

Effective Moore-Postnikov Factorization

```
;; Clock  
Computing  
<TnPr <TnPr  
End of computing.  
  
;; Clock -> 2002-01-17, 19h 25m 36s.  
Computing the boundary of the generator 19 (dimension 7) :  
<TnPr <TnPr <TnPr S3 <<Abar[2 S1][2 S1]>>> <<Abar>>> <<Abar>>>  
End of computing.
```

Homology in dimension 6 :

Component Z/12Z

---done---

```
;; Clock -> 2002-01-17, 19h 27m 15s
```

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Semantics of colours:

Blue = “Standard” Mathematics

Red = Constructive, effective,

algorithm, machine object, ...

Violet = Problem, difficulty,

obstacle, disadvantage, ...

Green = Solution, essential point,

mathematicians, ...

Dark Orange = Fuzzy objects.

Pale grey = Hyper-Fuzzy objects.

Plan.

1. Eilenberg-MacLane spaces.
2. Map \Rightarrow Mapping Cone.
3. Postnikov's fundamental Theorem.
4. Postnikov construction
for spaces with effective homology.
5. Postnikov construction for square diagrams.

1/5. Eilenberg-MacLane spaces.

Given: $\pi =$ Abelian group ; $2 \leq n \in \mathbb{N}$.

$K(\pi, n) :=$ CW-complex well-defined up to homotopy by:

$$\pi_n(K(\pi, n)) = \pi \quad \text{and} \quad \pi_k(K(\pi, n)) = 0 \quad \text{for } k \neq n$$

$K(\pi, n)$ represents the functor:

$$\underline{\text{Top}} \ni X \longmapsto H^n(X; \pi) \in \underline{\text{Ab}}$$

Characteristic class: $\chi_{\pi, n} \in H^n(K(\pi, n), \pi)$.

Canonical bijection:

$$[X, K(\pi, n)] \ni f \xrightarrow{\cong} f^*(\chi_{\pi, n}) \in H^n(X; \pi)$$

Canonical simplicial model of $K(\pi, n)$

by Eilenberg and MacLane.

$$E(\pi, n)_k := C^n(\Delta^k, \pi)$$

$$K(\pi, n)_k := Z^n(\Delta^k, \pi)$$

If $\alpha : [0 \dots k] \rightarrow [0 \dots \ell] = \Delta^k \rightarrow \Delta^\ell$ and $s \in E(\pi, n)_\ell$, then:

$$\alpha^*(s) := \alpha^*(s)$$

$$\alpha : 0 \dots k \rightarrow 0 \dots \ell \Rightarrow \alpha^* : E(\pi, n)_\ell \rightarrow E(\pi, n)_k$$

$$\alpha : \Delta^k \rightarrow \Delta^\ell \Rightarrow \alpha^* : C^n(\Delta^\ell, \pi) \rightarrow C^n(\Delta^k, \pi)$$

Normalized cochains (and cocycles):

$$C^n(\Delta^k, \pi) := \{c : \Delta_n^k \rightarrow \pi\}$$

$$\Delta_n^k := \{0 \leq i_0 \leq \dots \leq i_n \leq k\}$$

$$\Delta_{n,ND}^k := \{0 \leq i_0 < \dots < i_n \leq k\} \quad (= \emptyset \text{ if } k < n)$$

$$\Delta_{n,D}^k := \Delta_n^k - \Delta_{n,ND}^k$$

Cochain c **normalized** := $c|_{\Delta_{n,D}^k} \equiv 0$

$$\Rightarrow [(k < n) \Rightarrow C^n(\Delta^k, \pi) = 0 = Z^n(\Delta_k, \pi)]$$

\Rightarrow **Eilenberg-MacLane's** $K(\pi, n)$ begins only in dimension n .

$$K(\pi, n)_k = \{*_k\} = \{0 \in Z^n(\Delta^k, \pi)\} \text{ for } k < n$$

$$K(\pi, n)_n = \pi$$

Canonical identifications:

Given: $2 \leq n \in \mathbb{N}$, $\pi \in \underline{\text{Ab}}$, $X = \text{simplicial set}$.

$\underline{\text{SSet}}$:= Category of simplicial sets.

$\underline{\text{SSet}}(X, Y) := \{\text{Simplicial maps : } X \rightarrow Y\}$.

Then:

$$C^n(X, \pi) \cong \underline{\text{SSet}}(X, E^n(\pi, n))$$

$$Z^n(X, \pi) \cong \underline{\text{SSet}}(X, Z^n(\pi, n))$$

Proof: abstract nonsense.

Universal characteristic objects:

$$\underline{\text{SSet}}(X, E^n(\pi, n)) \leftrightarrow C^n(X, \pi)$$

$$\begin{aligned} \underline{\text{SSet}}(E(\pi, n-1), E(\pi, n-1)) \ni \text{id} &\leftrightarrow \\ &\leftrightarrow \chi_{E, \pi, n-1} \in C^{n-1}(E(\pi, n-1), \pi) \end{aligned}$$

$\Rightarrow \chi_{E, \pi, n-1} = \text{Universal characteristic cochain.}$

$$\underline{\text{SSet}}(X, Z^n(\pi, n)) \leftrightarrow Z^n(X, \pi)$$

$$\underline{\text{SSet}}(K(\pi, n), K(\pi, n)) \ni \text{id} \leftrightarrow \chi_{K, \pi, n} \in Z^n(K(\pi, n), \pi)$$

$\chi_{K, \pi, n} = \text{Universal characteristic cocycle.}$

Canonical Eilenberg-MacLane Kan-fibration:

$$K(\pi, n-1) \hookrightarrow E(\pi, n-1) \xrightarrow{p} K(\pi, n)$$

$$p : E(\pi, n-1)_k = C^{n-1}(\Delta^k, \pi) \ni c \mapsto dc \in Z^n(\Delta^k, \pi) = K(\pi, n)_k$$

Δ^k contractible

\Rightarrow every cocycle is a coboundary

\Rightarrow short exact sequence:

$$0 \rightarrow Z^{n-1}(\Delta^k, \pi) \xrightarrow{i} C^{n-1}(\Delta^k, \pi) \xrightarrow{d} Z^n(\Delta^k, \pi) \rightarrow 0$$

$\Rightarrow p := d : E(\pi, n-1) \rightarrow K(\pi, n)$ surjective.

In a diagram:

$$\begin{array}{ccc}
 & E(\pi, n-1) & \chi_E \in C^{n-1}(E(\pi, n-1), \pi) \\
 & \uparrow F & \downarrow p \\
 X & \xrightarrow{f} K(\pi, n) & \chi_K \in Z^n(K(\pi, n), \pi)
 \end{array}$$

$$\underline{\text{SSet}}(X, K(\pi, n)) \ni f \xrightarrow{\cong} f^*(\chi_K) \in Z^n(X, \pi)$$

$$\underline{\text{SSet}}(E(\pi, n-1), E(\pi, n-1)) \ni \text{id} \longleftrightarrow \chi_E \in C^{n-1}(E(\pi, n-1))$$

$$\underline{\text{SSet}}(E(\pi, n-1), K(\pi, n)) \ni p \longleftrightarrow p^*(\chi_K) \in Z^n(E(\pi, n-1), \pi)$$

$$\boxed{p^*(\chi_K) = d(\chi_E)}$$

$$f^*(\chi_K) = F^*p^*(\chi_K) = F^*d(\chi_E) = d(F^*(\chi_E))$$

Theorem: In the diagram:

$$\begin{array}{ccc}
 & E(\pi, n-1) & \chi_E \in C^{n-1}(E(\pi, n-1), \pi) \\
 & \downarrow p & \\
 X & \xrightarrow{f} K(\pi, n) & \chi_K \in Z^n(K(\pi, n), \pi) \\
 & \nearrow F ?? &
 \end{array}$$

a **lifting** F of f along p can be defined

if and only if $f^*(\chi_K) \in Z^n(X, \pi)$ is a **coboundary**:

$$f^*(\chi_K) = dc \text{ with } c \in C^{n-1}(X, \pi).$$

$$C^{n-1}(X, \pi) \ni c \leftrightarrow F \in \underline{\text{SSet}}(X, E(n-1, \pi))$$

Cocycle \Rightarrow **Principal Fibration:**

Given: $X \in \mathbf{SSet}$, $z \in Z^n(X; \pi)$.

