Constructive Homology Classes and Constructive Triangulations

;; Cloc Computing <TnPr <Tn End of computing.

;; Clock -> 2002-01-17, 19h 25m 36s. Computing the boundary of the generator 19 (dimension 7) : <TnPr <TnPr <TnPr S3 <<Abar[2 S1][2 S1]>>> <<Abar>>> <<Abar>>> End of computing.

Homology in dimension 6 :

Component Z/12Z

---done---

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Dedicated to <u>Mirian Andrès</u>

Francis Sergeraert, Institut Fourier, Grenoble Mathematics Algorithms and Proofs Logroño, Spain, 8-12 November, 2010 Semantics of colours:

Blue = "Standard" Mathematics Red = Constructive, effective, algorithm, machine object, ... Violet = Problem, difficulty, obstacle, disadvantage, ... Green = Solution, essential point, mathematicians, ... Dark Orange = Fuzzy objects.

Pale grey = Hyper-Fuzzy objects.

<u>Plan</u>.

- 1. Constructive Homological Algebra.
- 2. Triangulations and fundamental cycles.
- 3. Complex projective spaces.
- 4. Connection $P^n \mathbb{C} \iff P^\infty \mathbb{C}$.
- 5. Kenzo program + Constructive Homological Algebra \Rightarrow Constructive Triangulation of $P^n\mathbb{C}$.

1/5. Constructive Homological Algebra.

General style of Homological Algebra:

First step in the classification of angiosperms: Number of cotyledons = 1 or 2.

 $n = 1 \Rightarrow$ Monocotyledons (~ 60.000 species).

 $n = 2 \Rightarrow$ Dicotyledons (~ 200.000 species)

First step in the classification of topological spaces: $(\forall X \in \underline{\mathrm{Top}}) \Rightarrow [(\forall d \in \mathbb{N}) \Rightarrow H_d(X) \in \underline{\mathrm{AbGroup}}].$

Only partial classification !!!

Main problem:

Let $\Phi : \text{Top} \times \text{Top} \to \text{Top}$ be a constructor.

Example: $\Phi(X, Y) := X \times Y$.

Homological version of this constructor ??

$$\Phi_H: (H_*(X), H_*(Y)) \stackrel{???}{\longmapsto} H_*(\Phi(X, Y))$$

Sometimes possible, for example for the product constructor (Künneth formulas).



Example:

The loop space constructor $\Omega: X \mapsto \Omega X := \mathcal{C}(S^1, X)$

Example²:

 $X=S^2 \lor S^4$ $Y=P^2 \mathbb{C}$ $H_*(X)=H_*(Y)=(\mathbb{Z},0,\mathbb{Z},0,\mathbb{Z},0,0,0,\ldots)$

 $egin{aligned} H_*(\Omega X) &= \ (\mathbb{Z}, \ \mathbb{Z}, \ \mathbb{Z}, \ \mathbb{Z}^2, \ \mathbb{Z}^3, \ \mathbb{Z}^4, \ \mathbb{Z}^6, \ \mathbb{Z}^9, \ \mathbb{Z}^{13}, \ \ldots) \ H_*(\Omega Y) &= \ (\mathbb{Z}, \ \mathbb{Z}, \ 0, \ 0, \ \mathbb{Z}, \ \mathbb{Z}, \ 0, \ 0, \ \mathbb{Z}, \ \ldots) \end{aligned}$

Corollary: $\not\exists$ algorithm $\Omega_H : H_*(X) \mapsto H_*(\Omega X)$.

Analysis of the problem.

Ordinary homological algebra is not constructive.

 $H_4(X)$ "=" \mathbb{Z} means: \exists isomorphism $H_4(X) \xleftarrow{\cong} \mathbb{Z}$; But most often \exists is not constructive.

Reorganizing Homological Algebra

to make these \exists 's constructive

- \Rightarrow Constructive Homological Algebra
- \Rightarrow Algorithms:

 $\Phi_{CH}: (CH_*(X), CH_*(Y)) \mapsto CH_*(\Phi(X,Y)).$

 \Rightarrow (JR) Efficient solution of Adams' problem for loop spaces.

2/5. Triangulations and fundamental cycles.

