

Constructive Homological Algebra I - The Problem

```
;; Clock
Computing
<TnPr <TnPr
End of computing.
```

```
;; Clock -> 2002-01-17, 19h 25m 36s.
Computing the boundary of the generator 19 (dimension 7) :
<TnPr <TnPr <TnPr S3 <<Abar[2 S1][2 S1]>>> <<Abar>>> <<Abar>>>
End of computing.
```

Homology in dimension 6 :

Component Z/12Z

---done---

```
;; Clock -> 2002-01-17, 19h 27m 15s
```

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Semantics of colours:

Blue = “Standard” Mathematics

Red = Constructive, effective,
algorithm, machine object, ...

Violet = Problem, difficulty, obstacle, disadvantage, ...

Green = Solution, essential point, mathematicians, ...

Dark Orange = Fuzzy objects.

Pale grey = Hyper-Fuzzy objects.

1953 = Birth date of

the computational problem in Algebraic Topology.

Theorem (Serre): In simply connected Algebraic Topology,
the homology and homotopy groups
of the “reasonable” spaces
are \mathbb{Z} -modules of finite type.

\Rightarrow Problem: $\exists?$ algorithms:

$$X \mapsto H_*(X) \quad ???$$

$$X \mapsto \pi_*(X) \quad ???$$

Example of “reasonable” space:

$$X = \Omega(D^3 \cup_2 \Omega(D^4 \cup_4 \Omega(P^\infty(\mathbb{R})/P^3(\mathbb{R}))))$$

Example of problem: $H_*X = ???$

Kenzo program + 2 months of computation \Rightarrow

$$H_0(X) = \mathbb{Z}$$

$$H_1(X) = \mathbb{Z}/2$$

$$H_2(X) = (\mathbb{Z}/2)^2 + \mathbb{Z}$$

$$H_3(X) = (\mathbb{Z}/2)^4 + \mathbb{Z}/8$$

$$H_4(X) = (\mathbb{Z}/2)^{10} + \mathbb{Z}/4 + \mathbb{Z}^2$$

$$H_5(X) = (\mathbb{Z}/2)^{23} + \mathbb{Z}/8 + \mathbb{Z}/16$$

$$H_6(X) = (\mathbb{Z}/2)^{52} + (\mathbb{Z}/4)^3 + \mathbb{Z}^3$$

$$H_7(X) = (\mathbb{Z}/2)^{113} + \mathbb{Z}/4 + (\mathbb{Z}/8)^3 + \mathbb{Z}/16 + \mathbb{Z}/32 + \mathbb{Z}$$

First important result:

Theorem (Edgar Brown, 1956):

X = finite simply connected simplicial set

$\Rightarrow \pi_* X$ computable.

Edgar Brown's warning:

It must be emphasized that although the procedures developed for solving these problems are finite, they are much too complicated to be considered practical.

Brown's warning still **valid** today!

Main currently available **solutions**:

1. **Edgar Brown** \Rightarrow **Rolf Schön** \Rightarrow **Alain Clément** \Rightarrow ???
2. **Operadic solution** (**Justin Smith** + ... + **Michael Mandell**)
Implementation ???
3. **Effective Homology I** (\Rightarrow **Kenzo**) + **II** (2008).

Main obstacle: Infinite intermediate objects.

Typical simple example: $\pi_4(S^3) = ?$

Cartan-Serre-Whitehead method:

$H_2(S^3) = 0 + H_3(S^3) = \mathbb{Z} \Rightarrow$ fibration:

$$K(\mathbb{Z}, 2) \hookrightarrow X_4 \longrightarrow S^3$$

with $\pi_p(X_4) = \pi_p(S^3)$ for every $p \neq 3$ and $\pi_3(X_4) = 0$.

$\Rightarrow \pi_p(X_4) = 0$ for $p < 4$

\Rightarrow (Hurewicz' theorem) $\pi_4(X_4) = H_4(X_4)$

but the standard model for $K(\mathbb{Z}, 2)$ is infinite

and X_4 cannot be “totally” installed on a machine.

Main methods invented by the **topologists**

to **overcome** this **obstacle**:

Exact sequences and spectral sequences

Typically, if $F \hookrightarrow E \rightarrow B$ is a **fibration**,

the **Serre spectral sequence** is

a(n enormous) set of (very sophisticated) **relations**

connecting the **groups** $H_p(F)$, $H_p(E)$ and $H_p(B)$.