\Rightarrow **Principal Fibration:**

$$K(\pi, n-1) \hookrightarrow X \underset{z}{\overset{\sim}{\times}}_p E(\pi, n-1) \xrightarrow{p'} X$$

$$\begin{array}{ccc}
 \boxed{X \underset{z}{\overset{\sim}{\times}}_p E(\pi, n-1)} & \xrightarrow{c'} & E(\pi, n-1) \\
 \begin{array}{c} p' \downarrow \\ X \end{array} & & \begin{array}{c} \downarrow p \\ K(\pi, n) \end{array} \\
 & \xrightarrow{z} &
 \end{array}$$

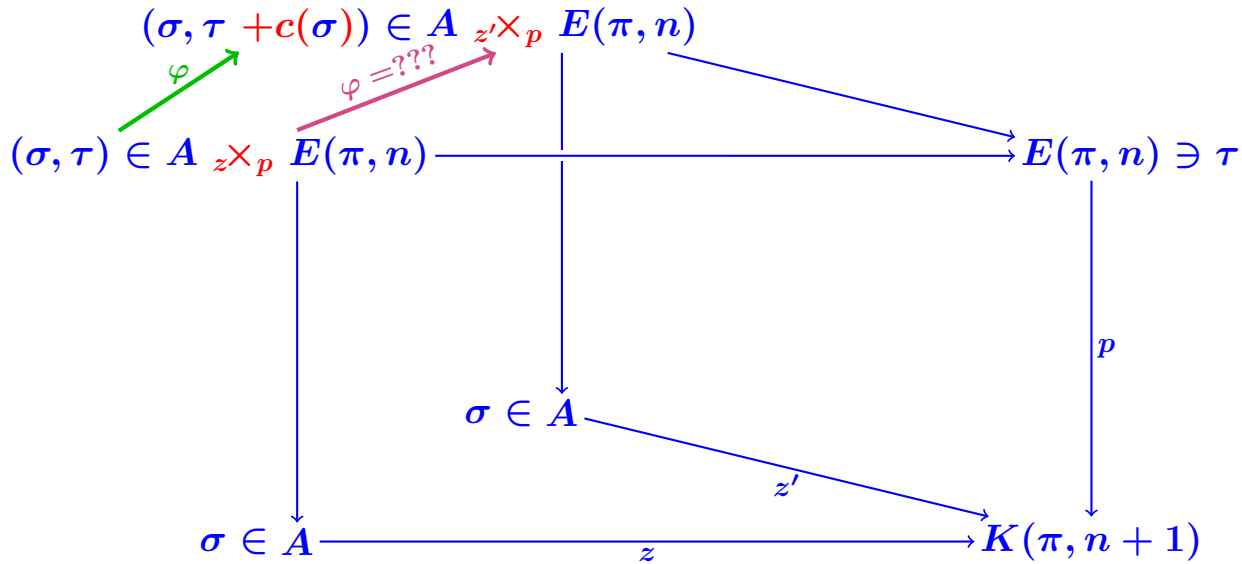
with $[X \underset{z}{\times}_p E(\pi, n-1)]_k :=$

$$\{(a, b) \in X_k \times E(\pi, n-1)_k \text{ st } z(a) = p(b)\}$$

Proposition: Given: $A =$ simplicial set;
 $\pi =$ abelian group;
 $z, z' \in Z^{n+1}(A, \pi)$.

Then the $K(\pi, n)$ -principal fibrations defined by z and z'
 are isomorphic if and only if z and z' are cohomologous.

A cochain c satisfying $dc = z' - z$
 canonically defines such an isomorphism.



Solution: $\varphi(\sigma, \tau) = (\sigma, \tau + c(\sigma))$.

$$(\sigma, \tau) \in z \times_p \Leftrightarrow z(\sigma) = p(\tau)$$

$$(\sigma, \tau + c(\sigma)) \stackrel{?}{\in} z' \times_p \Leftrightarrow z'(\sigma) \stackrel{?}{=} p(\tau + c(\sigma))$$

But $z'(\sigma) = z(\sigma) + dc(\sigma) = p(\tau) + p(c(\sigma)) \Rightarrow \text{OK.} \quad \text{QED}$

Generalized lifting problem.

$$\begin{array}{ccccc}
 X & \xrightarrow{z} & \tilde{X}_p & \xrightarrow{c'} & E(\pi, n-1) \\
 \uparrow F \text{ ??} & & \downarrow p' & & \uparrow p \\
 A & \xrightarrow{f} & X & \xrightarrow{z} & K(\pi, n) \ni \chi_K
 \end{array}$$

Theorem: A **lifting** F of f along p' can be defined

if and only if $f^*z^*\chi_K$ is a **coboundary** $f^*z^*\chi_K = dc$.

Proof: $c \in C^{n-1}(A, \pi) = \underline{\text{SSet}}(A, E(\pi, n-1))$.

Take:

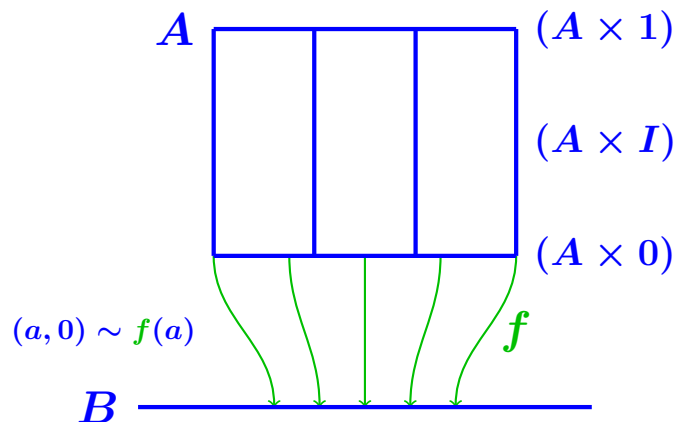
$$F := (f, c)$$

QED

2/5. Map \Rightarrow Mapping Cone.

Mapping cylinder of $f : A \rightarrow B$:

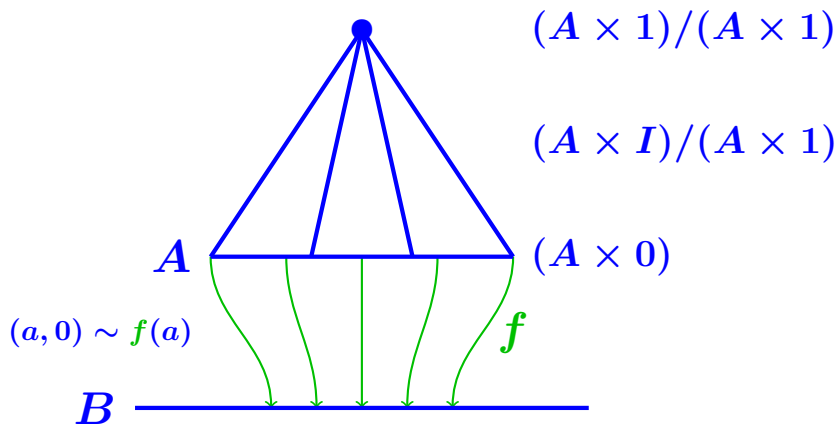
$$\text{Cyl}(f) := [(A \times I) \amalg B] / [(A, 0) \sim f(a)].$$



The canonical projection $\text{Cyl}(f) \rightarrow B$
is a homotopy equivalence.

