

Constructive

Homological Algebra

II - Homological Perturbations

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;; Clock
Computing
<TnPr <TnPr
End of computing.

;; Clock -> 2002-01-17, 19h 25m 36s.
Computing the boundary of the generator 19 (dimension 7) :
<TnPr <TnPr <TnPr S3 <<Abar[2 S1][2 S1]>>> <<Abar>>> <<Abar>>>
End of computing.

```

Homology in dimension 6 :

Component Z/12Z

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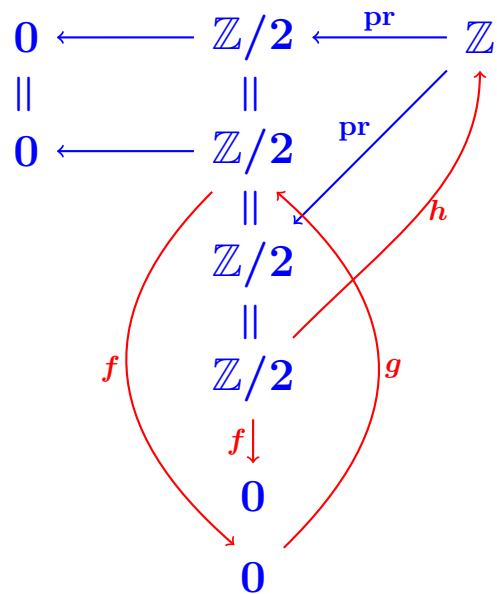
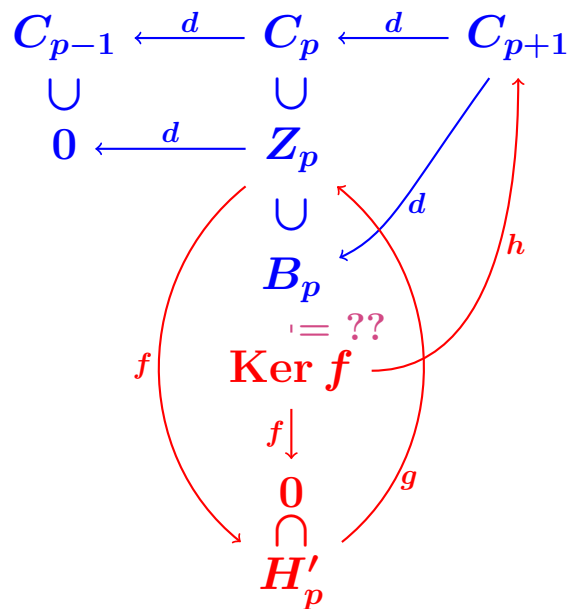
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Ana Romero, Universidad de La Rioja
Julio Rubio, Universidad de La Rioja
Francis Sergeraert, Institut Fourier, Grenoble
Workshop: Formal Methods in Commutative Algebra
Oberwolfach, November 9-14, 2009

Problem when implementing SHPs:



Implementation of $h : \mathbb{Z}/2 \rightarrow \mathbb{Z} ???$

The **Basic Perturbation Lemma (BPL)**

frequently allows to **efficiently overcome** this **difficulty**.

⇒ **Effective Homology I :**

- Notion of **object with effective homology**.
- Most **exact** and **spectral sequences** become **effective**.
- But **spectral sequences** coming from **exact couples** fail!

⇒ **Effective homology II = SHPs** only.

Effective Homology I
flow chart

Functional Programming



“Locally” effective objects



Connections between
Eff. & Loc.Eff. objects



Constructive
Homological Algebra

Reductions between
Chain Complexes



Basic Perturbation Lemma



Functional Programming:

The art of handling **functional** objects.

Examples of functional objects:

$$(\mathbb{Z}, +, -, \times) \quad (\mathbb{Z}[\mathbf{X}], +, -, \times)$$

Other example:

Kan model for the loop space $\Omega S^3 := \mathcal{C}(S^1, S^3)$:

$$(\mathcal{S}_{\Omega S^3}, \{\partial_i^n\}_{n \geq 1, 0 \leq i \leq n}, \{\eta_i^n\}_{n \geq 0, 0 \leq i \leq n})$$

with $\mathcal{S}_{\Omega S^3} =$ the **simplex set** of the **Kan** model.

= “**Locally**” **effective objects**.