In **some particular** cases this can be a tool:

$$\{H_p(B), H_p(F)\}_{p \in \mathbb{N}} \mapsto \{H_p(E)\}_{p \in \mathbb{N}}$$

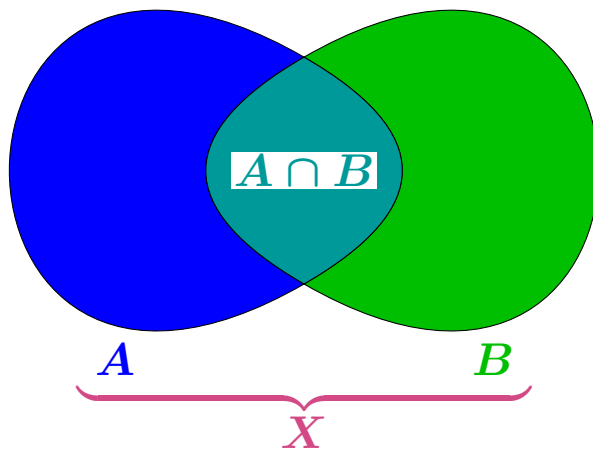
In general not!

Typical example:

Theorem (Mayer-Vietoris): $X = \text{space}$ covered by $\{A, B\}$.

\Rightarrow canonical exact sequence:

$$\begin{aligned} \dots \rightarrow H_p(A \cap B) &\xrightarrow{\alpha_p} H_p(A) \oplus H_p(B) \rightarrow H_p(X) \rightarrow \\ &\rightarrow H_{p-1}(A \cap B) \xrightarrow{\alpha_{p-1}} H_{p-1}(A) \oplus H_{p-1}(B) \rightarrow \dots \end{aligned}$$



How to use Mayer-Vietoris?

Long exact sequence:

$$A_* \xrightarrow{\alpha} B_* \rightarrow C_* \rightarrow D_* \xrightarrow{\beta} E_*$$

A_*, B_*, D_*, E_* given $\Rightarrow C_* = ???$

Long exact sequence \Rightarrow short exact sequence:

$$0 \rightarrow \text{Coker } \alpha \rightarrow C_* \rightarrow \text{Ker } \beta \rightarrow 0$$

But most often α and β unknown!

If α and β are known,

the extension class $\tau \in H^2(\text{Ker } \beta; \text{Coker } \alpha)$

giving the right extension $C_* = \text{Coker } \alpha \boxtimes_{\tau} \text{Ker } \beta$

can be very hard to be computed.

Analysis of the **obstacle**:

Standard **Algebraic Topology** is **not constructive!**

Example: Construction of $\alpha : H_p(A \cap B) \rightarrow H_p(A) \oplus H_p(B)$?

Let us assume $H_p(A \cap B) = \mathbb{Z}/6$, $H_p(A) = \mathbb{Z}/7$, $H_p(B) = \mathbb{Z}/8$.

Meaning of $H_p(A) = \mathbb{Z}/7$??

\exists isomorphism $H_p(A) \xrightarrow{\cong} \mathbb{Z}/7$.

Most often, this existence is **not constructive!!**

How to organize standard **Homological Algebra**

to **construct** and **manipulate**

such **constructive isomorphisms?**

“Actual” Homology Group = subquotient: $H_p = Z_p/B_p$.

Critical diagram:

$$\begin{array}{ccccc}
 C_{p-1} & \xleftarrow{d} & C_p & \xleftarrow{d} & C_{p+1} \\
 \cup & & \cup & & \cup \\
 0 & \xleftarrow{d} & Z_p & & \\
 & & \cup & & \uparrow \\
 & & B_p & & \\
 & & \vdots = ?? & & \\
 & & \text{Ker } f & & \\
 & & \downarrow f & & \\
 & & 0 & & \\
 & & \cup & & \\
 & & H'_p & &
 \end{array}$$

The diagram includes several red arrows:

- A red arrow labeled f from $\text{Ker } f$ to 0 .
- A red arrow labeled g from 0 to H'_p .
- A red arrow labeled h from H'_p to C_{p+1} .
- A red arrow labeled f from C_{p-1} to H'_p .

Definition:

Constructive isomorphism

= (f,g,h) with:

f morphism satisfies $fd = 0$.

g map satisfies $fg = \text{id}$.

h map satisfies $dh = \text{id}$.