Mapping cone of $f : A \rightarrow B$:

$$\text{Cone}(f) = [B \amalg (A \times I)] / [(A \times 1) \ \& \ (a, 0) \sim f(a)].$$



$$\text{Cone}(f) := \text{Cyl}(f)/(A \times 1)$$

Algebraic translation.

Definition: Given:

$$f : C'_* \rightarrow C_*$$

= chain complex morphism.

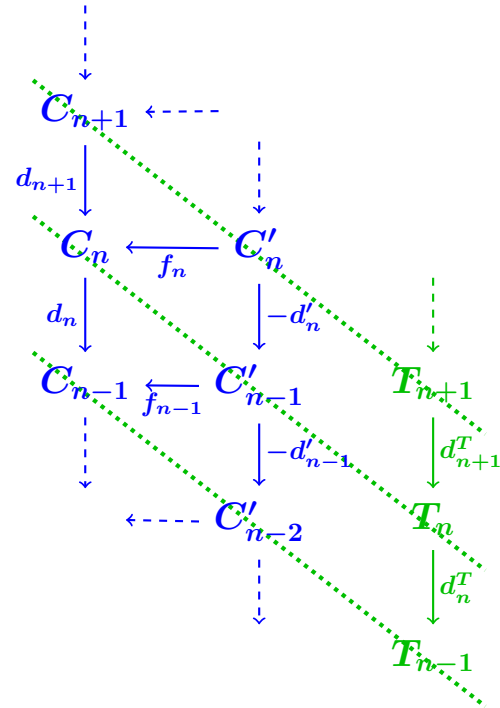
Then $\text{Cone}(f)$ is

the chain complex T_*

defined by:

$$T_n := C_n \oplus C'_{n-1}$$

$$d_n^T : T_n \rightarrow T_{n-1} := \begin{bmatrix} d_n & f_{n-1} \\ 0 & -d'_{n-1} \end{bmatrix}$$



Notation: $\text{Cone}(f) = C_* \oplus_f C_*'^{[1]}$

Remarks:

C_* subcomplex of

$$\text{Cone}(f) = C_* \oplus_f C_*'^{[1]}$$

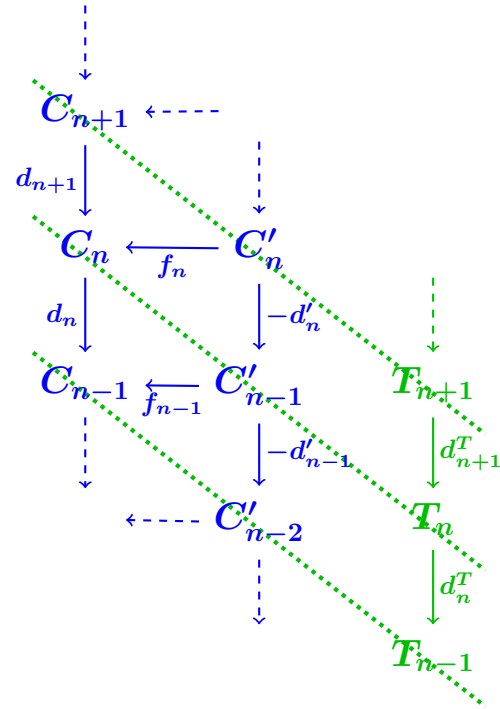
$C_*'^{[1]}$ not subcomplex of

$$\text{Cone}(f) = C_* \oplus_f C_*'^{[1]}$$

$$C_*'^{[1]} = \{C_* \oplus_f C_*'^{[1]}\} / C_*$$

\Leftrightarrow short exact sequence:

$$0 \rightarrow C_* \xrightarrow{i} \text{Cone}(f) \xrightarrow{p} C_*'^{[1]} \rightarrow 0$$



⇒ Long homology exact sequence:

$$\begin{array}{ccccccc}
 & & & & & \dashrightarrow & H_{n+1}(C'_*[1]) \\
 & & & & & \searrow & \curvearrowright \\
 & & & & & \rightarrow & H_n(C_*) \longrightarrow H_n(\text{Cone}(f)) \longrightarrow H_n(C'_*[1]) \\
 & & & & & \searrow & \curvearrowright \\
 & & & & & \dashrightarrow & H_{n-1}(C_*) \dashrightarrow
 \end{array}$$

Better viewed as:

$$\begin{array}{ccccccc}
 & & & & & \dashrightarrow & H_{n+1}(\text{Cone}(f)) \\
 & & & & & \searrow & \curvearrowright \\
 & & & & & \xrightarrow{p} & H_n(C'_*) \xrightarrow[f_*]{p} H_n(C_*) \xrightarrow{i} H_n(\text{Cone}(f)) \\
 & & & & & \searrow & \curvearrowright \\
 & & & & & \dashrightarrow & H_{n-1}(C'_*) \dashrightarrow
 \end{array}$$

Also $\text{Cone}(f) := \text{Cyl}(f)/(A \times 1 = A)$

$$\text{and } \widetilde{H}_*(\text{Cone}(f)) = H_*(\text{Cyl}(f), A)$$

Hurewicz Theorem \Rightarrow

$$\begin{aligned} f \text{ is } n\text{-connected} &:= \pi_i(\text{Cyl}(f), A) = 0 \text{ for } 0 \leq i \leq n \\ &\Leftrightarrow H_i(\text{Cyl}(f), A) = 0 \text{ for } 0 \leq i \leq n \\ &\Leftrightarrow \widetilde{H}_i(\text{Cone}(f)) = 0 \text{ for } 0 \leq i \leq n \\ &\Leftrightarrow \pi_i(\text{Cone}(f)) = 0 \text{ for } 0 \leq i \leq n. \end{aligned}$$

3/5. Postnikov's fundamental theorem.

* * * **Simplicial Set Category** * * *

* * * **All objects assumed 1-connected** * * *

Postnikov Theorem: Given: $f : A \rightarrow B =$

an n -connected map.

Then \exists a **factorization**:

$$\begin{array}{ccc}
 & & B' = B \times_{\tau} K(\pi, n) \\
 & \nearrow f' & \downarrow p \\
 A & \xrightarrow{f} & B
 \end{array}$$

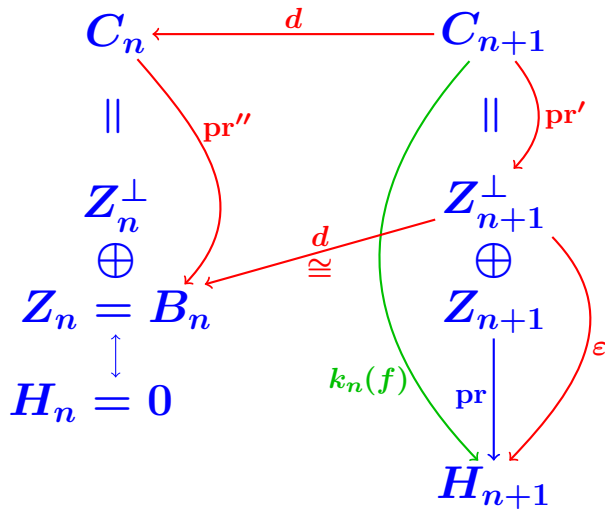
with: f' $(n + 1)$ -connected;

p a principal fibration:

$$K(\pi, n) \rightarrow [B \times_{\tau} K(\pi, n) = B'] \xrightarrow{p} B ;$$

$$\pi = H_{n+1}(\text{Cone}(f)).$$

Lemma: $k_n(f)$ as a cohomology class unambiguously defined.



$\Leftrightarrow \epsilon$ arbitrary

$\Rightarrow \epsilon \cdot pr' \stackrel{??}{=} \text{coboundary } ??$

But $\epsilon \cdot pr' = \epsilon d^{-1} pr'' d$

$$= d(\epsilon d^{-1} pr'').$$

QED

Remark: $H_n = 0 \Rightarrow pr''$ defined.