Main **problem**:

Designing **programs** $(f_1, \dots, f_n) \mapsto f$.

Example:

$$(\mathfrak{R}, +_{\mathfrak{R}}, -_{\mathfrak{R}}, \times_{\mathfrak{R}}) \mapsto (\mathfrak{R}[X], +_{\mathfrak{R}[X]}, -_{\mathfrak{R}[X]}, \times_{\mathfrak{R}[X]})$$

Topological example. $X = \text{topological space}$.

$$(\mathcal{S}_X, \{\partial(X)_i^n\}_{n \geq 1, 0 \leq i \leq n}, \{\eta(X)_i^n\}_{n \geq 0, 0 \leq i \leq n}) \\ \mapsto (\mathcal{S}_{\Omega X}, \{\partial(\Omega X)_i^n\}_{n \geq 1, 0 \leq i \leq n}, \{\eta(\Omega X)_i^n\}_{n \geq 0, 0 \leq i \leq n})$$

Solution = **λ -calculus, Lisp, ML, Axiom, Haskell...**

Definition: A (homological) reduction is a diagram:

$$\rho: \boxed{h \circlearrowleft \widehat{C}_* \begin{array}{c} \xleftarrow{g} \\ \xrightarrow{f} \end{array} C_*}$$

with:

1. \widehat{C}_* and $C_* =$ chain complexes.
2. f and $g =$ chain complex morphisms.
3. $h =$ homotopy operator (degree +1).
4. $fg = \text{id}_{C_*}$ and $d_{\widehat{C}}h + hd_{\widehat{C}} + gf = \text{id}_{\widehat{C}_*}$.
5. $fh = 0$, $hg = 0$ and $hh = 0$.

$$\begin{array}{c}
 \{ \dots \xrightarrow[h]{d} \widehat{C}_{m-1} \xrightarrow[h]{d} \widehat{C}_m \xrightarrow[h]{d} \widehat{C}_{m+1} \xrightarrow[h]{d} \dots \} = \widehat{C}_* \\
 \{ \dots \xrightarrow[h]{d} \begin{array}{c} A_{m-1} \\ \oplus \\ B_{m-1} \\ \oplus \\ C'_{m-1} \end{array} \xrightarrow[h]{d} \begin{array}{c} A_m \\ \oplus \\ B_m \\ \oplus \\ C'_m \end{array} \xrightarrow[h]{d} \begin{array}{c} A_{m+1} \\ \oplus \\ B_{m+1} \\ \oplus \\ C'_{m+1} \end{array} \xrightarrow[h]{d} \dots \} = \begin{array}{c} A_* \\ \oplus \\ B_* \\ \oplus \\ C'_* \end{array} \\
 \{ \dots \xrightarrow{d} C'_{m-1} \xrightarrow{d} C'_m \xrightarrow{d} C'_{m+1} \xrightarrow{d} \dots \} = C'_* \\
 \{ \dots \xrightarrow{d} C_{m-1} \xrightarrow{d} C_m \xrightarrow{d} C_{m+1} \xrightarrow{d} \dots \} = C_*
 \end{array}$$

$$A_* = \ker f \cap \ker h$$

$$B_* = \ker f \cap \ker d$$

$$C'_* = \operatorname{im} g$$

$$\widehat{C}_* = A_* \oplus B_* \text{ exact} \oplus C'_* \cong C_*$$

Let $\rho: \boxed{h \hookrightarrow \hat{C}_* \begin{matrix} \xleftarrow{g} \\ \xrightarrow{f} \end{matrix} C_*}$ be a **reduction**.

Frequently:

1. \hat{C}_* is a **locally effective chain complex**:
its **homology groups** are **unreachable**.
2. C_* is an **effective chain complex**:
its **homology groups** are **computable**.
3. The **reduction** ρ is an entire description of
the **homological nature** of \hat{C}_* .
4. Any **homological problem** in \hat{C}_* is **solvable**
thanks to the **information** provided by ρ .

$$\rho: \boxed{h \circlearrowleft \hat{C}_* \begin{matrix} \xleftarrow{g} \\ \xrightarrow{f} \end{matrix} C_*}$$

1. What is $H_n(\hat{C}_*)$? Solution: Compute $H_n(C_*)$.

2. Let $x \in \hat{C}_n$. Is x a cycle? Solution: Compute $d_{\hat{C}_*}(x)$.

3. Let $x, x' \in \hat{C}_n$ be cycles. Are they homologous?

Solution: Look whether $f(x)$ and $f(x')$ are homologous.