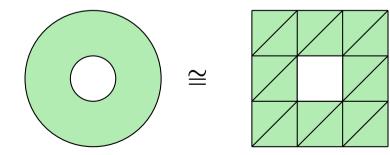
Amazing spin-off of Constructive Homological Algebra:

Using constructive isomorphisms

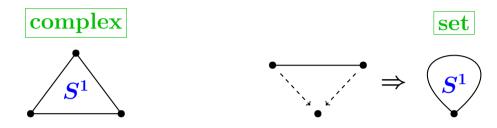
to produce difficult triangulations.

Notions of triangulation.

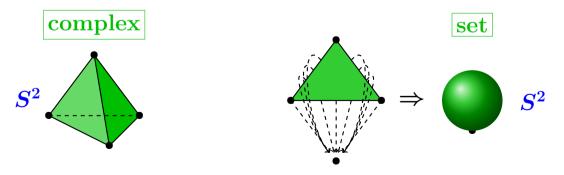
Triangulation as a simplicial complex of $S^1 \times I$.



Triangulations of S^1 as simplicial:



Triangulations of S^2 as simplicial:

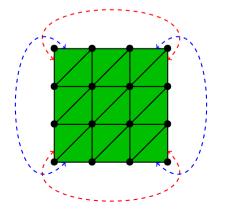


- $M = \text{closed } n \text{-manifold} \Rightarrow M \text{ is triangulable.}$
- We assume M orientable.
- Let \mathcal{T} be some triangulation

and T_n the corresponding collection of *n*-simplices. Then $H_n(M) = \mathbb{Z}$

and a cycle representing a generator of H_n is $z = \sum_{\tau \in T_n} \varepsilon_{\tau} \tau$.

Example for M = 2-torus:

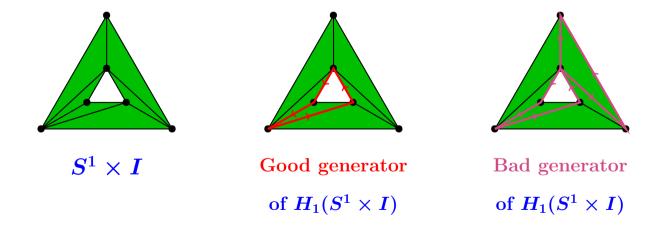


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In a context of Constructive Homological Algebra, the result can <u>sometimes</u> be <u>reversed</u>.

Toy example with $S^1 \times I \stackrel{H}{\sim} S^1$.

$$H_*(S^1 imes I)=H_*(S^1)=(\mathbb{Z},\mathbb{Z},0,0,0,\ldots)$$



3/5. Complex projective spaces.

Using this method to construct a triangulation of $\mathbb{P}^{n}\mathbb{C}$.

 $S^{2n+1} =$ unit sphere(\mathbb{C}^{n+1})

 $P^n\mathbb{C}:=S^{2n+1}/S^1$

 $S^1 \subset S^3 \subset S^5 \subset \cdots \subset S^\infty$

 $* \subset P^1 \mathbb{C} \subset P^2 \mathbb{C} \subset P^3 \mathbb{C} \subset \cdots \subset P^\infty \mathbb{C}$

Universal fibration:

$$K(\mathbb{Z},1)=S^1 \hookrightarrow S^\infty
ightarrow P^\infty \mathbb{C}$$

 $\Rightarrow P^{\infty}\mathbb{C} = K(\mathbb{Z},2)$

 $egin{aligned} K(\mathbb{Z},2) &= ext{``catalog'' space} &= \ ext{collection of all the possible configurations} \ ext{ of elements } z \in H^2(-,\mathbb{Z}) \end{aligned}$

Standard simplicial model for $K(\mathbb{Z},2)$

due to Eilenberg-MacLane.

 $K(\mathbb{Z},2) = ext{Monster:} \ K(\mathbb{Z},2)_n \sim \mathbb{Z}^{rac{n(n-1)}{2}}$

But the methods of Constructive Algebraic Topology can handle this monster. 4/5. Connection $P^nC \longleftrightarrow P^{\infty}\mathbb{C}$.