Postnikov Construction = Diagram of $K(\pi, n)$ -fibrations:

$$\begin{array}{ccccccc}
 E_A & \longrightarrow & E_B & \xrightarrow{\sim} & E_{\text{Cyl}} & \longrightarrow & E_{\text{Cone}} & \longrightarrow & E(\pi, n) \\
 \downarrow & & \downarrow & & \downarrow & & \downarrow & & \downarrow \\
 A & \xrightarrow{f} & B & \xrightarrow{\sim} & \text{Cyl} & \xrightarrow{\text{pr}} & \text{Cone} & \xrightarrow{z} & K(\pi, n+1)
 \end{array}$$

A red arrow points from E_A to A . A green box highlights the top-left part of the diagram, containing E_A , E_B , A , and B , with a green arrow labeled f' pointing from E_A to E_B .

Claim: f n -connected + $z = k_n(f) \Rightarrow f'$ $(n+1)$ -connected.

Postnikov Construction = Functor:

$\{[f : A \rightarrow B], z \in Z^{n+1}(\text{Cone}(f), \pi)\} \mapsto$
 $\mapsto \{\text{Postnikov Diagram}\}$

$$\begin{array}{ccccccc}
 A & \xrightarrow{f} & B & \longrightarrow & \text{Cyl}(f) & \longrightarrow & \text{Cone}(f) \\
 \downarrow & & \downarrow & & \downarrow & & \downarrow \\
 A' & \xrightarrow{f'} & B' & \longrightarrow & \text{Cyl}(f') & \longrightarrow & \text{Cone}(f') \\
 & & & & & & \nearrow^{z'} \\
 & & & & & & K(\pi, n+1)
 \end{array}$$

\searrow_z

+ the $K(\pi, n)$ -fibrations induced by $\{\cdots \rightarrow K(\pi, n+1)\}$.

Postnikov Construction \Rightarrow Diagram of exact sequences:

$$\begin{array}{ccccccc}
 \dashrightarrow \pi_{q+1}(A) & \longrightarrow & \pi_{q+1}(B) & \longrightarrow & \pi_{q+1}(f) & \longrightarrow & \pi_q(A) \dashrightarrow \\
 \downarrow f' & & \downarrow = & & \downarrow \lambda = ?? & & \downarrow f' \\
 \dashrightarrow \pi_{q+1}(E_B) & \longrightarrow & \pi_{q+1}(B) & \longrightarrow & \pi_q(K(\pi, n)) & \longrightarrow & \pi_q(E_B) \dashrightarrow
 \end{array}$$

Top sequence = Exact sequence of **pair** (Cyl, A)

with $\text{Cyl} \sim B$.

Bottom sequence =

Exact sequence of **fibration** $K(\pi, n) \hookrightarrow E_{\text{Cyl}} \rightarrow \text{Cyl}$.

Vertical arrow $\pi_{q+1}(f) \xrightarrow{\lambda = ??} \pi_q(K(\pi, n)) \quad ???$

Vertical arrow $\pi_{q+1}(f) \xrightarrow{??} \pi_q(K(\pi, n))$???

Fibration $E_A \rightarrow A$ **trivial** $\Rightarrow E_A = A \times K(\pi, n) \Rightarrow$

$$0 \longrightarrow \pi_q(K(\pi, n)) \xleftarrow[i]{\rho} \pi_q(E_A) \xleftarrow[p]{\sigma} \pi_q(A) \longrightarrow 0$$

with $\text{id}_{\pi_q(E_A)} = \sigma p + i \rho$.

$$\begin{array}{ccc} \Rightarrow & \pi_{q+1}(E_{\text{Cyl}}, E_A) & \xrightarrow{\partial} \pi_q(E_A) \xrightarrow{\rho} \pi_q(K(\pi, n)) \\ & \downarrow \cong p & \nearrow \lambda = \rho \partial p^{-1} \\ & \pi_{q+1}(f) := \pi_{q+1}(\text{Cyl}, A) & \end{array}$$

Three squares to be proved **commutative**.

$$\text{Square 1} = \begin{array}{ccc} \pi_{q+1}(A) & \xrightarrow{f} & \pi_{q+1}(B) \\ \downarrow f' & \textcircled{??} & \downarrow = \\ \pi_{q+1}(E_B) & \xrightarrow{p} & \pi_{q+1}(B) \end{array}$$

Variant of:

$$\begin{array}{ccc} E_A & \xrightarrow{f} & E_B \\ \sigma \left(\begin{array}{c} \downarrow p \\ \downarrow p \end{array} \right) & \nearrow f' := f\sigma & \downarrow p \\ A & \xrightarrow{f} & B \end{array}$$

Proof:

$$pf' = pf\sigma = fp\sigma = f$$

QED

$$\text{Square 2} = \begin{array}{ccc} \pi_{q+1}(B) & \xrightarrow{j} & \pi_{q+1}(f) \\ \downarrow = & \textcircled{??} & \downarrow \lambda \\ \pi_{q+1}(B) & \xrightarrow{\partial_{\text{fib}}} & \pi_q(K(\pi, n)) \end{array}$$

$$\begin{array}{ccccc} & & \xrightarrow{\partial_{\text{fib}}} & & \\ \pi_{q+1}(B) & \xleftarrow{\cong/p} & \pi_{q+1}(E_B, K(\pi, n)) & \xrightarrow{\partial_{\text{pair}}} & \pi_q(K(\pi, n)) \\ \cong \downarrow & & \cong \downarrow & & = \downarrow \\ \pi_{q+1}(\text{Cyl}) & \xleftarrow{\cong/p} & \pi_{q+1}(E_{\text{Cyl}}, K(\pi, n)) & \xrightarrow{\partial_{\text{pair}}} & \pi_q(K(\pi, n)) \\ j \downarrow & & j \downarrow & \xrightarrow{\lambda} & j \downarrow \uparrow \rho \\ \pi_{q+1}(f) = \pi_{q+1}(\text{Cyl}, A) & \xleftarrow{\cong/p} & \pi_{q+1}(E_{\text{Cyl}}, E_A) & \xrightarrow{\partial_{\text{pair}}} & \pi_q(E_A) \end{array}$$

$$\lambda j = \rho \partial_{\text{pair}} p^{-1} j = \rho \partial_{\text{pair}} j p^{-1} = \rho j \partial_{\text{pair}} p^{-1} = \partial_{\text{pair}} p^{-1} = \partial_{\text{fib}}$$

QED

$$\text{Square 3} = \begin{array}{ccc} \pi_{q+1}(f) & \xrightarrow{\partial_{\text{pair}}} & \pi_q(A) \\ \downarrow \lambda & \textcircled{??} & \downarrow f' \\ \pi_q(K(\pi, n)) & \xrightarrow{i} & \pi_q(E_B) \end{array}$$

$$\begin{array}{ccccc} \pi_{q+1}(\text{Cyl}) & \xleftarrow{\cong} & \pi_{q+1}(E_{\text{Cyl}}, K(\pi, n)) & \xrightarrow{\partial_{\text{pair}}} & \pi_q(K(\pi, n)) \\ \downarrow i & & \downarrow i & \nearrow \lambda & \downarrow \rho \\ \pi_{q+1}(f) = \pi_{q+1}(\text{Cyl}, A) & \xleftarrow{\cong} & \pi_{q+1}(E_{\text{Cyl}}, E_A) & \xrightarrow{\partial_{\text{pair}}} & \pi_q(E_A) & \xrightarrow{f} & \pi_q(E_B) \\ \downarrow \partial_{\text{pair}} & & \nearrow p & \searrow \sigma & \downarrow f & & \downarrow i \cong \\ \pi_q(A) & & & \xrightarrow{f'} & \pi_q(E_{\text{Cyl}}) & & \end{array}$$

$$\begin{aligned} 0 &= f \partial_{\text{pair}} p^{-1} = f(i\rho + \sigma p) \partial_{\text{pair}} p^{-1} = f i \rho \partial_{\text{pair}} p^{-1} + f \sigma p \partial_{\text{pair}} p^{-1} = \\ &= i \rho \partial_{\text{pair}} p^{-1} + f' \partial_{\text{pair}} p p^{-1} = i \lambda + f' \partial_{\text{pair}} = 0 \end{aligned}$$

