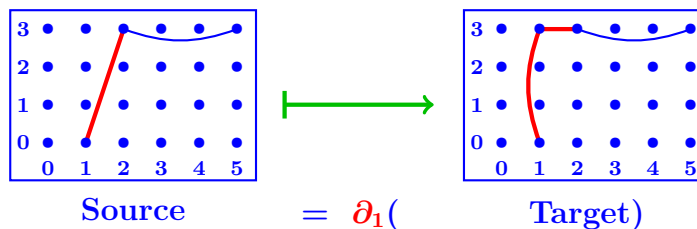
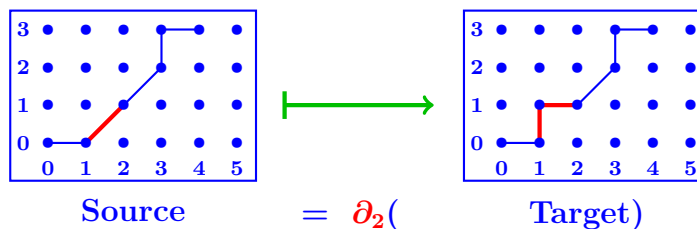
Structure of $F \times_{\tau} B$:

$$\partial_i(\sigma_f, \sigma_b) = (\partial_i \sigma_f, \partial_i \sigma_b) \quad \text{for } i > 0$$

$$\partial_0(\sigma_f, \sigma_b) = (\tau(\sigma_b) \cdot \partial_0 \sigma_f, \partial_0 \sigma_b)$$

\Rightarrow Only the **0-face** is modified in the **twisted product**.

Reminder about the **EZ-vector field** of $\Delta^5 \times \Delta^3$.



The **vector field** is concerned by faces ∂_i only if $i > 0$.

1. The **twisting function** τ modifies only $\boxed{0}$ -faces.
2. The **EZ-vector field** V_{EZ} of $X \times Y$
uses only \boxed{i} -faces with $i \geq 1$.

$\Rightarrow V_{EZ}$ is **defined** and **admissible** as well on $X \times_{\boxed{\tau}} Y$.

Fundamental theorem of admissible vector fields \Rightarrow

$$\begin{array}{ccc}
 C_*(X \times Y) & & C_*(X \times_{\boxed{\tau}} Y) \\
 V_{EZ} \Rightarrow \Downarrow & & V_{EZ} \Rightarrow \Downarrow \\
 C_*(X) \otimes C_*(Y) & & C_*(X) \otimes_{\boxed{t}} C_*(Y)
 \end{array}$$

Known as the **twisted Eilenberg-Zilber Theorem**.

Corollary: **Base B 1-reduced** \Rightarrow **Algorithm:**

$$[(F, C_*(F), EC_*^F, \varepsilon_F) + (B, C_*(B), EC_*^B, \varepsilon_B) + G + \tau] \\ \longmapsto (F \times_\tau B, C_*(F \times_\tau B), EC_*^{F \times_\tau B}, \varepsilon_{F \times_\tau B}).$$

Version of F with effective homology

+ Version of B with effective homology

+ $G + \tau$ describing the fibration $F \hookrightarrow F \times_\tau B \rightarrow B$

\Rightarrow Version with effective homology of the total space $F \times_\tau B$.

= Version with effective homology

of the Serre Spectral Sequence

Analogous result for the **Eilenberg-Moore spectral sequence**.

Key results:

$G = \text{Simplicial group} \Rightarrow BG = \text{classifying space.}$

$$BG = \dots (((SG \times_{\tau} SG) \times_{\tau} SG) \times_{\tau} SG) \times_{\tau} \dots$$

$X = \text{Simplicial set} \Rightarrow KX = \text{Kan loop space.}$

$$KX = \dots (((S^{-1}X \times_{\tau} S^{-1}X) \times_{\tau} S^{-1}X) \times_{\tau} S^{-1}X) \times_{\tau} \dots$$

Analogous process \Rightarrow **Algorithms:**

$$(G, C_*G, EC_*^G, \varepsilon_G) \mapsto (BG, C_*BG, EC_*^{BG}, \varepsilon_{BG})$$

$$(G, C_*X, EC_*^X, \varepsilon_X) \mapsto (KX, C_*KX, EC_*^{KX}, \varepsilon_{KX})$$

8/8. **Main result.**

“Standard” **Algebraic Topology** becomes **constructive**.

The most basic **standard spectral sequences**

(**Serre, Eilenberg-Moore**) become **effective**.

Example: $\pi_5(\Omega S^3 \cup_2 D^3) = (\mathbb{Z}/2)^4$.

$$\#[S^5, \text{Cont}(S^1, S^3) \cup_2 D^3] = 16$$

Computing plan: Five fibrations (twisted products):

$$X_2 = \Omega S^3 \cup_2 D^3$$

$$K(\mathbb{Z}/2, 1) \rightarrow X_3 \rightarrow X_2 \quad \pi_2(X_2) = \mathbb{Z}/2$$

$$K(\mathbb{Z}/2, 2) \rightarrow X_4 \rightarrow X_3 \quad \pi_3(X_3) = \mathbb{Z}/2$$

$$K(\mathbb{Z}, 3) \rightarrow X'_4 \rightarrow X_4 \quad \pi_4(X_4) = \mathbb{Z}/4 \oplus \mathbb{Z}$$

$$K(\mathbb{Z}/2, 3) \rightarrow X''_4 \rightarrow X'_4 \quad \pi_4(X'_4) = \mathbb{Z}/4$$

$$K(\mathbb{Z}/2, 3) \rightarrow X_5 \rightarrow X''_4 \quad \pi_4(X''_4) = \mathbb{Z}/2$$

Applying five times our effective version

of the Serre spectral sequence and

five times the effective version

of the Eilenberg-Moore spectral sequence

$$\Rightarrow \pi_5(X_2) = \pi_5(X_5) = H_5(X_5) = (\mathbb{Z}/2)^4.$$

The END

```
;; Clock  
Computing  
<TnPr <TnPr  
End of computing.  
  
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Computing the boundary of the generator 19 (dimension 7) :  
<TnPr <TnPr <TnPr S3 <<Abar[2 S1][2 S1]>>> <<Abar>>> <<Abar>>>  
End of computing.
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Homology in dimension 6 :

Component Z/12Z

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