4. Let $x, x' \in \hat{C}_n$ be homologous cycles.

Find $y \in \hat{C}_{n+1}$ satisfying $dy = x - x'$?

Solution:

(a) Find $z \in C_{n+1}$ satisfying $dz = f(x) - f(x')$.

(b) $y = g(z) + h(x - x')$.

Definition: $(C_*, d) =$ given chain complex.

A perturbation $\delta : C_* \rightarrow C_{*-1}$ is an operator of degree -1

satisfying $(d + \delta)^2 = 0$ ($\Leftrightarrow (d\delta + \delta d + \delta^2) = 0$):

$$(C_*, d) + (\delta) \mapsto (C_*, d + \delta).$$

Problem: Let $\rho: \boxed{h \hookrightarrow (\widehat{C}_*, \widehat{d}) \xrightleftharpoons[f]{g} (C_*, d)}$ be a given reduction and $\widehat{\delta}$ a perturbation of \widehat{d} .

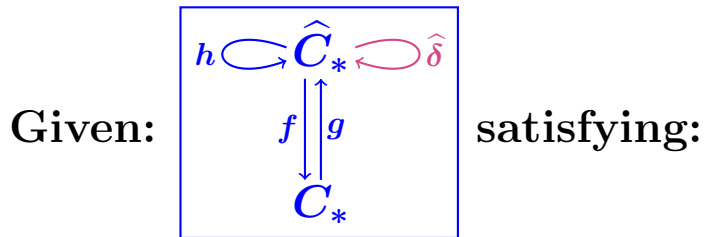
How to determine a new reduction:

$$???: \boxed{h_{+?} \hookrightarrow (\widehat{C}_*, \widehat{d} + \widehat{\delta}) \xrightleftharpoons[f_{+?}]{g_{+?}} (C_*, d_{+?})}$$

describing in the same way the homology of

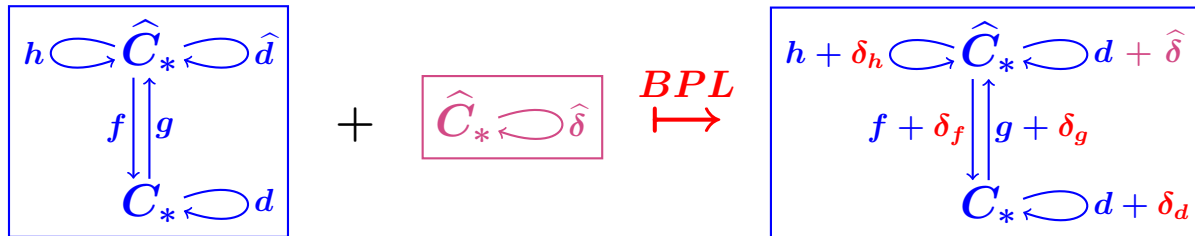
the chain complex with the perturbed differential?

Basic Perturbation “Lemma” (BPL):



1. $\hat{\delta}$ is a **perturbation** of the differential \hat{d} of \hat{C}_* ;
2. The operator $h \circ \hat{\delta}$ is **pointwise nilpotent**.

Then a **general algorithm BPL** constructs:



Proof:

$\phi := \sum_{i=1}^{\infty} (-1)^i (h\hat{\delta})^i$ and $\psi := \sum_{i=1}^{\infty} (-1)^i (\hat{\delta}h)^i$ are defined.

Then:

- $\delta_d := f\hat{\delta}(\text{id}_{\hat{C}} + \phi)g = f(\text{id}_{\hat{C}} + \psi)\hat{\delta}g$
- $\delta_f := f\psi$
- $\delta_g := \phi g$
- $\delta_h := \phi h = h\psi$

is the **solution**.

QED

Serre: “Everything” in Algebraic Topology

can be reduced to Fibration problems.

Examples: Loop spaces, Classifying spaces, Homogeneous spaces, Whitehead tower, Postnikov tower, ...

Remark: Fibration = Twisted Product

= Perturbation of Trivial Product.