QED up to sign

Remark: The **Postnikov construction** is a **functor**.

$$\begin{array}{ccc} \text{Canonical morphism : } A & \xrightarrow{f} & B & k_n(f) \in Z^n(\text{Cone}) \\ & \downarrow & \downarrow k'_n(f) & \\ & * & \longrightarrow K(\pi, n+1) & \text{id} \in Z^n(\text{Cone}) \end{array}$$

\Rightarrow Diagram:

$$\begin{array}{ccc} \pi_{n+1}(f) = \pi_{n+1}(\text{Cyl}, A) & \xrightarrow[\cong]{k'} & \pi_{n+1}(K(\pi, n+1), *) \\ \cong \uparrow p^{-1} & & \cong \uparrow p^{-1} \\ \pi_{n+1}(E_{\text{Cyl}}, E_A) & \xrightarrow{k'} & \pi_{n+1}(E(\pi, n), K(\pi, n)) \\ \downarrow \partial_{\text{pair}} & & \cong \downarrow \partial_{\text{pair}} \\ \pi_n(E_A) & \xrightarrow{k'} & \pi_n(K(\pi, n)) \\ \downarrow \rho & & \cong \downarrow \rho \\ \pi_n(K(\pi, n)) & \xrightarrow{=} & \pi_n(K(\pi, n)) \end{array}$$

λ (green arrow) points from $\pi_{n+1}(f)$ to $\pi_n(K(\pi, n))$.

\Rightarrow $\lambda : \pi_{n+1}(f) \rightarrow \pi_n(K(\pi, n)) = \text{isomorphism.}$

Using the morphism of exact sequence – I.

$$\begin{array}{ccccccccc}
 \pi_{n+1}(B) & \longrightarrow & \pi_{n+1}(f) & \longrightarrow & \pi_n(A) & \longrightarrow & \pi_n(B) & \longrightarrow & \cancel{\pi_n(f)} \\
 \downarrow = & & \downarrow \lambda \cong & & \downarrow f' & & \downarrow = & & \downarrow \lambda \\
 \pi_{n+1}(B) & \longrightarrow & \pi_n(K(\pi, n)) & \longrightarrow & \pi_n(E_B) & \longrightarrow & \pi_{n+1}(B) & \longrightarrow & \cancel{\pi_{n-1}(K(\pi, n))}
 \end{array}$$

$\Rightarrow f' : \pi_n(A) \rightarrow \pi_n(E_B) = \text{isomorphism}$

and the same for $f' : \pi_i(A) \rightarrow \pi_i(E_B)$ for $i < n$.

Using the morphism of exact sequence – II.

$$\begin{array}{ccccccc}
 \pi_{n+1}(A) & \longrightarrow & \pi_{n+1}(B) & \longrightarrow & \pi_{n+1}(f) & & \\
 & & \downarrow f' & & \downarrow = & & \downarrow \lambda \cong \\
 \pi_{n+1}(K(\pi, n)) & \longrightarrow & \pi_{n+1}(E_B) & \longrightarrow & \pi_{n+1}(B) & \longrightarrow & \pi_n(K(\pi, n))
 \end{array}$$

Standard diagram chasing \Rightarrow

$f' : \pi_{n+1}(A) \rightarrow \pi_{n+1}(E_B)$ is epimorphism.

Proved: $f' : A \rightarrow E_B$ is $(n + 1)$ -connected.

QED

Corollary: Given: A, B simply connected and $f : A \rightarrow B$.

Then:

Moore-Postnikov decomposition

of f with:

- f_i is i -connected.
- p_i is the fibration:

$$\begin{array}{ccc} K(\pi_{i+1}(f_i), i) & \hookrightarrow & B_{i+1} \\ & & \downarrow p_i \\ & & B_i \end{array}$$

$$\begin{array}{ccccc} \varprojlim B_i = B_\infty & & & & \\ \longleftarrow & & & & \\ & & & & \downarrow \text{---} \\ & & & & B_3 \\ & & & & \downarrow p_2 \\ & & & & B_2 \\ & & & & \downarrow p_1 \\ & & & & B \\ & & & & =: B_1 \\ & & & & \\ A & \xrightarrow{f =: f_1} & B & & \\ & \nearrow f_2 & & & \\ & \nearrow f_3 & & & \\ & \nearrow f_\infty & & & \end{array}$$

defined by the **Postnikov class** $k_i(f_i) \in H^{i+1}(B_i, \pi_{i+1}(f_i))$.

In particular f_∞ is a **homotopy equivalence**.

Particular case 1. $f : A \rightarrow *$ \Rightarrow Postnikov tower of A .

$$\begin{array}{c}
 \lim_{\leftarrow} A_i = A_\infty \\
 \downarrow \text{dashed} \\
 A_3 = K(\pi_3, 3) \times_{k_3} A_2 \\
 \downarrow p_2 \\
 A_2 = K(\pi_2, 2) \\
 \downarrow p_1 \\
 A \xrightarrow{f_1 := *} *
 \end{array}$$

The diagram shows a Postnikov tower starting from a space A at the bottom. A horizontal arrow labeled $f_1 := *$ points from A to a point $*$. From A , three diagonal arrows labeled f_2 , f_3 , and f_∞ point upwards to the stages A_2 , A_3 , and A_∞ respectively. Vertical arrows labeled p_1 and p_2 point downwards from A_2 to $*$ and from A_3 to A_2 . A dashed vertical arrow points downwards from A_∞ to A_3 . The stages are defined as $A_2 = K(\pi_2, 2)$ and $A_3 = K(\pi_3, 3) \times_{k_3} A_2$.

$A_{i-1} = (i-1)$ -th stage of the Postnikov tower.

$$\pi_i(A) = H_{i+1}(\text{Cone}(f_{i-1}), \mathbb{Z}).$$

$k_i \in H^{i+1}(A_{i-1}, \pi_i(A)) = \text{Postnikov class} = \text{intrinsic}.$

$k_i \in H^{i+1}(A_{i-1}, \pi_i) \text{ depending on } \pi_i \cong \pi_i(A).$

$$A_i = K(\pi_i(A), i) \times_{k_i} A_{i-1}.$$

General case:

$$\begin{array}{ccc}
 \lim B_i := B_\infty & \longleftarrow & F_\infty := \lim F_i \\
 \swarrow f_\infty & & \downarrow \\
 K(\pi_3(\text{Cone}(f_2)), 2) \times_{k_3} B_2 := B_3 & \longleftarrow & F_2 = K(\pi_3(\text{Cone}(f_2)), 2) \times_{k_3} F_1 \\
 \swarrow f_3 & & \downarrow p_2 \\
 K(\pi_2(\text{Cone}(f_1)), 1) \times_{k_2} B_1 := B_2 & \longleftarrow & F_1 = K(\pi_2(\text{Cone}(f_1)), 1) \\
 \swarrow f_2 & & \downarrow p_1 \\
 A \xrightarrow{f} B_1 := B & \longleftarrow & *
 \end{array}$$

The restriction of the Postnikov tower of f above B

to the base point $*$

is the Postnikov tower of the homotopy fiber F of $f : A \rightarrow B$.