 $P^{\infty}\mathbb{C} = \lim_{n \to \infty} P^n\mathbb{C}$ has a good homological translation:

$$H_{*}(P^{\infty}\mathbb{C}) = (\overset{0}{\mathbb{Z}}, 0, \overset{2}{\mathbb{Z}}, 0, \overset{4}{\mathbb{Z}}, 0, \overset{6}{\mathbb{Z}}, 0, \overset{8}{\mathbb{Z}}, 0, \overset{10}{\mathbb{Z}}, 0, \dots)$$

$$H_{*}(P^{1}\mathbb{C}) = (\mathbb{Z}, 0, \mathbb{Z}, 0, 0, 0, 0, 0, 0, 0, 0, 0, \dots)$$

$$H_{*}(P^{2}\mathbb{C}) = (\mathbb{Z}, 0, \mathbb{Z}, 0, \mathbb{Z}, 0, 0, 0, 0, 0, 0, 0, \dots)$$

$$H_{*}(P^{3}\mathbb{C}) = (\mathbb{Z}, 0, \mathbb{Z}, 0, \mathbb{Z}, 0, \mathbb{Z}, 0, 0, 0, 0, 0, \dots)$$

$$\cdots = \cdots$$

Also the inclusion $P^n \mathbb{C} \hookrightarrow P^\infty \mathbb{C}$

induces an inclusion $H_*(P^n\mathbb{C}) \hookrightarrow H_*(P^\infty\mathbb{C})$.

So that a generator g_{2n} of $H_{2n}(P^{\infty}\mathbb{C})$ corresponds to a generator g_{2n} of $H_{2n}(P^{n}\mathbb{C})$ which could correspond to a triangulation of $P^{n}\mathbb{C}$.

5/5. Kenzo calculations.

- 1. kz2 := $K(\mathbb{Z}, 2)$
- 2. "Local" calculations are possible.
- 3. The effective homology is computable:

 $[C_*(K(\mathbb{Z},2))=$ K86] \ll K216 \Longrightarrow K212

- 4. G4 = generator of H_4 (K212) = \mathbb{Z} .
- 5. GP4 = generator of $H_4(K86) = H_4(K(\mathbb{Z},2)) = \mathbb{Z}$.
- 6. P2C? = finite simplicial subset of $K(\mathbb{Z}, 2)$

generated by GP4.

Kenzo calculations (continuation):

5. GP4 = generator of $H_4(K86) = H_4(K(\mathbb{Z},2)) = \mathbb{Z}$.

6. P2C? = finite simplicial subset of $K(\mathbb{Z}, 2)$

generated by GP4.

Next question: P2C? $\stackrel{???}{=} P^2 \mathbb{C}$

Proposition: P2C? = $P^2\mathbb{C}$ \Leftarrow the inclusion P2C? $\hookrightarrow K(\mathbb{Z}, 2)$

induces isomorphisms:

$$H_k(t{P2C?}) \stackrel{\cong \ref{eq: red}}{\longrightarrow} H_k(K(\mathbb{Z},2))$$

for $k \leq 4$.

Proof: Hurewicz-Whitehead Theorem.

$$H_k(t{P2C?}) \stackrel{\cong \regan grade }{\longrightarrow} H_k(K(\mathbb{Z},2))$$

<u>Cone constructor</u>:

P2C? = $P^2\mathbb{C}$

Doctor inclusion $\mathbf{T}_{\mathcal{L}}(\mathbf{T}_{\mathcal{L}}, \mathbf{O})$

$$\mathrm{Cone}(\mathrm{inclusion}) := C_*(\mathbb{P2C?})^{[+1]} \oplus_{\mathrm{inclusion}} C_*(K(\mathbb{Z},2))$$

 $\begin{array}{l} \underline{\text{Proposition:}} \ H_k(\mathbb{P2C?}) \xrightarrow{\cong ??} H_k(K(\mathbb{Z},2)) \ \text{for} \ k \leq 4 \\ \Leftrightarrow \\ H_k(\text{Cone(inclusion)}) = 0 \ \text{for} \ k \leq 5 \end{array}$

 \Leftrightarrow

<u>Proof</u>: Elementary homological algebra.

Kenzo calculations (continuation):

- 5. GP4 = generator of $H_4(K86) = H_4(K(\mathbb{Z},2)) = \mathbb{Z}$.
- 6. P2C? = finite simplicial subset of $K(\mathbb{Z}, 2)$
 - generated by GP4.
- 7. Construction of Cone $\left\{ C_*(P2C?) \stackrel{\text{inclusion}}{\longrightarrow} C_*(K(\mathbb{Z},2)) \right\}$
- 8. Calculation of $H_k(\text{Cone}\{\cdots\})$ for $k \leq 6$.
- 9. $H_k(\text{Cone}) = 0 \text{ for } k \leq 5 \implies P2C? = P^2\mathbb{C}.$

 \Rightarrow a triangulation of $P^2\mathbb{C}$ as P2C? is obtained.

10. The same for higher dimensions.

The END

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