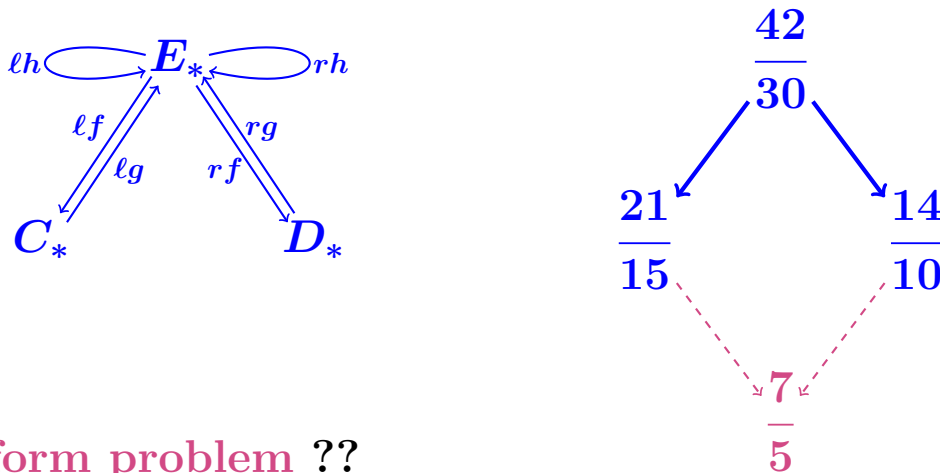
Corollary: BPL is effective

+ Fibration = Perturbation of Trivial Product

+ Everything is Fibration

⇒ Alg. Topology becomes Constructive.

Definition: A (strong chain-) equivalence $\varepsilon : C_* \rightleftarrows D_*$ is a pair of reductions $C_* \xrightarrow{\ell\rho} E_* \xrightarrow{r\rho} D_*$:



Normal form problem ??

More structure often necessary in C_* .

Most often: no possible choice for C_* .

Definition: An **object with effective homology** X
is a 4-tuple:

$$X = \boxed{X, C_*(X), EC_*, \varepsilon}$$

with:

1. X = an arbitrary **object** (simplicial set, simplicial group, differential graded algebra, ...)
2. $C_*(X)$ = “the” **chain complex** “traditionally” associated with X to define the **homology groups** $H_*(X)$.
3. EC_* = some **effective chain complex**.
4. ε = some **equivalence** $C_*(X) \overset{\varepsilon}{\rightleftarrows} EC_*$.

Main result of effective homology:

Meta-theorem: Let X_1, \dots, X_n be a collection of **objects** with **effective homology** and ϕ be a **reasonable construction process**:

$$\phi : (X_1, \dots, X_n) \mapsto X.$$

Then **there exists a version with effective homology** ϕ_{EH} :

$$\phi_{EH}: \left(\boxed{X_1, C_*(X_1), EC_{1*}, \varepsilon_1}, \dots, \boxed{X_n, C_*(X_n), EC_{n*}, \varepsilon_n} \right) \mapsto \boxed{X, C_*(X), EC_*, \varepsilon}$$

The process is **perfectly stable**

and can be **again used** with X for **further calculations**.

Example:

Julio Rubio's solution of Adams' problem.

$$X = (X, C_*(X), EC_*^X, \epsilon^X)$$



Eil.-Moore_{EH}

$$\Omega X = (\Omega X, C_*(\Omega X), EC_*^{\Omega X}, \epsilon^{\Omega X})$$

⇒ Trivial iteration now available.

⇒ Very simple solution of Adam's problem :

Indefinite iteration of the Cobar construction ???

$$X = (X, C_*(X), EC_*^X, \varepsilon^X)$$

$$\Downarrow \Omega_{EH}$$

$$\Omega X = (\Omega X, C_*(\Omega X), EC_*^{\Omega X}, \varepsilon^{\Omega X})$$

$$\Downarrow \Omega_{EH}$$

$$\Omega^2 X = (\Omega^2 X, C_*(\Omega^2 X), EC_*^{\Omega^2 X}, \varepsilon^{\Omega^2 X})$$

$$\Downarrow \Omega_{EH}$$

$$\Omega^3 X = (\Omega^3 X, C_*(\Omega^3 X), EC_*^{\Omega^3 X}, \varepsilon^{\Omega^3 X})$$

$$\Downarrow \Omega_{EH}$$

$$\Omega^4 X = \dots$$

“Cobar” 3 (EC_*^X)

Example: **Effective homology version of**
the Serre spectral sequence.