Particular case of the general case:

$$\begin{array}{ccc}
 \lim_{\leftarrow} B_i =: B_\infty & \hookrightarrow & F_\infty := \lim_{\leftarrow} F_i \\
 \uparrow f_\infty & \dashv & \downarrow \\
 K(\pi_3(B), 2) \times_{k_3} B_2 =: B_3 & \hookrightarrow & F_2 = K(\pi_3(B), 2) \times_{k_3} F_1 \\
 \uparrow f_3 & \downarrow p_2 & \downarrow p_2 \\
 K(\pi_2(B), 1) \times_{k_2} B_1 =: B_2 & \hookrightarrow & F_1 = K(\pi_2(B), 1) \\
 \uparrow f_2 & \downarrow p_1 & \downarrow p_1 \\
 * & \xrightarrow{f} & B_1 := B \hookrightarrow *
 \end{array}$$

$F_\infty = \text{Homotopy fiber of } [* \rightarrow B] =: \Omega B = \text{loop space of } B.$

Corollary: The **Postnikov classes** of ΩB

are the **Whitehead classes** of $B.$

4/5. Postnikov construction

for spaces with effective homology

Given: $f : A \rightarrow B$ an n -connected simplicial map ($n \geq 1$)
between 1-reduced simplicial sets with effective homology.

Then: A general algorithm produces the diagram:

$$\begin{array}{ccc}
 & & B' \\
 & \nearrow f' & \downarrow p \\
 A & \xrightarrow{f} & B
 \end{array}$$

with:

- f' is $(n + 1)$ -connected;
- p is a fibration $K(\pi, n) \hookrightarrow B' \xrightarrow{p} B$;
- B' is a simplicial set with effective homology.

Let C_* be a free \mathbb{Z} -chain complex,

n -connected ($n \geq 1$), with effective homology.

\Rightarrow Equivalence $C_* \iff EC_*$

with $EC_* =$ free \mathbb{Z} -chain complex of finite type.

$\Rightarrow \pi := H_{n+1}(C_*) = H_{n+1}(EC_*)$ is computable.

$$\begin{array}{ccccc}
 EC_n & \xleftarrow{d} & EC_{n+1} & \xleftarrow{d} & EC_{n+2} \\
 \cup & & \cup & & \searrow \\
 0 & \xleftarrow{d} & Z_{n+1} & & \\
 & & \cup & & \\
 & & B_{n+1} & & \\
 & & \searrow & & \\
 & & \pi & &
 \end{array}$$

pr (blue arrow from Z_{n+1} to B_{n+1})
 Eh (red arrow from 0 to π)

$\Rightarrow \text{Eh} =$ fund. cohomology class
of EC_* computable.

Effective homology of C_* :

$$\begin{array}{ccccccc}
 lh \circlearrowleft & C_* & \begin{array}{c} \xleftarrow{lf} \\ \xrightarrow{lg} \end{array} & \widehat{C}_* & \begin{array}{c} \xleftarrow{rg} \\ \xrightarrow{rf} \end{array} & EC_* & \circlearrowright rh \\
 & & \searrow \mathfrak{h} = E\mathfrak{h} \circ rf \circ lg & & & \downarrow E\mathfrak{h} & \\
 & & & & & \pi &
 \end{array}$$

\Rightarrow The fundamental cohomology class of C_* :

$$\mathfrak{h} := E\mathfrak{h} \circ rf \circ lg$$

is computable.

Cone Theorem: Given: $f : A_* \rightarrow B_*$

a chain complex morphism

between two chain complexes

with effective homology.

Then: A general algorithm computes

a version with effective homology of $\text{Cone}(f)$.

Proof: Particular case of the SES_2 Theorem.

Given $f : A \rightarrow B$ n -connected +

$$C_*(A) \xrightarrow{f_*} C_*(B) \xrightarrow{i} \text{Cone}(f_*)$$

+ A and B with effective homology

\Rightarrow

$$C_*(A) \xrightarrow{f_*} C_*(B) \xrightarrow{i} \text{Cone}(f_*) \xrightarrow{\mathfrak{h}} \pi$$

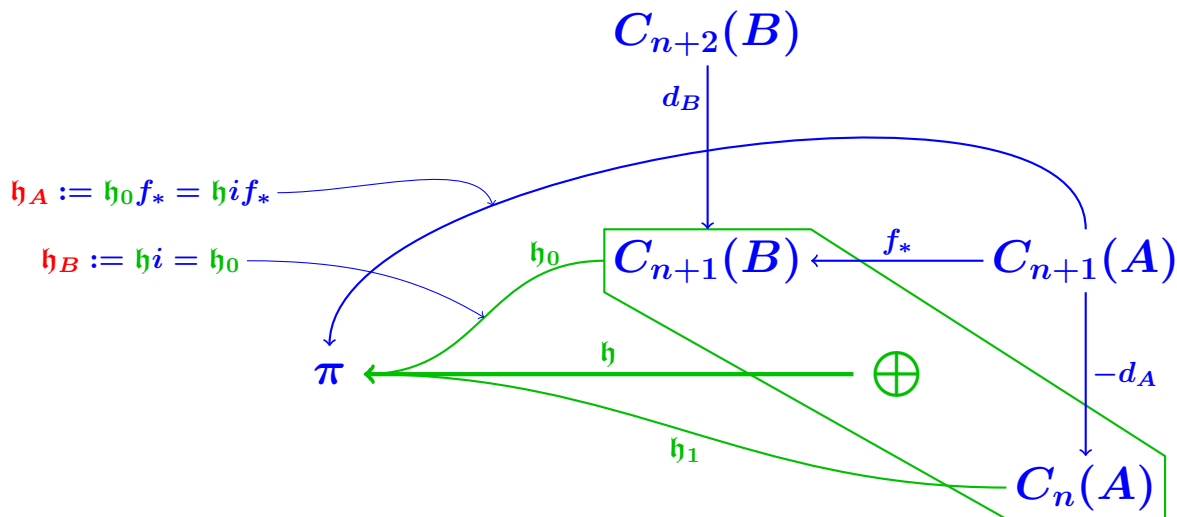
with $\pi := H_{n+1}(\text{Cone}(f_*), \mathbb{Z})$

\Rightarrow π -cochains:

$$\mathfrak{h} \in Z^{n+1}(\text{Cone}(f_*), \pi)$$

$$\mathfrak{h}_B := \mathfrak{h}i \in Z^{n+1}(B, \pi)$$

$$\mathfrak{h}_A := \mathfrak{h}if_* \in Z^{n+1}(A, \pi)$$



$h = [h_0 \quad h_1]$ cocycle \Leftrightarrow

$$[h_0 \quad h_1] \begin{bmatrix} d_B & f_* \\ 0 & -d_A \end{bmatrix} = [h_0 d_B \quad h_0 f_* - h_1 d_A] = 0$$

$$\Rightarrow h_A = h_0 f_* = h_1 d_A = d_A(h_1)$$

$\Rightarrow h_A$ is a cohomology class **constructively null**.

5/5. Postnikov construction for morphisms.

Given: A commutative diagram:

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ \alpha \downarrow & & \downarrow \beta \\ C & \xrightarrow{g} & D \end{array}$$

with A, B, C, D 1-connected, and f, g n -connected.