H'_p hyp. $\cong H_p$.

Chain complex: $C_* = \{\cdots \leftarrow C_{q-1} \leftarrow C_q \leftarrow C_{q-1} \leftarrow \cdots\}$

Definition: A **SHP** (**S**olution for the **H**omological **P**roblem)

for C_* is a family $(H'_p, f_p, g_p, h_p)_{p \in \mathbb{Z}}$ satisfying:

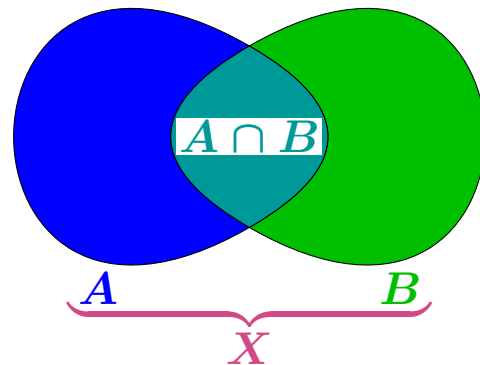
- $\{H'_p\}_{p \in \mathbb{Z}}$ = family of **effective groups**.
- $\{f_p : Z_p \rightarrow H'_p\}_{p \in \mathbb{Z}}$ = **morphism family** with $f_p d_{p+1} = 0$.
- $\{g_p : H'_p \rightarrow Z_p\}_{p \in \mathbb{Z}}$ = **map family** with $f_p g_p = \text{id}$.
- $\{h_p : \text{Ker } f_p \rightarrow C_{p+1}\}_{p \in \mathbb{Z}}$ = **map family**
satisfies $d_{p+1} h_p = \text{id}$.

Mayer-Vietoris revisited.

C_*A , C_*B , $C_*(A \cap B)$

with SHPs given.

$H_*X = ???$



$$\begin{aligned} \dots \rightarrow H_p(A \cap B) &\xrightarrow{\alpha_p} H_p(A) \oplus H_p(B) \rightarrow H_p(X) \rightarrow \\ &\rightarrow H_{p-1}(A \cap B) \xrightarrow{\alpha_{p-1}} H_{p-1}(A) \oplus H_{p-1}(B) \rightarrow \dots \end{aligned}$$

Constructing: $\alpha_p : H_p(A \cap B) \rightarrow H_pA \oplus H_pB$???

Solution:

$$H'_p(A \cap B) \xrightarrow{g_p} Z_p(A \cap B) \xrightarrow{i_p} Z_pA \oplus Z_pB \xrightarrow{f_p} H'_pA \oplus H'_pB$$

\Rightarrow OK !!

Moral:

Constructive Homological Algebra

is a matter of being able to
 appropriately **compute**, **handle** and “**proof**”
representants for **homology classes**.

Effective Homology I (BPL) \Rightarrow

“Simple” **exact** and **spectral sequences** become **effective**.

Effective Homology II (SHPs) \Rightarrow

Sophisticated spectral sequences

(**Bockstein**, **Bousfield-Kan**, ...)

defined through **exact couples**

become **effective**.

The main ingredients of **Homological Algebra** are
a **ground ring R**
and **chain complexes** made of **R -modules** and **R -morphisms**.

Three classes of **modules**:

- **Effective modules.**
- **Locally Effective modules.**
- **Fuzzy modules.**

1. Effective module M .

The membership property of an arbitrary object $a \in? M$
is decidable.

The module is discrete:
equality between objects is decidable.

Ordinary computations can be executed.

If $\alpha \in R$ and $a, b \in M$, algorithms compute $a + b$ and αa .

The global structure of M is known.

“Reasonable” questions about M can be answered.

Depending on R , the isomorphism problem
between two effective R -modules M and M'
is or is not decidable.