$$\begin{aligned}
 & F = (F, C_*(F), EC_*^F, \varepsilon^F) \\
 + & B = (B, C_*(B), EC_*^B, \varepsilon^B) \\
 + & \tau : B \rightarrow F \\
 & \Downarrow \Downarrow \Downarrow \Downarrow \Downarrow \text{Serre}_{EH} \\
 & E = F \times_{\tau} B = (E, C_*(E), EC^E, \varepsilon^E)
 \end{aligned}$$

(Serre + G. Hirsch + H. Cartan + Shih W.

+ Szczarba + Ronnie Brown + J. Rubio + FS)

Proof.

$$\begin{array}{ccc}
 C_*(F \times B) & \xrightarrow{\text{id}} & C_*(F \times B) \xrightarrow{EZ} C_*F \otimes C_*B \\
 & \simeq & \\
 C_*F \otimes C_*B & \xrightarrow{\otimes} & \widehat{C}^F \otimes \widehat{C}^B \xrightarrow{\otimes} EC^F \otimes EC^B
 \end{array}$$

↓↓↓↓↓↓ Serre_{EH}

$$\begin{array}{ccc}
 C_*(F \times_{\tau} B) & \xrightarrow{\text{id}} & C_*(F \times_{\tau} B) \xrightarrow{\text{Shih}} C_*F \otimes_{\tau} C_*B \\
 & \simeq & \\
 C_*F \otimes_{\tau} C_*B & \xrightarrow{EPL} & \widehat{C}^F \otimes_{\tau'} \widehat{C}^B \xrightarrow{BPL} EC^F \otimes_{\tau''} EC^B
 \end{array}$$

+ Composition of equivalences \implies O.K.

Combining these ingredients \Rightarrow

Homological Algebra becomes **constructive**.

Corollary: The “standard” **exact and spectral sequences**
of **Homological Algebra**

really become **computational tools**.

\Rightarrow Concrete **computer programs** (**EAT**, **Kenzo**).

Warning about the **right chronology**. Example, let:

$$\phi : F \hookrightarrow E \twoheadrightarrow B$$

be a fibration with B simply connected.

1. The ordinary Serre Spectral Sequence **is not** constructive.

2. Methods of **Effective Homology** give an **algorithm**:

$$[EH_*(F) + EH_*(B) + \phi] \mapsto EH_*(E).$$

3. Methods of **Effective Homology** can **then** compute:

$$[EH_*(F) + EH_*(B) + \phi + EH_*(E) \mapsto \text{SSS}(\phi)].$$

That is, the **SSS** is a **byproduct** of **Effective Homology**.

(Ana Romero)

Are your **computer programs** efficient?

What about **benchmarks**?

Do you **compute** new **sphere homotopy groups**?

Non-relevant question!

To be compared with **prime number chasing**.

Two **different** activities:

1. Searching for **very big prime numbers**.
2. Designing methods applicable to **arbitrary numbers**.

Most efficient current methods for big prime numbers:

Specific tests (Lucas-Lehmer):

$\Rightarrow (2^{32,582,657} - 1)$ prime (2006).

Most “efficient” general method:

Agrawal-Kayal-Saxena n^{12} -algorithm.

Both methods have totally different scopes.

Analogous situation in **Algebraic Topology**.

In case of **spheres** (\sim **Mersenne numbers**),

specific methods go very far.

But these methods are **inapplicable in general situations**.

Typical situation:

Modifying a **loop space** by “**pre-attaching**” a **cell**.

What **influence about**

the **homology groups** of the **new loop space**?