Then: A **morphism** between **Postnikov constructions**

is induced:

$$\begin{array}{ccccc} & & B' & \xrightarrow{\beta'} & D' \\ & \nearrow f' & \downarrow p & & \downarrow q \\ A & \xrightarrow{f} & B & \xrightarrow{g} & D \\ & \searrow & \nearrow & & \\ & & C & & \end{array}$$

α β

Remember the functor:

$$\{[f : A \rightarrow B], z \in Z^{n+1}(\text{Cone}(f), \pi)\} \longmapsto \\ \longmapsto \{\text{Postnikov Diagram}\}$$

$$\begin{array}{ccccccc} E_A & \longrightarrow & E_B & \xrightarrow{\sim} & E_{\text{Cyl}} & \longrightarrow & E_{\text{Cone}} & \longrightarrow & E(\pi, n) \\ & \searrow & \nearrow & & \downarrow & & \downarrow & & \downarrow \\ & & & & \text{Cyl} & \xrightarrow{\text{pr}} & \text{Cone} & \xrightarrow{z} & K(\pi, n+1) \\ & \swarrow & \searrow & & & & & & \\ A & \xrightarrow{f} & B & \xrightarrow{\sim} & & & & & \end{array}$$

The diagram shows a commutative structure with two rows of spaces. The top row consists of $E_A \rightarrow E_B \xrightarrow{\sim} E_{\text{Cyl}} \rightarrow E_{\text{Cone}} \rightarrow E(\pi, n)$. The bottom row consists of $A \xrightarrow{f} B \xrightarrow{\sim} \text{Cyl} \xrightarrow{\text{pr}} \text{Cone} \xrightarrow{z} K(\pi, n+1)$. Vertical arrows connect $E_A \rightarrow A$, $E_B \rightarrow B$, $E_{\text{Cyl}} \rightarrow \text{Cyl}$, $E_{\text{Cone}} \rightarrow \text{Cone}$, and $E(\pi, n) \rightarrow K(\pi, n+1)$. A red curved arrow points from E_A to A . A green parallelogram highlights the region containing E_A, E_B, A, B and the map f' from E_A to E_B .



Serious general **coherence problem**

with **cocycle coherence**.

Problem visible in this simple case:

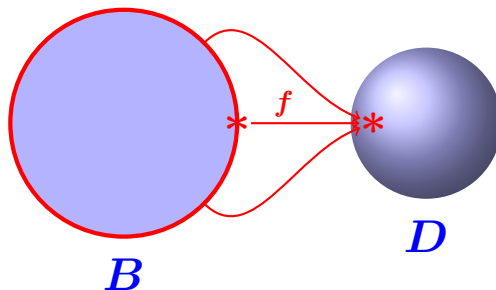
$A = C = * \Rightarrow \text{Cone} = B \text{ or } D.$

$B = \text{Disk}$ bounded by a **circle**.

$$\Rightarrow C_*(B) = [0 \leftarrow \mathbb{Z} \xleftarrow{0} \mathbb{Z} \xleftarrow{1} \mathbb{Z} \leftarrow 0]$$

$D = \text{2-sphere} \Rightarrow C_*(D) = [0 \leftarrow \mathbb{Z} \leftarrow 0 \leftarrow \mathbb{Z} \leftarrow 0]$

$f = \text{the map [circle of } B] \mapsto [\text{base point of } D].$



B , D and f are 1-connected.

$\Rightarrow \pi = H_2(B, \mathbb{Z}) = 0$ and $\pi' = H_2(D, \mathbb{Z}) = \mathbb{Z}$.

$\Rightarrow z = 0 \in Z^2(B, 0)$ and $z' = \text{id} \in Z^2(D, \mathbb{Z})$ [no choice !].

$$\Rightarrow \begin{array}{ccc} B \times_0 K(0, 1) & \xrightarrow{???} & D \times_1 K(\mathbb{Z}, 1) \\ \downarrow p & & \downarrow p \\ B & \xrightarrow{f} & D \end{array}$$

Only one natural intermediate object:

$$\begin{array}{ccccc} B \times_0 K(0, 1) & \xrightarrow{???} & B \times_{?z''?} K(\mathbb{Z}, 1) & \xrightarrow{???} & D \times_1 K(\mathbb{Z}, 1) \\ \downarrow p & & \downarrow p & & \downarrow p \\ B & \xrightarrow{\text{id}} & B & \xrightarrow{f} & D \end{array}$$

with in particular $?z''?$ to be determined.

$$\begin{array}{ccccc}
 B \times_0 K(0, 1) & \xrightarrow{???) & B \times_{?z''?) & K(\mathbb{Z}, 1) & \xrightarrow{???) & D \times_1 K(\mathbb{Z}, 1) \\
 \downarrow p & & \downarrow p & & \downarrow p \\
 B & \xrightarrow{\text{id}} & B & \xrightarrow{f} & D
 \end{array}$$

Left-coherence $\Rightarrow z'' = f_*(z) = f_*(0) = 0$

with $f_* : Z^2(B, \pi) \rightarrow Z^2(B, \pi')$.

Right-coherence $\Rightarrow z'' = f^*(z') = f^*(1) = 1$

with $f^* : Z^2(D, \pi') \rightarrow Z^2(B, \pi)$.

$0 \neq 1 \Rightarrow \text{Problem!}$

At the level of chain complexes:

$$\begin{array}{ccccccc}
 & & & & f_* & & \\
 & & & & \text{Non-comm!} & & \\
 \pi = H_2 = 0 & \xleftarrow{z=0} & C_2 = \mathbb{Z} & \xrightarrow{\frac{1}{f}} & C'_2 = \mathbb{Z} & \xrightarrow{z'=1} & \mathbb{Z} = H'_2 = \pi' \\
 & & \downarrow 1 & & \downarrow 0 & & \\
 & & C_1 = \mathbb{Z} & \xrightarrow{\frac{0}{f}} & C'_1 = 0 & & \\
 & & \downarrow 0 & & \downarrow 0 & & \\
 & & C_0 = \mathbb{Z} & \xrightarrow{\frac{1}{f}} & C'_0 = \mathbb{Z} & &
 \end{array}$$

C_* and C'_* chain complexes, $f : C_* \rightarrow C'_*$ morphism.

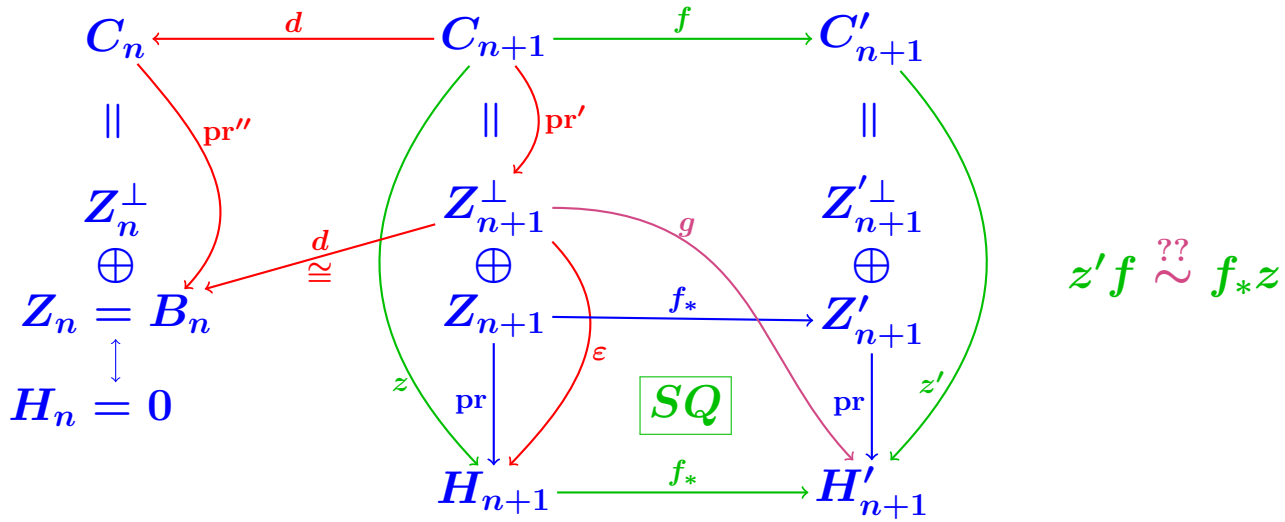
$z \in Z^2(C_*, H_2)$, $z' \in Z^2(C'_*, H'_2) = \text{characteristic cocycles}$.

$f_* : Z^2(C_*, H_2) \rightarrow Z^2(C_*, H'_2) : z = 0 \mapsto 0 = f_*(z)$.