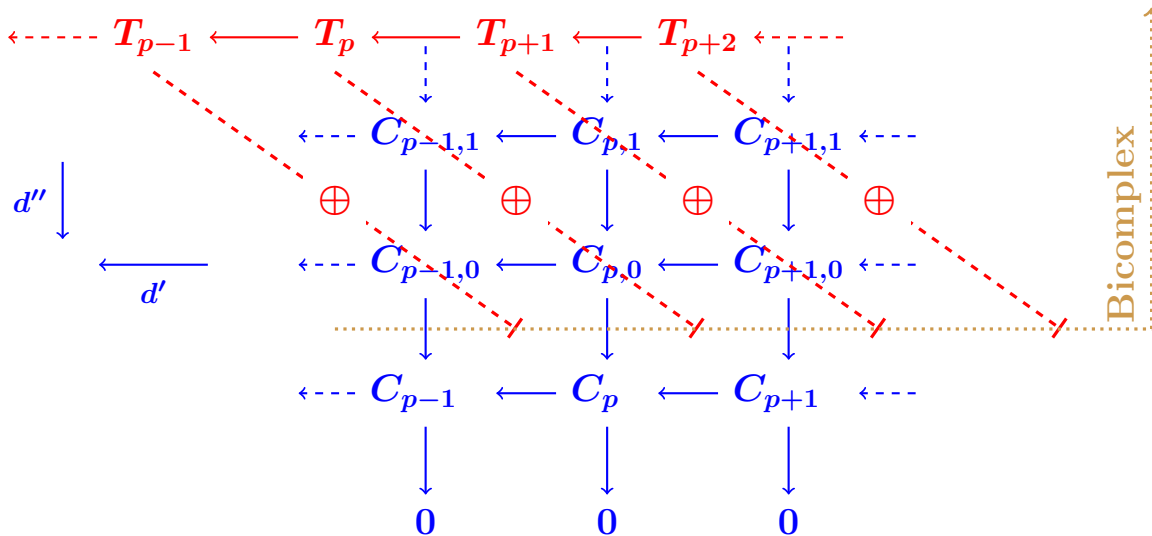
$R = \text{Cramer ring}$ (= strongly discrete and coherent):

- Any linear system $LV = 0$,
with $L \in \text{Hom}(R^m, R^n)$ given, $V \in R^m$ unknown,
has a complete “solution” $S \in \text{Hom}(R^k, R^m)$:
 $\text{Ker } L = \text{Im } S$.

$$R^n \xleftarrow{L} R^m \xleftarrow{S} R^k$$

- For any $L \in \text{Hom}(R^m, R^n)$,
an algorithm $\sigma : R^n \rightarrow \{\perp\} \amalg R^m$ satisfies:
 - $\sigma(V) = \perp \Leftrightarrow V \notin \text{Im } L$.
 - $\sigma(V) = U \in R^m \Leftrightarrow V = LU$.

Bicomplex theorem:



$$\left. \begin{array}{l} \text{Every column exact} \\ T_p = \bigoplus_{a+b=p} C_{a,b} \\ d'd' = d''d'' = d'd'' + d''d' = 0 \end{array} \right\} \Rightarrow H_*(C_*, d') \cong H_*(T_*, d' \oplus d'')$$

$R =$ Cramer ring.

Definition: An **effective module** M is an R -module

with a **finite presentation**:

$$M = (R^{b_0}, R^{b_1}, d) \Leftrightarrow 0 \leftarrow M \leftarrow R^{b_0} \xleftarrow{d} R^{b_1}$$

R^{b_0} = type of the elements of M .

$m \sim m' \bmod d(R^{b_1})$ **decidable** $\Rightarrow M$ discrete.

Proposition: $M =$ **effective module**.

$\Rightarrow M$ admits a **free resolution of finite type**.

$$0 \leftarrow M \leftarrow R^{b_0} \leftarrow R^{b_1} \leftarrow R^{b_2} \leftarrow R^{b_3} \leftarrow \dots$$

Proof: Cramer.

QED

Proposition: An algorithm produces a SHP

for a chain complex of free ft modules

$$\begin{array}{ccccc}
 C_{p-1} & \xleftarrow{d} & C_p & \xleftarrow{d} & C_{p+1} \\
 \cup & & \cup & & \cup \\
 0 & \xleftarrow{d} & Z_p & & \\
 & & \cup & & \\
 & & B_p & & \\
 & & \downarrow f & & \\
 & & \text{Ker } f & & \\
 & & \downarrow f & & \\
 & & 0 & & \\
 & & \cup & & \\
 & & H'_p & &
 \end{array}$$

$f: C_p \rightarrow \text{Ker } f$, $g: \text{Ker } f \rightarrow H'_p$, $h: C_{p+1} \rightarrow H'_p$, $d: C_p \rightarrow Z_p$, $d: Z_p \rightarrow B_p$, $d: B_p \rightarrow \text{Ker } f$.