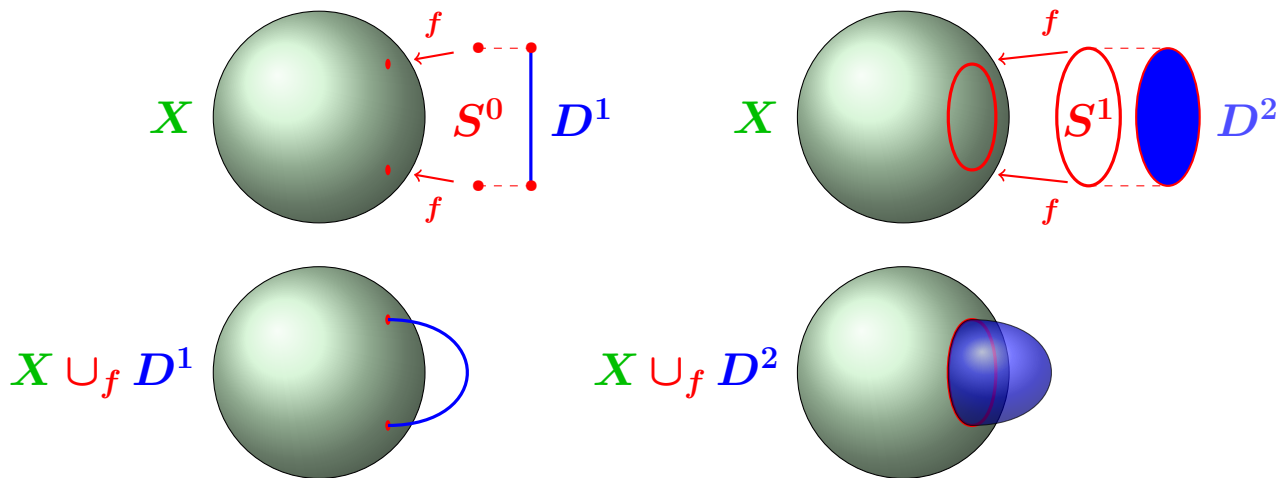
Attaching a cell D^n to a topological space

along the boundary S^{n-1} :

$X =$ Topological space.

$f : S^{n-1} \rightarrow X =$ continuous map.

$\Rightarrow X \cup_f D^n := (X \amalg D^n) / (X \ni f(x) \sim x \in S^{n-1})$.



Given:

- $X =$ Topological space.
- $H_*(X) =$ Homology of X .
- $H_*(\Omega X) =$ Homology of the loop space ΩX .
- $f : S^{n-1} \rightarrow X$ continue.

Problem:

- Determine $H_*(\Omega(X \cup_f D^n)) = ???$

Example:

$H_*(\Omega S^3)$ computed by **J.-P. Serre** (1950).

$H_*(\Omega^2 S^3 := \Omega(\Omega S^3))$ computed by **W. Browder** (1958).

Modifying $\Omega S^3 \mapsto \Omega S^3 \cup_2 D^3$.

\Rightarrow **Problem:**

$$H_*(\Omega(\Omega S^3 \cup_2 D^3)) = ???$$

Remark: $H_*(\Omega S^3 \cup_2 D^3)$ direct consequence of **Serre's** result.

“Standard” **Algebraic Topology** \Rightarrow ???

In **Effective Homology**:

$$S^3 = \text{Finite simplicial set} \Rightarrow S^3 = \text{OEH}^{1)}.$$

$$\text{EMSSEH}^{2)} \Rightarrow \Omega S^3 = \text{OEH}^{1)}.$$

$$\text{MVESEH}^{3)} \Rightarrow \Omega S^3 \cup_2 D^3 = \text{OEH}^{1)}.$$

$$\text{EMSSEH}^{2)} \Rightarrow \Omega(\Omega S^3 \cup_2 D^3) = \text{OEH}^{1)}.$$

$$\Omega(\Omega S^3 \cup_2 D^3) = \text{OEH}^{1)} \Rightarrow H_*(\Omega(\Omega S^3 \cup_2 D^3)) \boxed{\text{computable}}.$$

\Rightarrow **OK.**

1) OEH = Object with Effective Homology.

2) EMSSEH = Eilenberg-Moore Spectral Sequence with Effective Homology.

3) MVESEH = Mayer-Vietoris Exact Sequence with Effective Homology

A more complicated analogous computation:

$$X = \Omega(\Omega(\Omega(P^\infty\mathbb{R}/P^3\mathbb{R}) \cup_4 D^4) \cup_2 D^3))$$

$$H_*X = ???$$

$$H_0(X) = \mathbb{Z}.$$

$$H_1(X) = \mathbb{Z}/2\mathbb{Z}.$$

$$H_2(X) = (\mathbb{Z}/2\mathbb{Z})^2 + \mathbb{Z}.$$

$$H_3(X) = (\mathbb{Z}/2\mathbb{Z})^4 + \mathbb{Z}/8\mathbb{Z}.$$

$$H_4(X) = (\mathbb{Z}/2\mathbb{Z})^{10} + \mathbb{Z}/4\mathbb{Z} + \mathbb{Z}^2.$$

$$H_5(X) = (\mathbb{Z}/2\mathbb{Z})^{23} + \mathbb{Z}/8\mathbb{Z} + \mathbb{Z}/16\mathbb{Z}.$$

$$H_6(X) = (\mathbb{Z}/2\mathbb{Z})^{52} + (\mathbb{Z}/4\mathbb{Z})^3 + \mathbb{Z}^3.$$

$$H_7(X) = (\mathbb{Z}/2\mathbb{Z})^{113} + \mathbb{Z}/4\mathbb{Z} + (\mathbb{Z}/8\mathbb{Z})^3 + \mathbb{Z}/16\mathbb{Z} + \mathbb{Z}/32\mathbb{Z} + \mathbb{Z}.$$

The **longest Kenzo computation** (2 months).