$f^* : Z^2(C'_*, H'_2) \rightarrow Z^2(C_*, H'_2) : z' = 1 \mapsto 1 = f^*(z')$.

Bug: $0 = f_*(z) \neq f^*(z') = 1$!!!

Theorem: In this context, $f_*(z)$ and $f^*(z')$ are **cohomologous**.



1. The square SQ is **commutative**.
2. For g arbitrary:

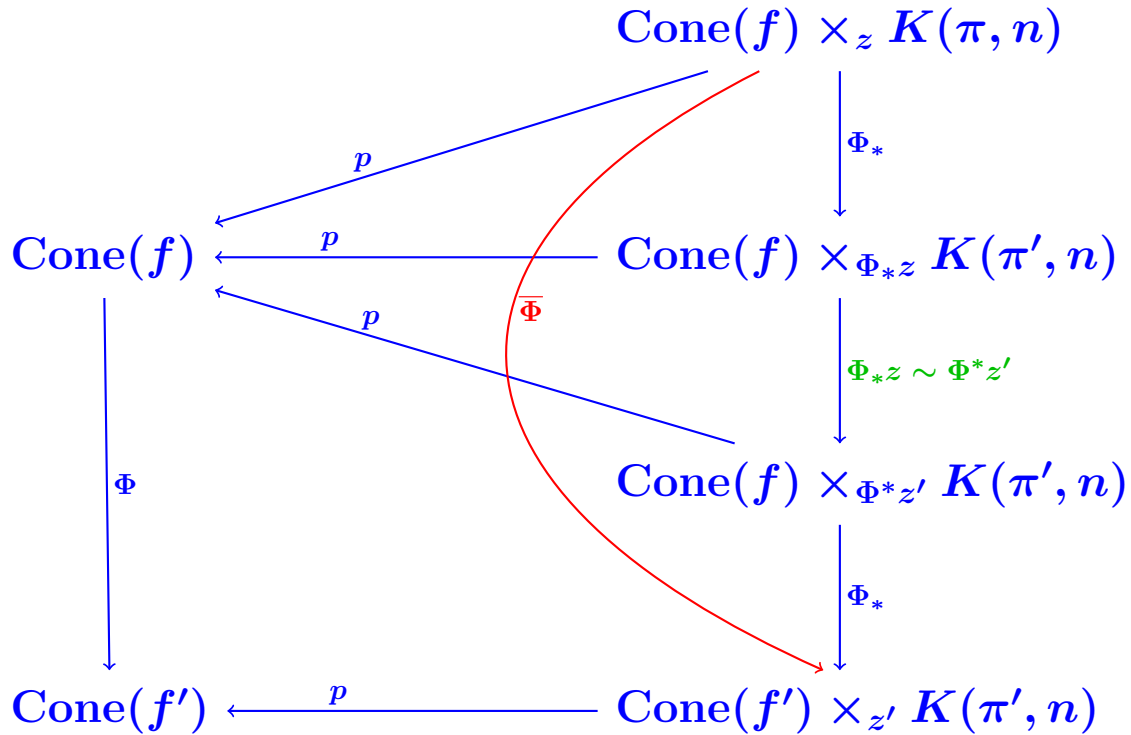
$$g \cdot \text{pr}' = g \cdot d^{-1} \cdot \text{pr}'' \cdot d = d(g \cdot d^{-1} \cdot \text{pr}'')$$

QED

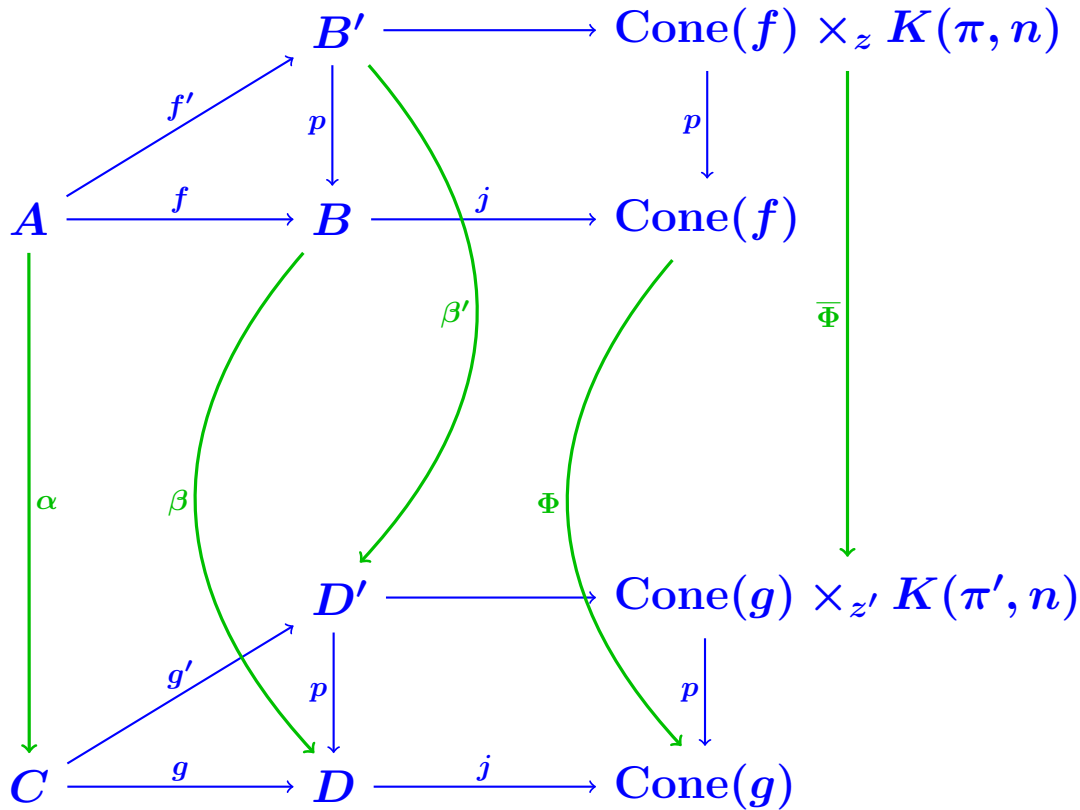
⇒ Diagram-1:

$$\begin{array}{ccccccccc}
 A & \xrightarrow{f} & B & \xrightarrow{j} & \text{Cyl}(f) & \xrightarrow{p} & \text{Cone}(f) & \xrightarrow{z} & K(\pi, n+1) \\
 \downarrow \alpha & \text{Com.} & \downarrow \beta & \text{Com.} & \downarrow \Phi & \text{Com.} & \downarrow \Phi & z'\Phi \sim \Phi_*z & \downarrow \Phi_* \\
 C & \xrightarrow{g} & D & \xrightarrow{j} & \text{Cyl}(g) & \xrightarrow{p} & \text{Cone}(g) & \xrightarrow{z'} & K(\pi', n+1)
 \end{array}$$

⇒ Diagram-2:



⇒ Diagram-3:



QED

Corollary: All the **Postnikov constructions**:

- **Moore-Postnikov factorization**;
- **Postnikov tower**;
- **Whitehead tower**;
- **Postnikov tower** of the **homotopy fiber**;

can be **functorially** organized.

In an **effective way** in context of **effective homology**.

The END

```
;; Clock  
Computing  
<TnPr <TnPr  
End of computing.  
  
;; Clock -> 2002-01-17, 19h 25m 36s.  
Computing the boundary of the generator 19 (dimension 7) :  
<TnPr <TnPr <TnPr S3 <<Abar[2 S1][2 S1]>>> <<Abar>>> <<Abar>>>  
End of computing.  
  
Homology in dimension 6 :  
  
Component Z/12Z  
  
---done---  
;; Clock -> 2002-01-17, 19h 27m 15s
```

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March 21-25, 2011*