$$\begin{array}{ccccc}
 R^\alpha & \xleftarrow{d} & R^\beta & \xleftarrow{d} & R^\gamma \\
 \cup & & \cup & & \cup \\
 0 & \xleftarrow{d} & Z & & \\
 & & \cup & & \\
 & & B & & \\
 & & \downarrow f & & \\
 & & R^\delta & & \\
 & & \downarrow \varphi & & \\
 & & R^\varepsilon & &
 \end{array}$$

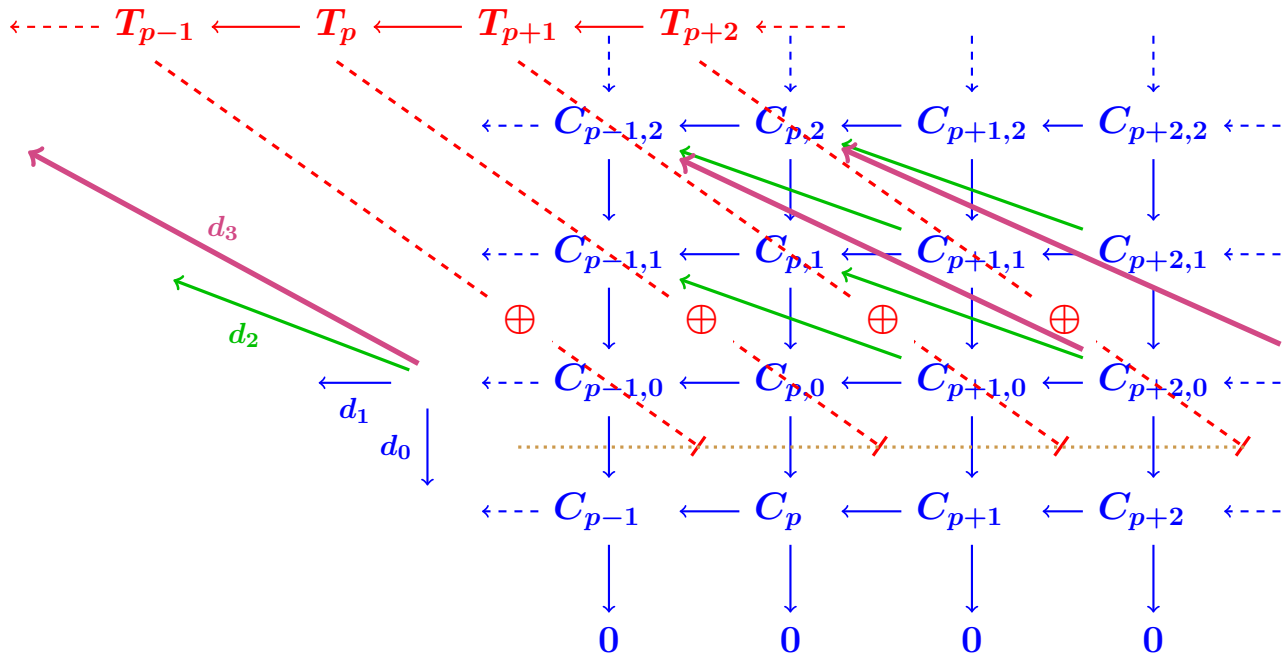
$f: B \rightarrow R^\delta$, $g: R^\gamma \rightarrow R^\delta$, $h: R^\gamma \rightarrow R^\delta$, $\psi: R^\gamma \rightarrow R^\delta$, $\phi: R^\delta \rightarrow R^\varepsilon$, $d: R^\beta \rightarrow Z$, $d: Z \rightarrow B$, $d: R^\alpha \rightarrow 0$.

$$\begin{array}{ccccc}
 0 & \leftarrow & Z & \leftarrow & R^\gamma \\
 & & \uparrow & & \uparrow \\
 & & R^\delta & \leftarrow & R^\gamma \\
 & & \uparrow & & \uparrow \\
 & & R^\varepsilon & \leftarrow & 0
 \end{array}$$