Two main steps in a **Kenzo calculation** $H(d) = ???$

- 1) “Automatic” **writing** of a sophisticated **highly functional program** P .
- 2) Using **program** P to compute $P(d) = H(d)$.

1) \Rightarrow Always **very fast** (< 1 sec.).

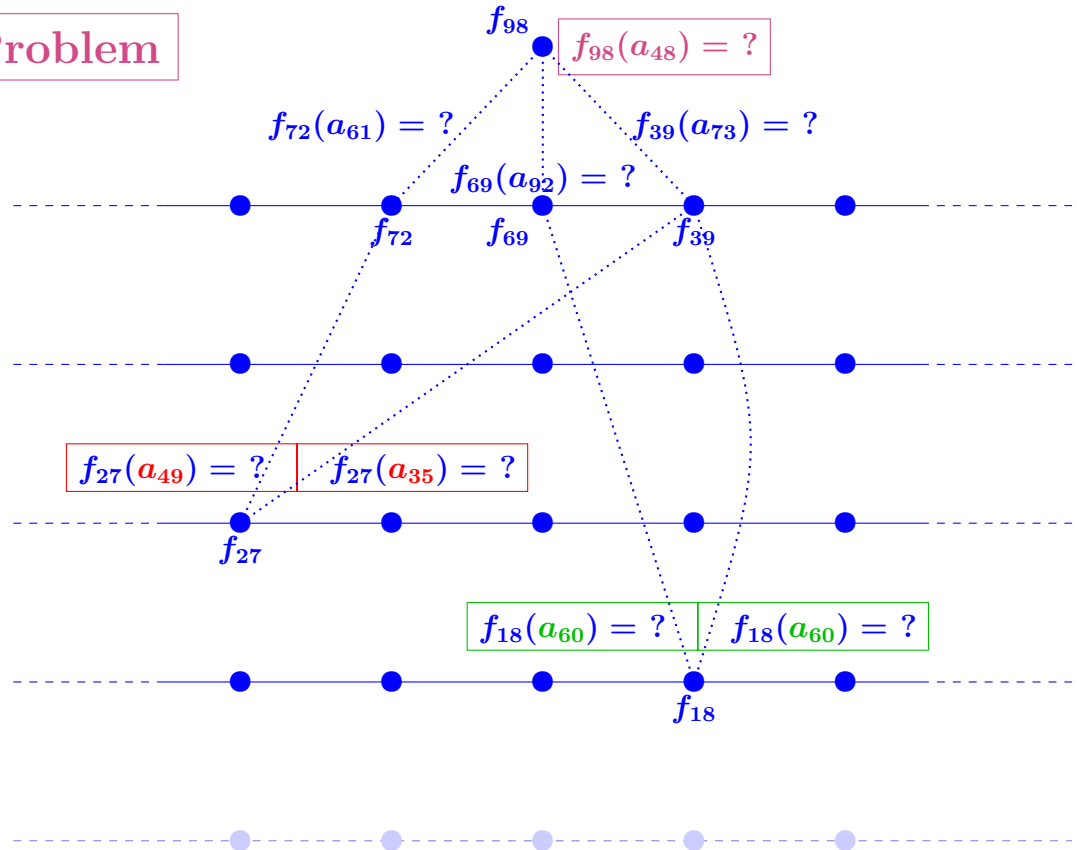
2) \Rightarrow Can be **very long** (hours, days, months, years, centuries, ...)

High efficiency in **functional programming** with **Common-Lisp**

\Rightarrow **No technical difficulty** in 1).

Terrible problem of memory management in 2).