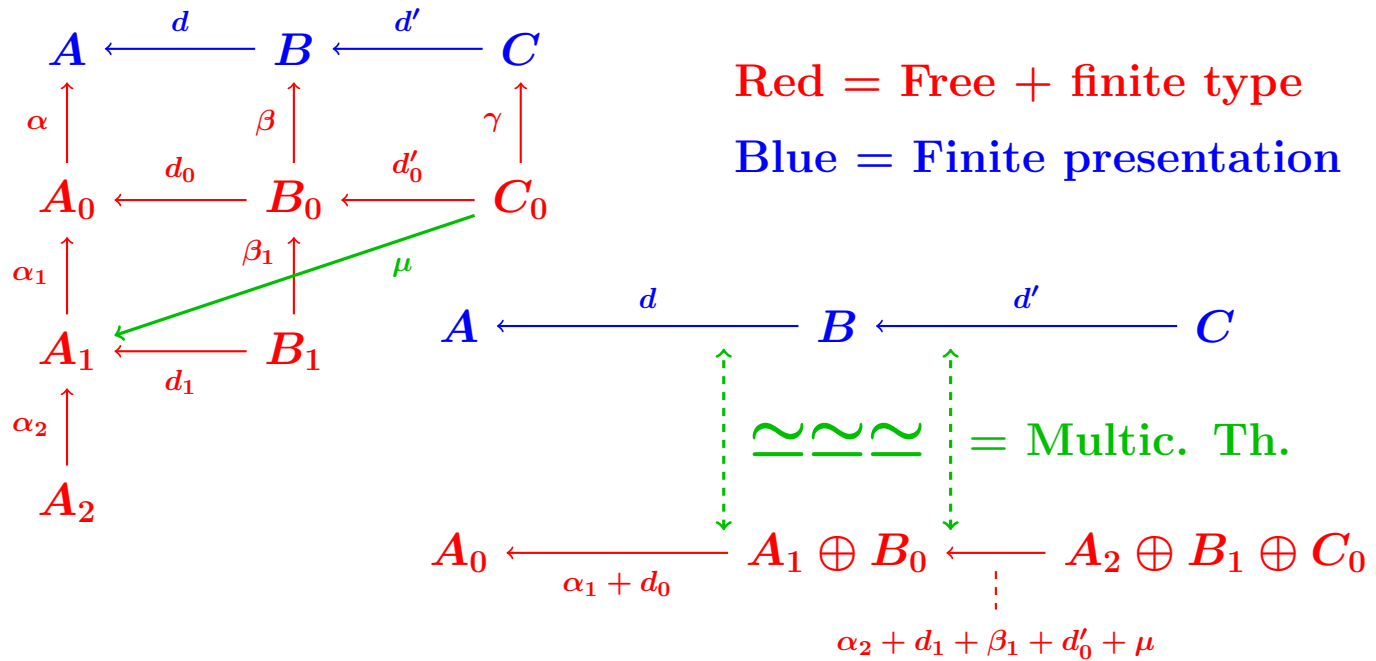
$$H' = (R^\delta, R^\gamma \oplus R^\varepsilon, \psi \oplus \phi)$$

Theorem: An algorithm produces a SHP for a chain complex of modules with finite presentation.

Multicomplex theorem:



Proof.



QED

2. Locally effective module M .

The membership property of an arbitrary object $a \in M$ is decidable.

The module is discrete:

equality between objects is decidable.

Ordinary computations can be executed.

If $\alpha \in R$ and $a, b \in M$, algorithms compute $a + b$ and αa .

The global structure of M is unknown.

In particular M maybe is not of finite type.

No computation

involving the whole “knowledge” of M can be done.

3. Fuzzy module M .

The membership property of an arbitrary object $a \in? M$ is decidable.

The module is not necessarily discrete:

equality between objects is in general undecidable.

Ordinary computations can be executed.

If $\alpha \in R$ and $a, b \in M$, algorithms compute $a + b$ and αa .

The global structure of M is unknown.

In particular M maybe is not of finite type.

No computation

involving the whole “knowledge” of M can be done.

Typical example of **fuzzy module**.

Given a **chain complex**:

$$\cdots \leftarrow C_{p-1} \xleftarrow{d_p} C_p \xleftarrow{d_{p+1}} C_{p+1} \leftarrow \cdots$$

made of **locally** effective modules.

Then $H_p = \text{Ker } d_p / \text{Im } d_{p+1}$ is a **fuzzy module**.

The element type for H_p is Z_p :

a **homology class** is implemented as a **cycle**.

C_{p-1} **discrete** \Rightarrow Membership to H_p **decidable**.

C_{p+1} **not of finite type**

\Rightarrow Membership to $\text{Im } d_{p+1}$ **undecidable**.

$\Rightarrow H_p$ **not discrete**.

The END

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Homology in dimension 6 :

Component Z/12Z

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