Functional Problem



Only **heuristic** methods available:

EAT-1: No result stored \Rightarrow **poor time computing efficiency.**

EAT-2: “**Non-trivial**” results stored
 \Rightarrow Computing time **divided by ~ 10 .**

Kenzo-1: Strong improvement in **storing-searching** methods.
 \Rightarrow Computing time **divided by ~ 10 .**

Kenzo-2: For overcoming **space complexity:**
Periodic **cleaning** of stored results.
 \Rightarrow Computing time **divided by ~ 10 .**

Theoretical framework for a **rational** study ??? **Open !!!**

“Vertical” vs “Horizontal” time complexity.

Computing $H_n(X)$, $\pi_n(X)$...

Two parameters: n and X .

“Horizontal” complexity := wrt X .

“Vertical” complexity := wrt n .

Effective homology \Rightarrow Horizontal complexity = P .

David Annick (1986) \Rightarrow

Vertical complexity \geq NP -complete.

Back to “standard” Mathematics.

Traditional main problem of Algebraic Topology:

Classifying the homotopy types.

1. Only “reasonable” spaces: CW-complexes \cong Simp. sets.
2. Non-simply connected topology excluded (word problem).
3. Classification “up to homeomorphism” out of scope \Rightarrow
Only classification “up to homotopy equivalence”.
4. “Standard” solution = Postnikov “invariants”.

Main problem of Algebraic Topology:

Algebraic Models for Topological Spaces ?

Main idea: Topology is difficult, Algebra is easy (!?).

Subquestion: what does the word “Algebra” means?

Answer: No meaning at all, only a “cultural” tradition.

Correct question:

Computable Models for Topological Spaces ?

Three solutions:

1. Rolf Schön.
2. Effective Homology.
3. Operads.

Schön's solution =

Intensive use of inductive limits

to approximate infinite objects.

Only one computer application:

Alain Clément, Lausanne, Haskell program.

Comparison: **Effective Homology** \leftrightarrow **Operads** ???

Object with effective homology = Triple: (X, HX, ε) with:

X = **locally effective** version of the object.

HX = **Effective chain complex**

describing the **ordinary homology** of X .

ε = Strong connection $X \xleftrightarrow{\varepsilon} HX$.

Theorem: The triple (X, HX, ε) is a computable model
of the **homotopy type** of X .

Operadic model for a topological space X :

$$(HX, \mathcal{M})$$

with:

$HX =$ Effective chain complex

describing the ordinary homology of X .

$\mathcal{M} = E_\infty$ -operadic structure over HX .

Theorem (Mandell): The pair (HX, \mathcal{M}) is

an “algebraic” model

of the homotopy type of X .

Connection between (X, HX, ε) and (HX, \mathcal{M}) ?

Theorem: There exists a **canonical correspondance**:

$$(X, HX, \varepsilon) \longleftrightarrow (HX, \mathcal{M})$$

1. “ \longrightarrow ” = **Berger-Fresse**.
2. “ \longleftarrow ” = **S.-Mandell**.

Far from concrete implementations !!

Main open problems in **Effective Homology**:

1. **Eilenberg-Zilber** = “ $\square \longleftrightarrow \square$ ”.

Problem: General **formula**

of unavoidable **exponential complexity**.

How to design an **efficient algorithm**

for concrete particular cases?

2. **Twisted Eilenberg-Zilber**.

New important results **experimentally** discovered in 98.

Not yet proved!

3. Spectral sequences.

Filtrated chain complexes vs Exact Couples.

Particular cases of Bousfield-Kan, Adams,

May, Adams-Novikov... spectral sequences.

Cf recent thesis by Ana Romero.

4. Commutative Algebra.

Recent result:

Canonical correspondance between:

Effective resolution of $K[x_1, \dots, x_m]$ -modules

\cong

Effective homology of Koszul complex.

\Rightarrow New algorithms producing effective resolutions.

5. Concrete implementation of the canonical correspondance:

$$(X, HX, \varepsilon) \longleftrightarrow (HX, \mathcal{M})$$

6. Efficient memory management

for high level functional programming ???

7. Program proof, theorem proving.

Recent result (Jesus Aransay):

Isabelle-certified proof of the Basic-Perturbation-Lemma.

Competing work by Coq people.

The END

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